

6.849

Class 2

Feb. 15, 2017

NP-hard  $\approx$  "computationally intractable"

- if a problem is NP-hard, then there's no efficient algorithm to solve it unless  $P=NP$   
(famous unsolved problem, worth \$1M+)

-  $P \neq NP \approx$  "computers can't simulate lucky guessing, say heads vs. tails, without trying both options"

$\hookrightarrow$  almost everyone believes it

Examples of NP-hard problems:

- SAT: given Boolean formula  $(x \text{ OR NOT } y) \text{ AND } z$   
can you set the variables  $x, y, z$  true/false so that formula is true?

- Partition: given  $n$  integers, can you split them into two halves of equal sum?

(e.g. equalizing teams for a game)

- weakly NP-hard: hard only when integers are exponential in  $n$

(still polynomial number of bits)

- 3-Partition: given  $n$  integers, can you split them into  $n/3$  triples of equal sum?

- strongly NP-hard: hard even when integers are polynomial in  $n$

Approach: show e.g. Partition is easier / a special case of your problem: any Partition problem can be converted into a problem of your type  
 $\Rightarrow$  your problem is NP-hard too

### Simple example: Ruler Folding

- given 1D hinge pattern  
 crease pattern where all creases are optional
  - given length bound  $M$
  - decide: is there a flat folding of length  $\leq M$
- is weakly NP-hard:

- given Partition problem  $a_1, a_2, \dots, a_n$
- choose  $L > a_1 + a_2 + \dots + a_n$  "HUGE"
- hinge pattern:



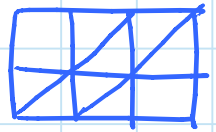
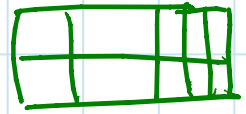
- target length  $M = 2L \Rightarrow$  forced folds
- intended folding:
- looking at travel in 1D:
  - at each hinge can crease  $\rightarrow$  change direction
  - or not  $\rightarrow$  same direction

$\Rightarrow$  can choose  $+a_i$  or  $-a_i$  for each  $i$

- must have  $\sum_i \pm a_i = 0$   
 i.e.  $\sum_i +a_i = -\sum_i -a_i$

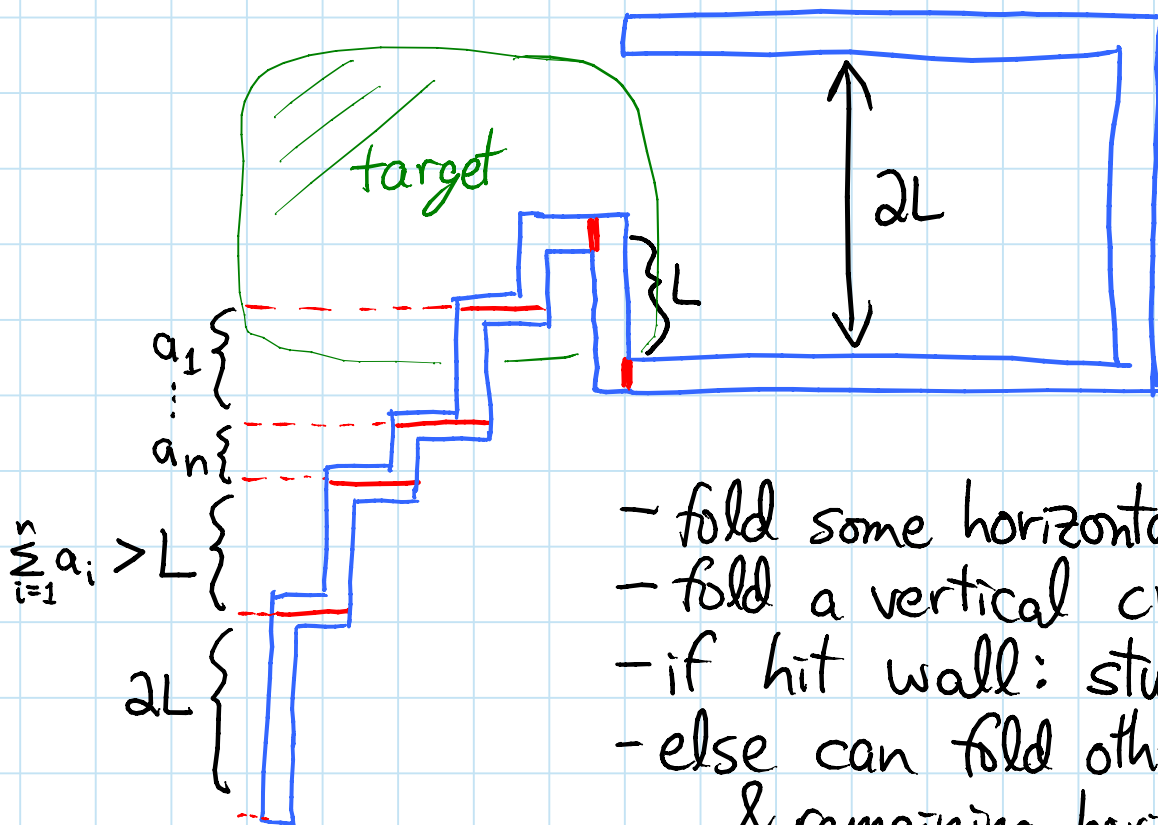
- YES to Partition  $\Leftrightarrow$  YES to Ruler Folding  $\square$

Simple folds: can given crease pattern be folded flat by sequence of simple folds?  
 - saw how to solve for 1D patterns & 2D orthogonal maps:



Weakly NP-hard if we add  $45^\circ$  diagonal creases or allow orthogonal paper

[Arkin, Bender, Demaine, Demaine, Mitchell, Sethia, Skiena 2000]  
 - reduction from Partition



- fold some horizontal creases
- fold a vertical crease
- if hit wall: stuck
- else can fold other vertical & remaining horiz. creases

□

Strongly NP-hard by reduction from 3-Partition  
 [Akitaya, Demaine, Ku 2016]