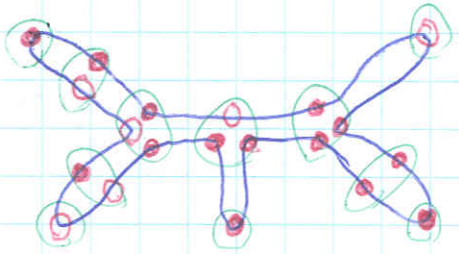
State-of-the-art for gluings:

	Enumeration alg./ # gluings	Decision algorithm
General	$2^{\Theta(n)}$	$2^{O(n)}$
Edge-to-edge gluings	$2^{\Theta(n)}$	$n^{\Theta(1)}$] L13.4
Bounded-sharpness polyg. (e.g. convex)	$n^{\Theta(1)}$	$n^{O(1)}$

General upper bound on # gluings (combinatorially distinct)

- saw $2^{\Omega(n)}$ lower bound [L13.3]
- $2^{O(n)}$:
  - each leaf in gluing tree is a polygon vertex or a fold point
  - $\Rightarrow O(n)$ leaves
  - internal nodes (junctions) have ≥ 1 polygon vertex
  - $\Rightarrow O(n)$ nodes in gluing tree
  - $2^{\Theta(n)}$ unlabeled trees on $O(n)$ nodes
  - $2^{O(n)}$ labelings of whether each point of a junction is a vertex or along an edge.
  - $O(n)$ choices for first vertex
  - $\Rightarrow 2^{O(n)} \times 2^{O(n)} \times O(n) = 2^{O(n)}$ comb. distinct gluing trees



● = vertex
○ = point along an edge
◯ = node of gluing tree