

- o Algorithm for testing flat foldability of 1D mountain-valley pattern (NEW) [C2]
 - - search (left to right) for segment that's crimpable  OR end foldable 
 - if none found: STOP ~ not flat foldable
 - else: do fold
 merge segments $x, y, z \rightarrow x-y+z$
 go back one segment (left of x)
 continue search

Correctness:

- doing fold changes foldability only of adjacent segments
 \Rightarrow enough to back up 1 step

Running time: $O(n)$

\uparrow # creases

$$\begin{aligned} - \# \text{ right steps} &= n + \underbrace{\# \text{ left steps}}_{=\# \text{ folds done}} \\ &\leq 2n \end{aligned}$$

Amortization: charge left steps to fold just done

- o Algorithm for testing flat foldability (NEW)
 - of single-vertex mountain-valley pattern
 - mark any start segment (angle)
 - - Search (clockwise) for crimpable segment
 - if found: do fold, merge segments
 - go back (ccw) one segment
 - (addition during class) if next to start segment:
move start one segment cw
 - continue search
 - if we "loop around" (return to start seg.):
 - foldable \Leftrightarrow just 2 M's or 2 V's
 - & 2 equal angles

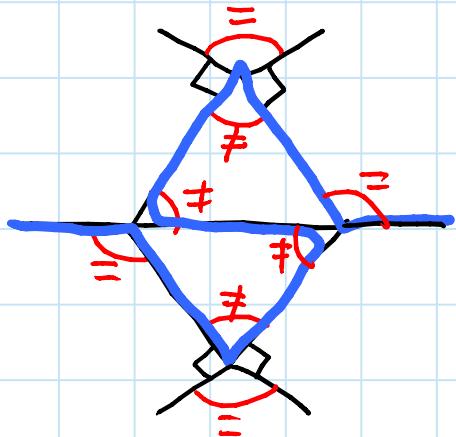
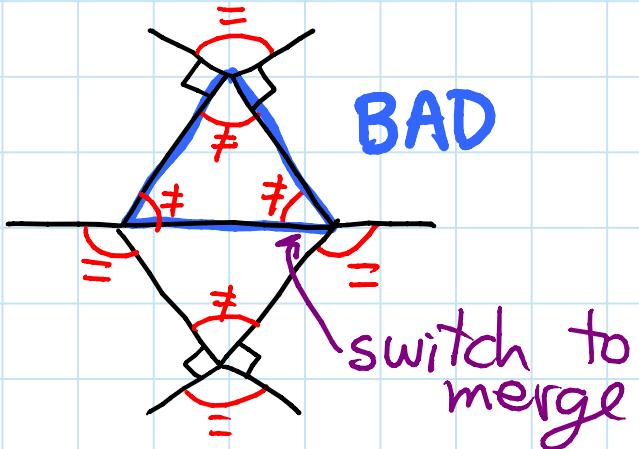
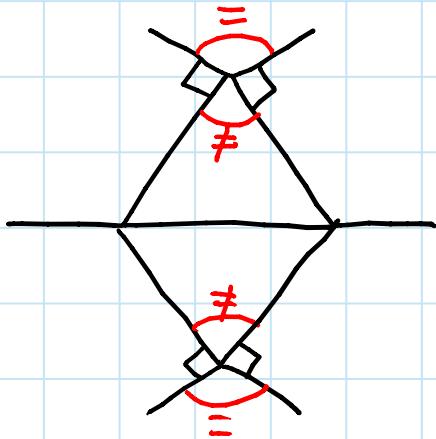
Correctness: maintain invariant that interval from start segment clockwise up to but not including current seg. is uncrimatable

Running time: $O(n)$

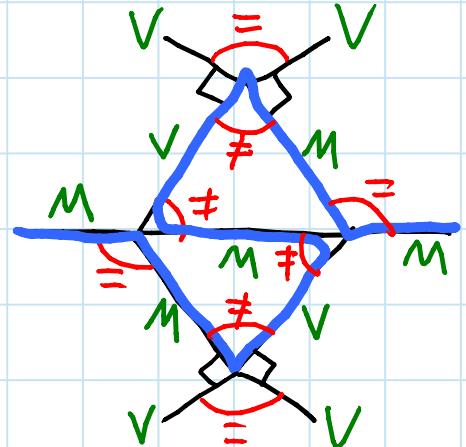
- interval shrinks (on either side) only when doing a fold

○ Local foldability examples

- generic case: unique pairing (crane)
- each cycle should have even # ≠'s
(i.e. parity(cycle) = parity(# ≠'s) as in L3)
- general case: merging



PATH,
NOT
CYCLE



○ T-shirt has $> 360^\circ$ of material at armpit
→ T-shirt folding video

○ Higher dimensions:

- definition: like  (add layer dim.)
- necessary condition & examples [Kawasaki 1989]
- graphics examples [Inoue, Itohara, Yajima, Kaino 2005]

OPEN: characterize single-vertex flat foldability

○ Why flat foldability?

- art e.g. tessellations
- practical storage e.g. airbag folding
- math e.g. Tachi's rigid-foldable quad mesh