

Admin:

Pset #4 due Friday

Projects: meet with TA's this week

Today:

Skype call with Danny Weitzner
re projects & legalities

Digital Signature Standard

"Gap groups"

Bilinear maps

BLS digital signatures

IBE (if time)



PEOPLE

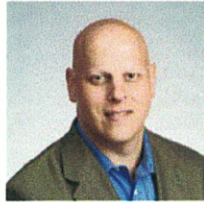
Principal Investigators

All Members

Student Spotlights

Home » People » Daniel Weitzner

DANIEL WEITZNER



Position: Principal Research Scientist
Office: [32-G516](#)
Phone: +1 (617) 253-8036
Email: djweitzner@csail.mit.edu
Areas of Study: Privacy, Internet Policy,
Web Architecture
Last Update: February 19, 2014
[Download vCard](#)

AWARDS

IAPP: Privacy Leadership Award (2013)
Newsweek/Daily Beast: Digital Power Index
(2012)

submit new awards here: [Award
Registration Form](#)

BIOGRAPHY

Daniel Weitzner is the Director of the MIT CSAIL Decentralized Information Group and teaches Internet public policy in MIT's Computer Science Department. His research includes development of accountable systems architectures to enable the Web to be more responsive to policy requirements.

From 2011-2012, Weitzner was the United States Deputy Chief Technology Officer for Internet Policy in the White House. He led initiatives on privacy, cybersecurity, Internet copyright, and trade policies promoting the free flow of information. He was responsible for the Obama Administration's [Consumer Privacy Bill of Rights](#) and the [OECD Internet Policymaking Principles](#).

Weitzner has been a leader in the development of Internet public policy from its inception, making fundamental contributions to the successful fight for strong online free expression protection in the United States Supreme Court, and for laws that control government surveillance of email and web browsing data.

Weitzner is a founder of the Center for Democracy and Technology, led the World Wide Web Consortium's public policy activities, and was Deputy Policy Director of the Electronic Frontier Foundation. In 2012 he was named to the Newsweek/Daily Beast Digital Power Index as a top 'Navigator' of global Internet public policy and in 2013 he received the International Association of Privacy

Professional's Leadership Award.

Digital Signature Standard (DSS - NIST 1991)

Public parameters (same for everyone):

q prime

$|q| = 160$ bits

$p = nq + 1$ prime

$|p| = 1024$ bits

g_0 generates Z_p^*

$g = g_0^n$ generates subgroup G_q of Z_p^* of order q

Keygen:

$x \xleftarrow{R} Z_q$ SK $|x| = 160$ bits

$y \leftarrow g^x \pmod{p}$ PK $|y| = 1024$ bits

Sign(m):

$k \xleftarrow{R} Z_q^*$ (i.e. $1 \leq k < q$)

$r = (g^k \pmod{p}) \pmod{q}$ $|r| = 160$ bits

$m = h(M)$

$s = (m + rx) / k \pmod{q}$ $|s| = 160$ bits

redo if $r=0$ or $s=0$

$\sigma(M) = (r, s)$

Note: if k is reused for different messages m , one could solve for x so it is not secure.

If k is reused for the same m , we obtain the same signature so this is not a problem. If k is different for the same m , it should be random and unknown (any known relation between the two k -s allows to solve for x)

Bottomline: All of the above are enforced by k chosen at random from Z_q^* for large enough q

Verify:

Check $0 < r < q$ & $0 < s < q$

Check $y^{r/s} g^{m/s} \pmod{p} \pmod{q} = r$

where $m = h(M)$

Correctness:

$$g^{(rx+m)/s} \stackrel{?}{=} r \pmod{p} \pmod{q}$$

$$\equiv g^k = r \pmod{p} \pmod{q} \quad \checkmark$$

As it stands, existentially forgeable for $h = \text{identity}$.

Provably secure (as with Modified El Gamal)

if we replace $m = h(m)$ by $m = h(M || r)$, as before.

Note: As with El Gamal, secrecy & uniqueness of k
is essential to security.

"Gap group" is one in which

• DDH is easy ("Decision Diffie Hellman")

[Recall: given (g, g^a, g^b, g^c) , to
decide if $ab = c \pmod{\text{order}(g)}$]

but • CDH is hard ("Computational Diffie Hellman")

[Recall: given (g, g^a, g^b) , to
compute g^{ab}]

(Note that CDH easy \Rightarrow DDH easy)

This difference in difficulty between DDH ("easy")
and CDH ("hard") forms a "gap".

- How can one construct a "gap group"?
- What good would that be?