## Problem Set 9

This problem set is due Wednesday, November 16 at noon.

1. Prove that the Steiner tree spanner as constructed in the lecture actually contains a Steiner tree of the given terminals of weight at most  $(1 + c\epsilon)$ OPT (where c is an absolute constant).

**Hint:** Recall that the structure theorem guarantees that we can transform any forest inside a brick into another forest of slightly larger weight that includes E and W and that preserves connectivity among the boundary and has at most  $\alpha$  joining vertices, where  $\alpha = 2c'\kappa\epsilon^{-2.5} = O(\epsilon^{-5.5})$ . However, we do not know the joining vertices; hence we greedily selected  $\theta = 2\alpha\epsilon^{-1} \cdot \ell(\text{MG})$  portals around each brick.

- 2. Let G be a plane graph, P be a  $(1 + \epsilon)$ -short path in G and T be binary tree in G with no vertices of degree 2, rooted at r, and having all leaves on P.
  - (a) How many joining vertices does T have with P if its depth is at most  $\epsilon^{-1}$ ?
  - (b) Show that if the depth of T is at least  $\epsilon^{-1}$  then there is a level k such that the total weight of all edges of T at level k is small and give an upper bound on that weight.
  - (c) Show how to transform T into another tree T' of weight at most  $(1+2\epsilon)\ell(T)$  that spans r and  $V(T) \cap V(P)$  and has  $f(\epsilon)$  joining vertices with P. Give an upper bound on  $f(\epsilon)$ .