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Admin:

- pset #1 grading:
- lecture notes for last week (#6 & #7) up soon

Note: RSA Conference; discussion of dual-ec-drbg

Project idea: "Security of routers"

Today:

Block ciphers:

DES (Data Encryption Standard)

AES (Advanced Encryption Standard)

Ideal Block Cipher

"Modes of Operations"
- ECB
- CTR
- CBC
- CFB

IND-CCA security def

UFE mode

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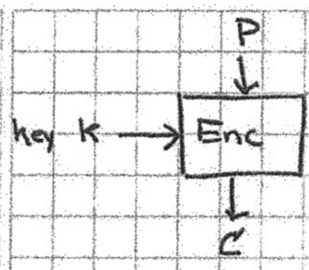
Block ciphers:



plaintext block

ciphertext block

fixed-length P, C, K

DES: $|P| = |C| = 64$ bits $|K| = 56$ bits

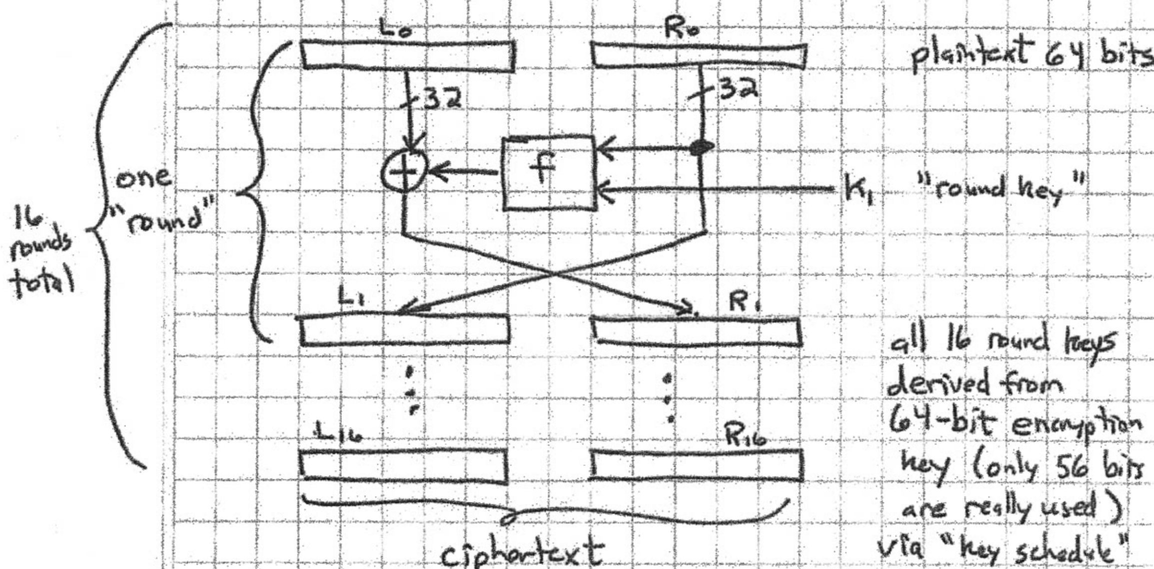
AES: $|P| = |C| = 128$ bits $|K| = 128, 192, 256$ bits

Use a "mode of operation" to handle variable-length input.

DES

"Data Encryption Standard"
Standardized in 1976. Now deprecated in favor of AES.

"Feistel structure":

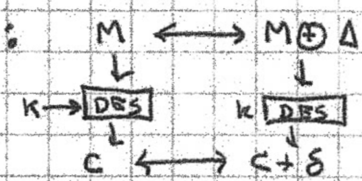


<u>Note:</u> Invertible for <u>any</u> f and <u>any</u> key schedule.

f uses 8 "S-boxes" mapping 6-bits ⟹ 4 bits nonlinearly.

Key is <u>too short</u>! (Breakable now quite easily by brute-force)

Subject to <u>differential attacks</u>:



2^{47} chosen pairs (Biham/Shamir)

Subject to <u>linear attacks</u>:

e.g. if $M_3 \oplus M_{15} \oplus C_2 \oplus K_{14} = 0$ (eqn on bits)
with prob $p = 1/2 + \varepsilon$

then need $1/\varepsilon^2$ samples to break (Matsui, 2^{43} PT/CT pairs)

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AES

"Advanced Encryption Standard" (U.S. govt)

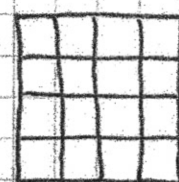
Replaces DES

AES "contest" 1997–1999:
15 algorithms submitted: RC6, Mars, Twofish, Rijndael, ...
winner = Rijndael (by Joan Daemen & Vincent Rijmen, (Belgians))

Specs: 128-bit plaintext/ciphertext blocks
128, 192, or 256-bit key
10, 12, or 14 rounds (dep. on key length)

Byte-oriented design (some math done in Galois field $GF(2^8)$)

View input as 4x4 byte array:



$4 \times 4 \times 8 = 128$

For version with 128-bit keys, 10 rounds:

- Derive 11 "round keys", each 128 bits (4x4x byte)
- In each round:
  1. XOR round key
  2. Substitute bytes (lookup table)
  3. Rotate rows (by different amts)
  4. Mix each column (by linear opn)
- Output final state

See readings for details.
There are very fast implementations. Also Intel has put supporting hardware into its CPU's.

Security: Good; perhaps # rounds should be a bit larger...

For practical purposes, can treat AES as <u>ideal block cipher</u>:

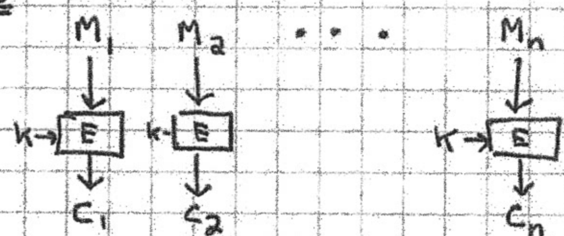
[For each key, mapping Enc(K, •) is a random independent permutation of $\{0,1\}^{128}$ to itself.

Modes of Operation:

How to encrypt variable-length messages? (using AES)

- "ECB" = "Electronic code book"
- "CTR" = "Counter mode"
- "CBC" = "Cipher-block chaining" (& CBC-MAC)
- "CFB" = "Cipher feedback"
- ... (others...)

ECB:



← divide data into b-bit blocks, where b = input block size

To handle data that is not a multiple of b bits in length:

- Append a "1" bit (<u>always</u>)
- Append enough "0" bits to make length a multiple of b bits.

This gives <u>invertible</u> (1-1) "padding" operation.
Pad before encryption; unpad after decryption.

ECB preserves many patterns: repeated message blocks
⇒ repeated ciphertext blocks

ECB really only good for encrypting <u>random</u> data (e.g. keys)