TOPIC: 6.857 - Computer & Network Security

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Admin:

Pset #4 out.

Quiz in-class Wed 4/16. Open notes (No laptops or books)

Projects: Next week: meet with TA to review progress.

Presentation schedule out soon (let us know if conflict.)

Today:

- ☐ Digital signatures
- ☐ Security defn for digital signatures
- ☐ Hash & sign
- ☐ RSA - PKCS
- ☐ RSA - PSS
- ☐ El Gamal dig. sigs.
- ☐ DSA - (NIST standard)

Digital Signatures (compare "electronic signature", "cryptographic signature")

- Invented by Diffie & Hellman in 1976
  ("New Directions in Cryptography")
- First implementation: RSA (1977)
- Initial idea: switch PK/SK
  (enc with secret key $\Rightarrow$ signature)
  (if PK decrypts it & looks ok then sig ok??)

Current way of describing digital signatures

- $\text{Keygen}(1^\lambda) \longrightarrow (\underline{PK}, \underline{SK})$
  verification key → PK; signing key → SK
- $\text{Sign}(SK, m) \rightarrow \underbrace{\sigma_{SK}(m)}_{\text{signature}}$ [may be randomized]
- $\text{Verify}(PK, m, \sigma) =$ True/False (accept/reject)

Correctness:

$(\forall m)\ \text{Verify}(PK, m, \text{Sign}(SK, m)) = \text{True}$

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Security of digital signature scheme:

Def: (weak) existential unforgeability under adaptive chosen message attack.

① Challenger obtains (PK, SK) from Keygen(1^λ)
Challenger sends PK to Adversary

② Adversary obtains signatures to a sequence

$$m_1, m_2, \ldots, m_q$$

of messages of his choice. Here $q = \text{poly}(\lambda)$, and m_i may depend on signatures to $m_1, m_2, \ldots, m_{i-1}$.

Let $\sigma_i = \text{Sign}(SK, m_i)$.

③ Adversary outputs pair (m, σ_*)

Adversary <u>wins</u> if $\text{Verify}(PK, m, \sigma_*) = \text{True}$ and $m \notin \{m_1, m_2, \ldots, m_q\}$

Scheme is secure (i.e. weakly existentially unforgeable under adaptive chosen message attack) if

$$\text{Prob}[\text{Adv wins}] = \text{negligible}$$

Scheme is <u>strongly</u> secure if adversary can't even produce new signature for a message that was previously signed for him.

I.e. Adv wins if Verify(PK, m, σ_*) = True and $(m, \sigma_*) \notin \{(m_1, \sigma_1), (m_2, \sigma_2), \ldots, (m_q, \sigma_q)\}$.

Digital signatures

- Def of digital signature scheme
- Def of weak/strong existential unforgeability under adaptive chosen message attack.

$\Big\}$ see notes from last lecture

Hash & Sign:

For efficiency reasons, usually best to sign cryptographic hash $h(M)$ of message, rather than signing M. Modular exponentiations are slow compared to (say) SHA-256.

Hash function h should be collision-resistant.

Signing with RSA — PKCS

- PKCS = "Public key cryptography standard"
  (early industry standard)
- Hash & sign method. Let H be C.R. hash fn.
- Given message M to sign:

  Let $m = H(M)$

  Define $pad(m) =$

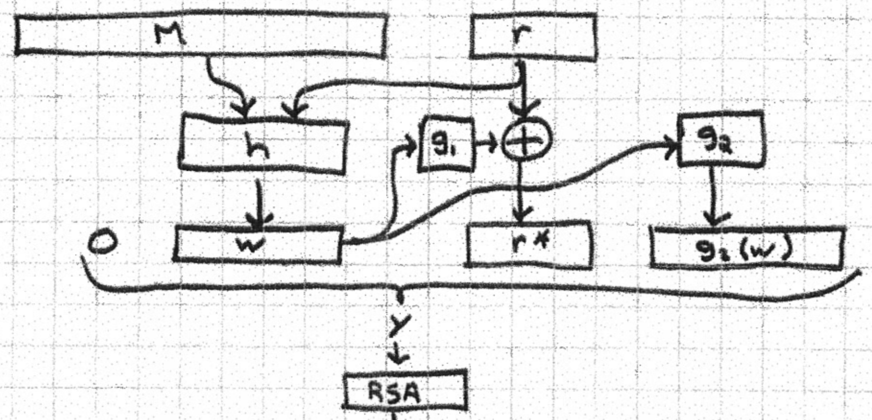
  $$0x\ 00\ 01\ FF\ FF \ldots FF\ 00 \,\|\, \text{hash-name} \,\|\, m$$

  where # FF bytes enough to make $|pad(m)| = |n|$ in bytes.

  where hash-name is given in ASN.1 syntax (ugly!)
- Seems secure, but no proofs (even assuming H is CR and RSA is hard to invert)
- $\sigma(M) = (pad(m))^d \pmod n$

PSS - Probabilistic Signature Scheme [Bellare & Rogaway 1996]

- RSA-based
- "Probabilistic" ≡ randomized [one M has <u>many</u> sigs]



$$\sigma(m) = y^d \pmod n$$

Sign(M): $r \xleftarrow{R} \{0,1\}^{k_0}$

$w \longleftarrow h(M \| r)$ $\qquad |w| = k_1$

$r^* \longleftarrow g_1(w) \oplus r$ $\qquad |r^*| = k_0$

$y \longleftarrow 0 \| w \| r^* \| g_2(w)$ $\qquad |y| = |n|$

output $\sigma(M) = y^d \pmod n$

Verify(M, σ): $y \leftarrow \sigma^e \pmod n$

Parse y as $b \| w \| r^* \| \gamma$

$r \leftarrow r^* \oplus g_1(w)$

return True iff $b = 0$ & $h(M \| r) = w$ & $g_2(w) = \gamma$

- We can model h, g_1, and g_2 as random oracles.

<u>Theorem:</u>

PSS is (weakly) existentially unforgeable against a chosen message attack in random oracle model if RSA is not invertible on random inputs.

El Gamal digital signatures

Public system parameters: prime p

generator g of $\mathbb{Z}_p^*$

Keygen: $x \xleftarrow{R} \{0,1,\ldots,p-2\}$ $\quad SK = x$

$y = g^x \pmod p$ $\quad PK = y$

Sign(M):

$m = \text{hash}(M)$ $\quad$ CR hash fn into $\mathbb{Z}_{p-1}$

$k \xleftarrow{R} \mathbb{Z}_{p-1}^*$ $\quad$ [$\gcd(k, p-1) = 1$]

$r = g^k$ $\quad$ [hard work is indep of M]

$s = \frac{(m - rx)}{k} \pmod{p-1}$

$\sigma(M) = (r, s)$

Verify(M, y, (r, s)):

Check that $0 < r < p$ $\quad$ (else reject)

Check that $y^r r^s = g^m \pmod p$

where $m = \text{hash}(M)$

Correctness of El Gamal signatures:

$$\underbrace{y^r r^s = g^{rx} g^{sk} = g^{rx+sk} \stackrel{?}{=} g^m \pmod{p}}_{\equiv}$$

$$rx + ks \stackrel{?}{=} m \pmod{p-1}$$

or $$s \stackrel{?}{=} \frac{(m - rx)}{k} \pmod{p-1}$$

(assuming $k \in \mathbb{Z}_{p-1}^*$) $\blacksquare$

Theorem: El Gamal signatures are existentially forgeable (without h, or h = identity (note: this is CR!))

Proof: Let $e \xleftarrow{R} Z_{p-1}$

$r \leftarrow g^e \cdot y \pmod p$

$s \leftarrow -r \pmod p$

Then (r,s) is valid El Gamal sig. for message $m = e \cdot s \pmod{p-1}$.

Check: $y^r r^s \stackrel{?}{=} g^m$

$g^{xr}(g^e y)^{-r} = g^{-er} = g^{es} = g^m$ ✓ ▨

But: It is easy to fix.

Modified El Gamal (Pointcheval & Stern 1996)

Sign (M): $k \xleftarrow{R} Z_p^*$

$r = g^k \pmod p$

$m = h(M \| r)$ $\Longleftarrow$ ***

$s = (m - rx)/k \pmod{p-1}$

$\sigma(M) = (r,s)$

Verify: check $0<r<p$ & $y^r r^s = g^m$ where $m = h(M\|r)$.

Theorem: Modified El Gamal is existentially unforgeable against adaptive chosen message attack, in ROM, assuming DLP is hard.

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Digital Signature Standard (DSS - NIST 1991)

Public parameters (same for everyone):

q prime $\qquad |q| = 160$ bits

$p = nq + 1$ prime $\qquad |p| = 1024$ bits

g_0 generates Z_p^*

$g = g_0^n$ generates subgroup G_q of Z_p^* of order q

Keygen:

$x \xleftarrow{R} Z_q$ $\qquad$ SK $\qquad |x| = 160$ bits

$y \leftarrow g^x \pmod p$ $\qquad$ PK $\qquad |y| = 1024$ bits

Sign(m):

Note: if k is reused for different messages m, one could solve for x so it is not secure. If k is reused for the same m, we obtain the same signature so this is not a problem. If k is different for the same m, it should be random and unknown (any known relation between the two k-s allows to solve for x)

Bottomline: All of the above are enforced by k chosen at random from Z_q^* for large enough q

$k \xleftarrow{R} Z_q^*$ $\qquad$ (i.e. $1 \le k < q$)

$r = (g^k \bmod p) \pmod q$ $\qquad |r| = 160$ bits

$m = h(M)$

$s = (m + rx)/k \pmod q$ $\qquad |s| = 160$ bits

redo if $r = 0$ or $s = 0$

$\sigma(M) = (r, s)$

Verify:

Check $0 < r < q$ & $0 < s < q$

Check $y^{r/s} g^{m/s} \pmod{p} \pmod{q} = r$

where $m = h(M)$

Correctness:

$$g^{(rx+m)/s} \stackrel{?}{=} r \pmod{p} \pmod{q}$$

$$\equiv \quad g^k = r \pmod{p} \pmod{q} \quad \checkmark$$

As it stands, existentially forgeable for h = identity.

Provably secure (as with Modified El Gamal)

if we replace $m = h(M)$ by $m = h(M \| r)$, as before.

Note: As with El Gamal, secrecy & uniquness of k is essential to security.