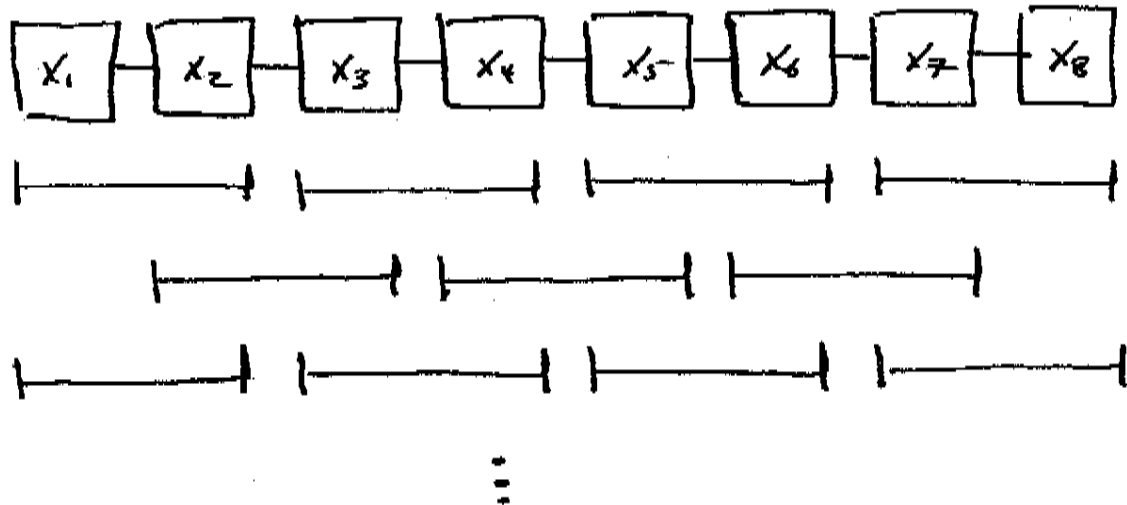


Mon 3/15/04
Michael Bender Lecturing

Sorting on 1 and 2D Arrays

Linear array: odd-even transposition sort



~~This odd-even trans sort runs in N steps (with 1 of opt).~~

Def: Oblivious comparison-exchange alg. Comparisons prespecified.

Independent of results of prev comparisons.

<<eg. quicksort not oblivious>>

Thm: If an obliv comparison-exchange alg sorts all 2^N sequences of 0's and 1's, it sorts all sequences of arbitrary #s

3/15/04

Proof: In 2 parts:

(1) Let f be monotonically increasing function.

$$\begin{aligned} \text{Then } \min\{f(x), f(y)\} &= f(\min\{x, y\}) \\ \max\{f(x), f(y)\} &= f(\max\{x, y\}). \end{aligned}$$

By induction on time steps, if alg transforms
 $\langle a_1, a_2, \dots, a_n \rangle \rightarrow \langle b_1, b_2, \dots, b_n \rangle$,

then it transforms

$$\langle f(a_1), f(a_2), \dots, f(a_n) \rangle \rightarrow \langle f(b_1), f(b_2), \dots, f(b_n) \rangle$$

«see CLR»

(2) Suppose false. i.e., network sorts all 0-1 seq,
 but $\exists \langle a_1, a_2, \dots, a_n \rangle$ st $a_i < a_j$, but a_i comes
 after a_j in output.

$$\text{Define } f(x) = \begin{cases} 0, & \text{if } x \leq a_i \\ 1, & \text{if } x > a_i \end{cases}$$

But network fails to sort $\langle f(a_1), f(a_2), \dots, f(a_n) \rangle$

Contradiction. ■

«threshold induction»

\Rightarrow Need only construct 0-1 sorting algo!

3/15/04

Thm: Odd-even transposition sort runs in N steps

(with \pm of OPT).

\ll Result less interesting than proof method \Rightarrow

Pf: Consider movement of rightmost \pm .

1st step: may not move $\underline{\pm}$

During subsequent steps, moves forward.

\Rightarrow cannot block other \pm s.

\Rightarrow k 'th leftmost \pm begins moving by step $k+1$.
Must reach position $N-k+1$.

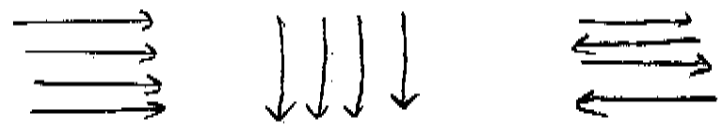
\Rightarrow All elmts in final position by time N

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Sorting on 2D Grid

Lower bounds: $2\sqrt{N}-2$ (diameter)
 $\sqrt{2N}$ (bisection)

Natural Grid Orders:

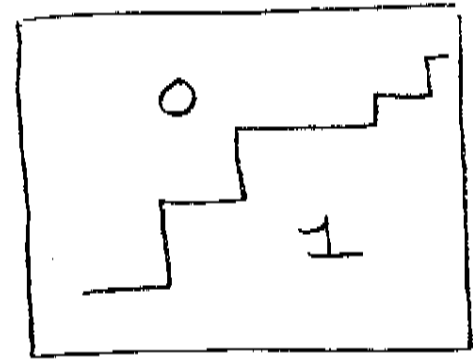


"Broken" Alg:

Repeat:

(1) ↓↓↓↓

(2) ⇒⇒⇒



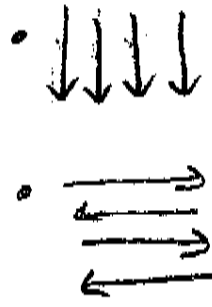
Doesn't yield unique order:

$$\begin{matrix} 0 & 1 \\ 1 & 0 \end{matrix} \Rightarrow \left\{ \begin{array}{l} 00 \quad \downarrow \Rightarrow \\ 11 \quad \downarrow \Rightarrow \\ 01 \quad \Rightarrow \downarrow \end{array} \right.$$

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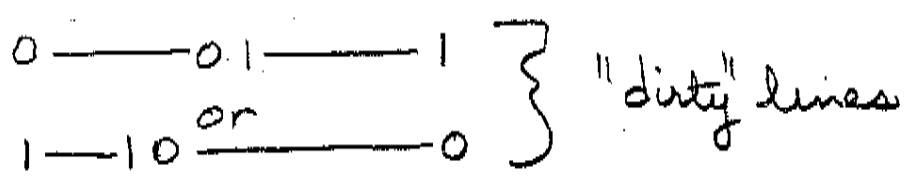
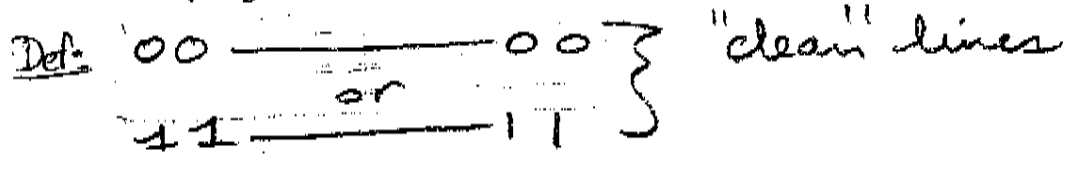
Shearsort

Repeat



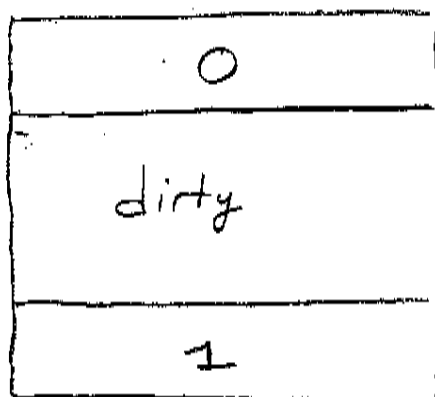
Thm: Shearsort produces unique sorting order after time $O(N \lg N)$. Eg, $O(\lg N)$ phases sufficient to sort.

Pf: Apply 0-1 lemma.

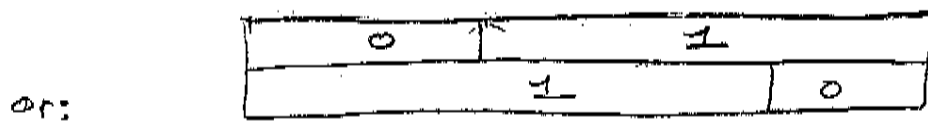
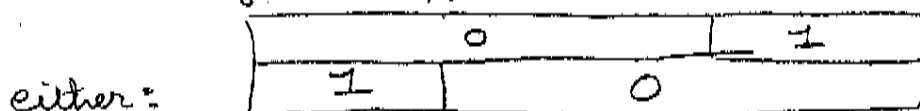


Claim: After each phase, # dirty rows decreases by at least half.

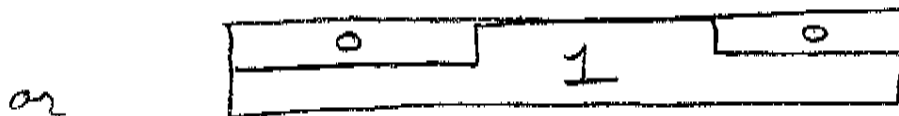
Grid has 3 regions:



Divide dirty into pairs of rows:



After sorting columns:

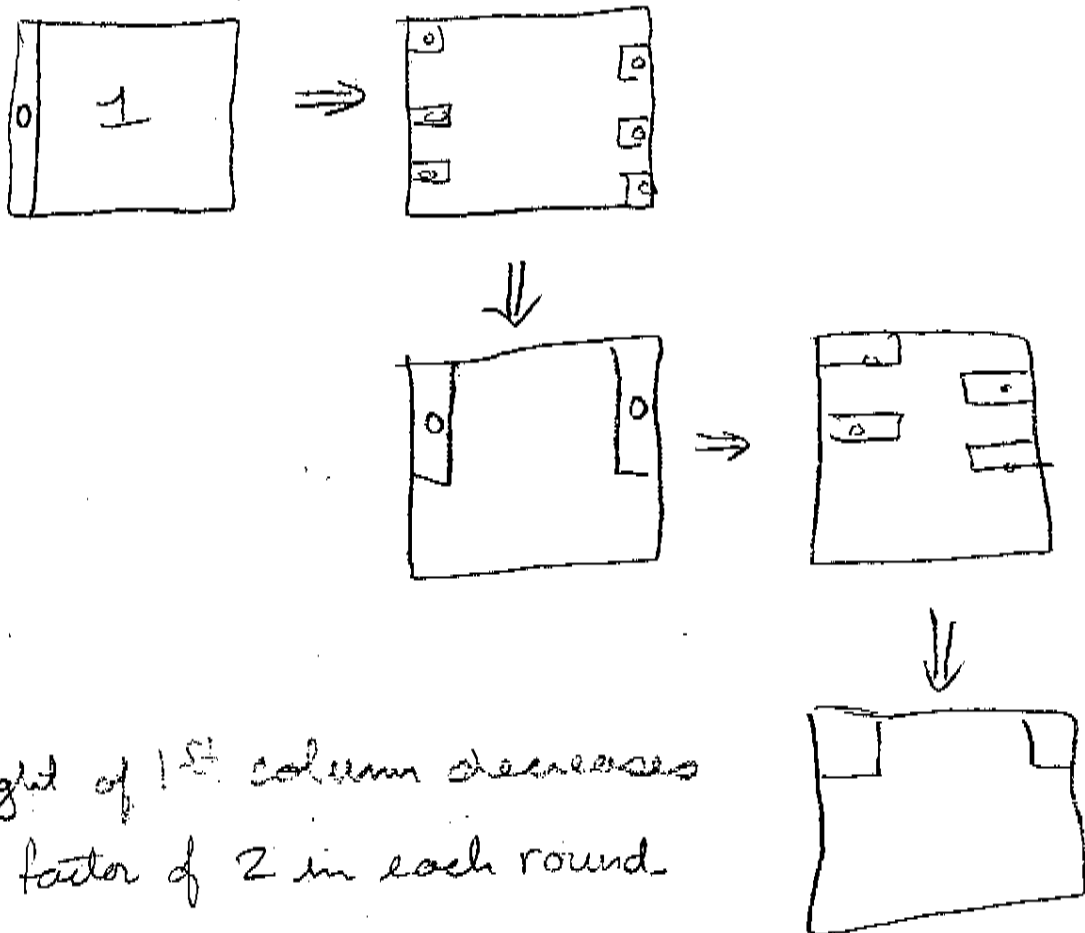


⇒ dirty region decreases by $\geq 1/2$.

⇒ after $\lg N$ phases $[\Theta(\sqrt{N} \lg N) \text{ time}]$
all sorted.

Lemma: Shear sort runs in $\Omega(\lg N)$ phases.

Pf: Bad example 0's in 1st column.



height of 1st column decreases
by factor of 2 in each round

Average Case:

Substitute 0's for \sqrt{N} smallest elms

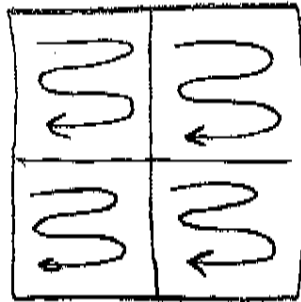
- row sort first: $E[\# \text{0's in 1}^{\text{st}} \text{ column}] = \Theta(N)$.

- column sort first: not true.
best LB = $\Omega(N \lg \lg N)$.

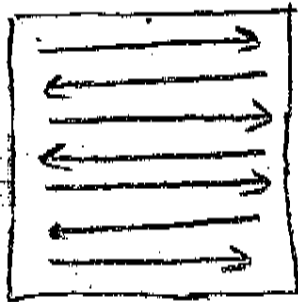
$O(\sqrt{N})$ Algorithm ($\leq 8\sqrt{N}$)

Assume \sqrt{N} is power of 2

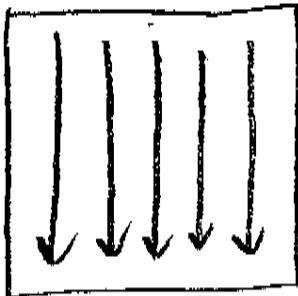
(1) Recursively sort each quadrant in snake order.



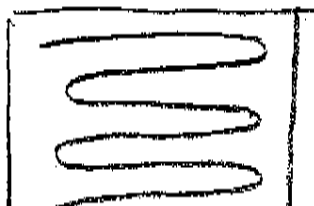
(2) Sort rows in alternate order



(3) Sort columns



(4) Do $2\sqrt{N}$ steps of 1D odd-even transposition on overall snake order.



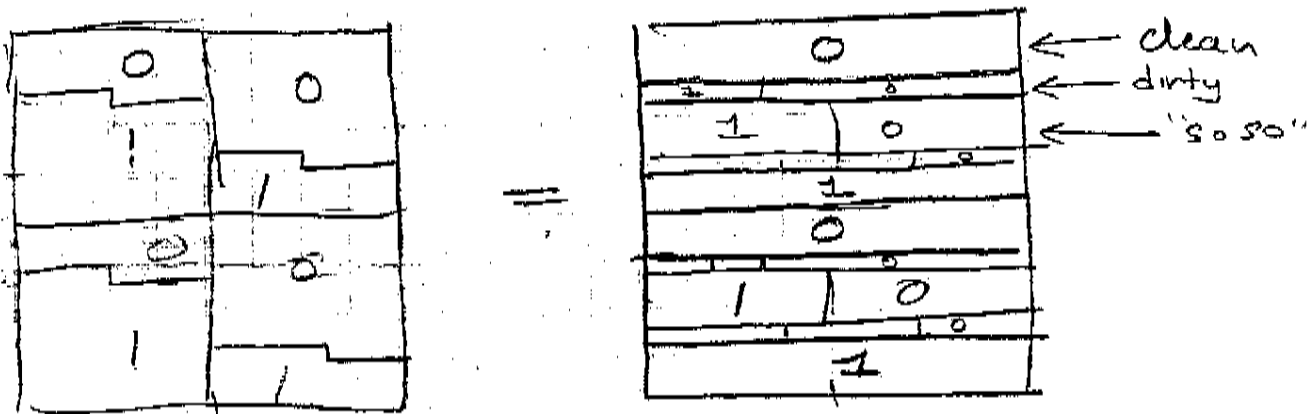
Running time:

$$T(N) \leq T(N/4) + \sqrt{N} + \sqrt{N} + 2\sqrt{N}$$

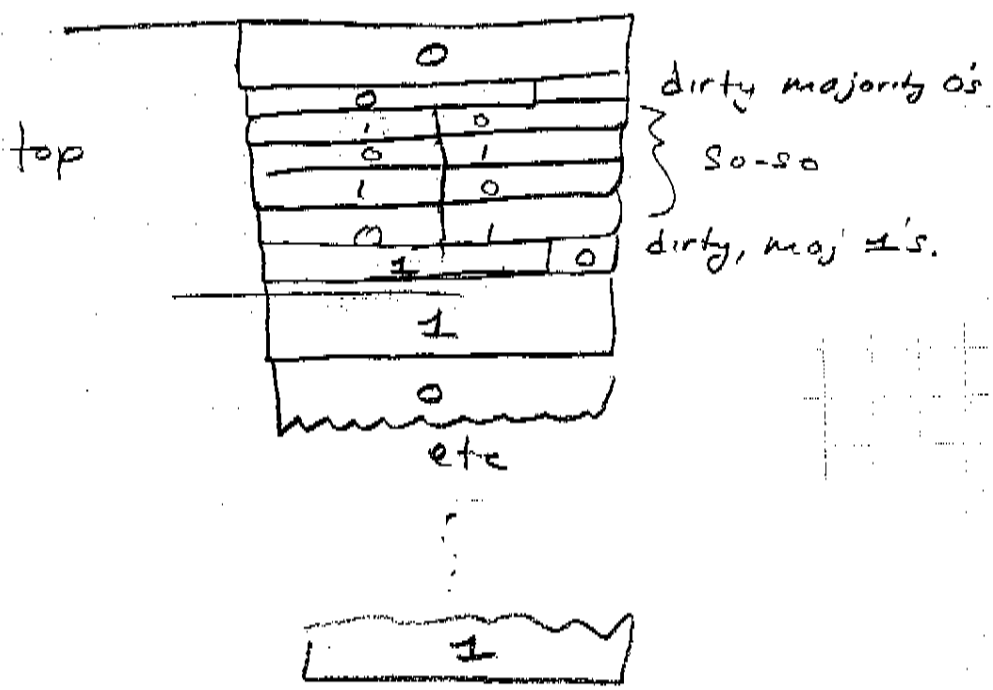
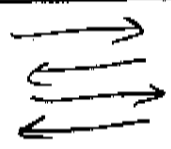
$$\leq 8\sqrt{N}$$

Proof of Correctness:

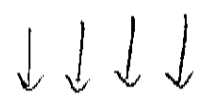
Phase 1 In each quadrant at most one of rows is dirty and rest are clean:



phase 2



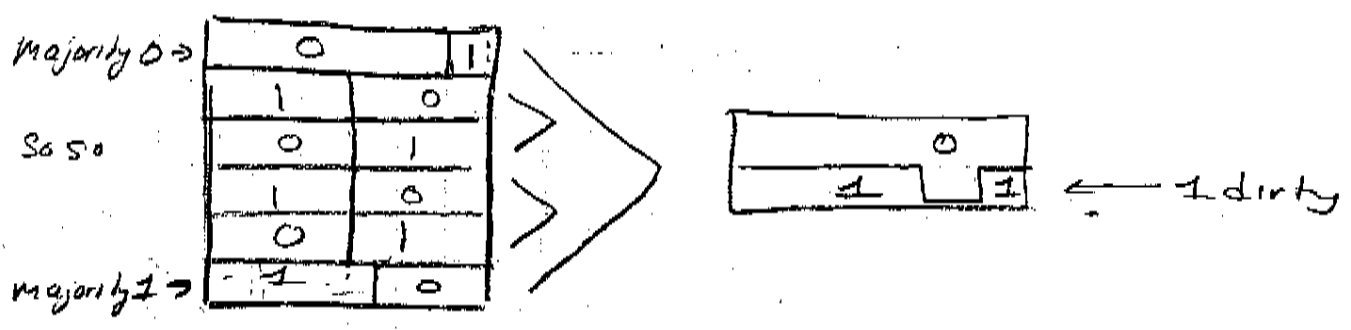
phase 3



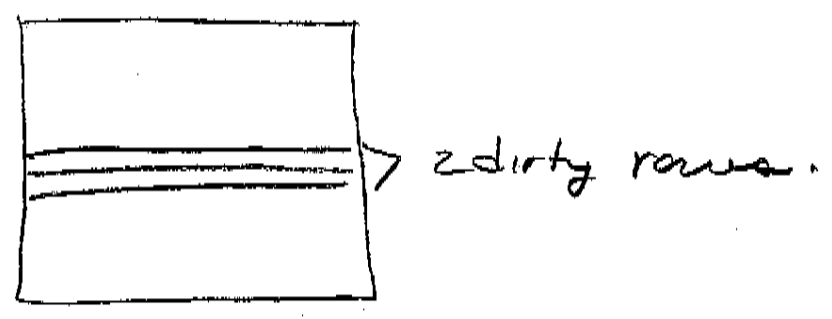
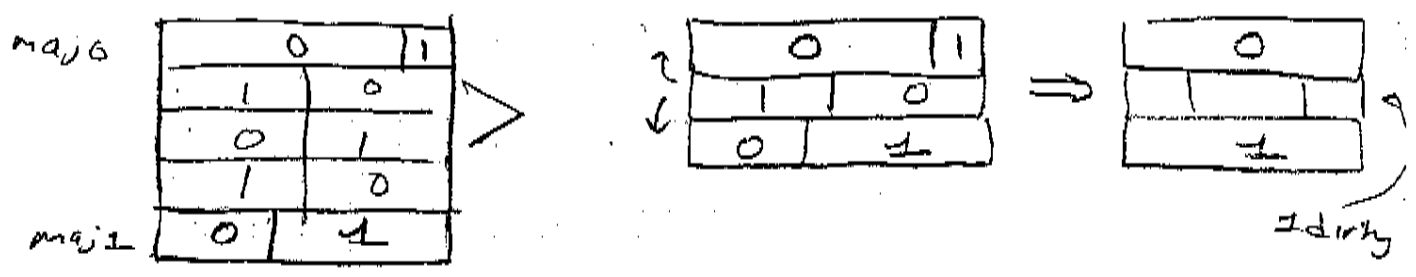
clm: ≤ 1 dirty row in top half &
" " " " bottom half.

<< gilding the lily - we already have a $\Theta(\sqrt{n})$ alg >>

Pf: Case 1: # so-so rows even:



Case 2: # so-so rows odd:



\Rightarrow Sorted after phase 4