

Configuration Space for Motion Planning

RSS Lecture 9

Wednesday, 5 March 2014

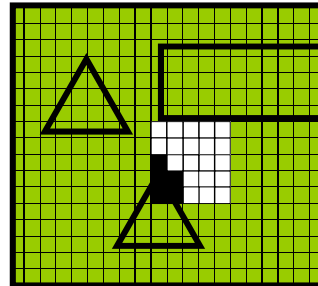
Prof. Seth Teller

Siegwart & Nourbakhsh S 6.2

(Thanks to Nancy Amato, Rod Brooks, Vijay Kumar,
and Daniela Rus for some of the figures)

Last Time

- Planning for point robots
 - Visibility graph method
 - Intermittent obstacle contact
- Ad hoc method of handling robots with positive area
 - Represent robot as a (2-DOF) disk
 - Discretize Cartesian space, conservatively
 - But: some feasible paths
- Today: “configuration space” methods
 - Reason directly in a space with dimension = #DOFs
 - Transform there; solve problem; transform back

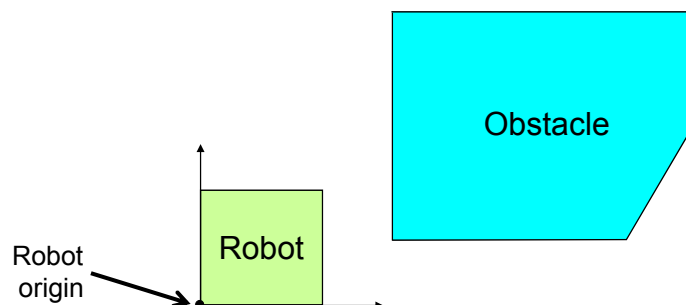


Today

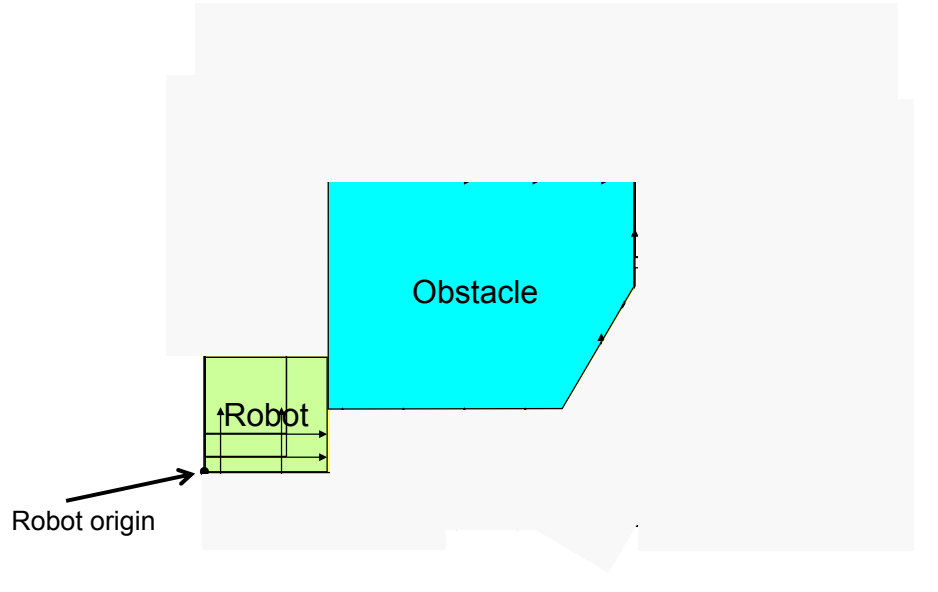
- Configuration space
 - Intuition
 - Preliminaries
 - Minkowski sums
 - Convexity, convex hulls
 - Definition
 - Construction
- Rigid (low-DOF) motion planning
 - Deterministic methods
- Articulated (high-DOF) motion planning
 - Randomized methods

Intuition

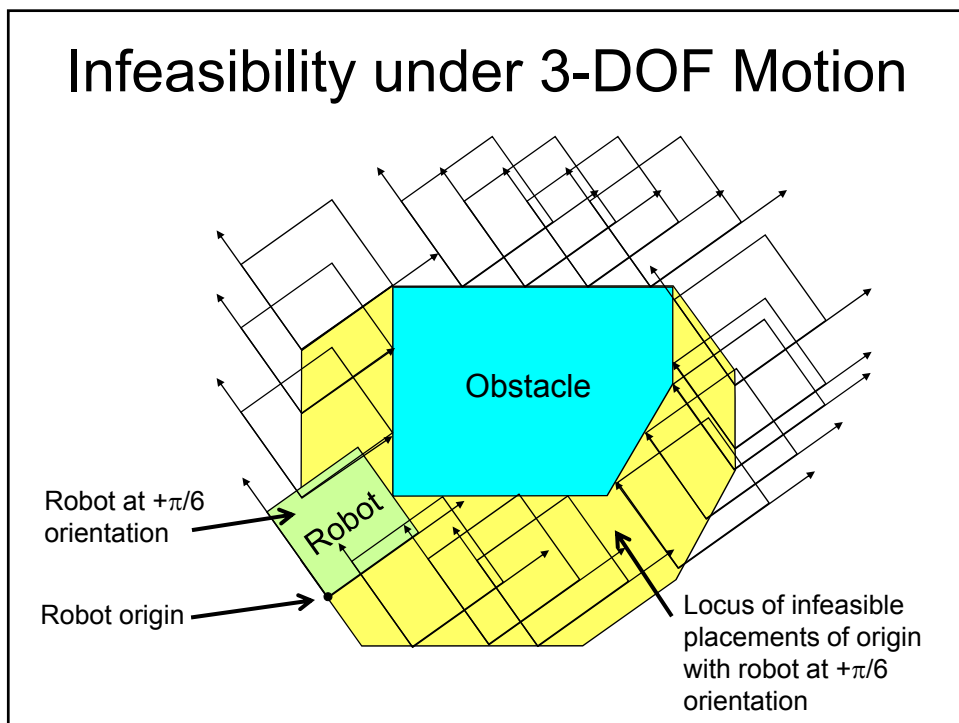
- Suppose robot can move only by translating in 2D
- How can it move in the presence of an *obstacle*?
 - Represent robot by its origin (how many DOFs?)
 - How to describe *infeasible placements* of robot origin?



Infeasibility Under Translation

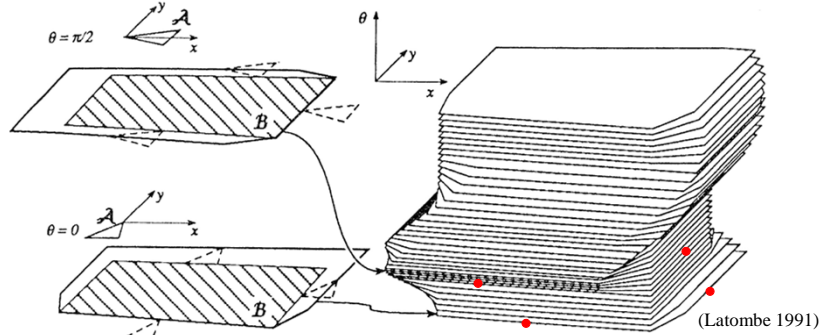


Infeasibility under 3-DOF Motion



Configuration Space

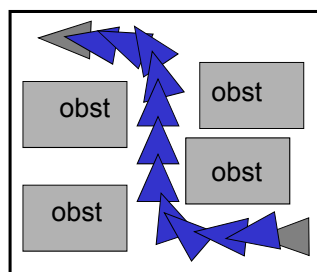
For a robot with k total motion DOFs, C-space is a coordinate system with *one dimension per DOF*



In C-space, a robot “pose” is simply
... and a workspace obstacle is a

Motion Planning Transformation

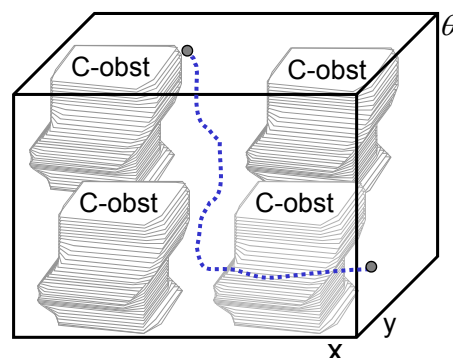
Workspace
(x, y)



▲ Robot

Path is hard to express

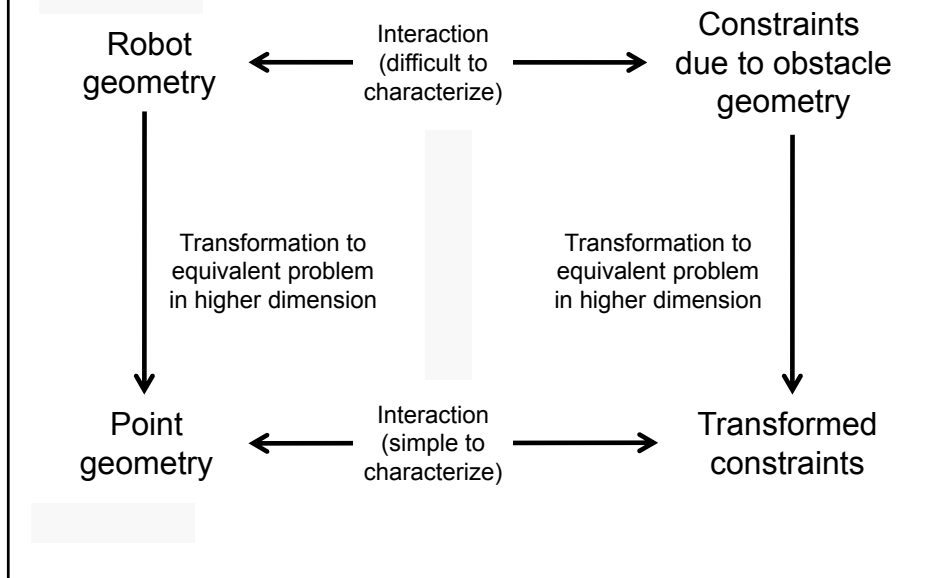
C-space
(x, y, θ)



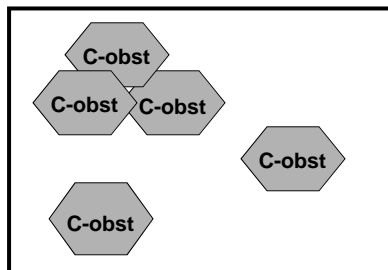
● Robot

Path is just a space curve

Configuration Space Idea



C-space Summary, Examples



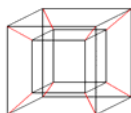
- Define space with one dimension per DOF of robot motion / pose
- Map robot to a *point* in this space
- C-space = *all* robot configurations
- C-obstacle = locus of *infeasible configurations* due to obstacle

Some example configuration spaces:

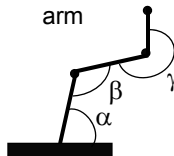
Translation + rotation in 2D



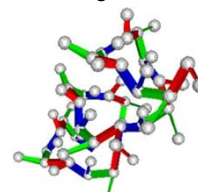
Translation + rotation in 3D



3-link arm

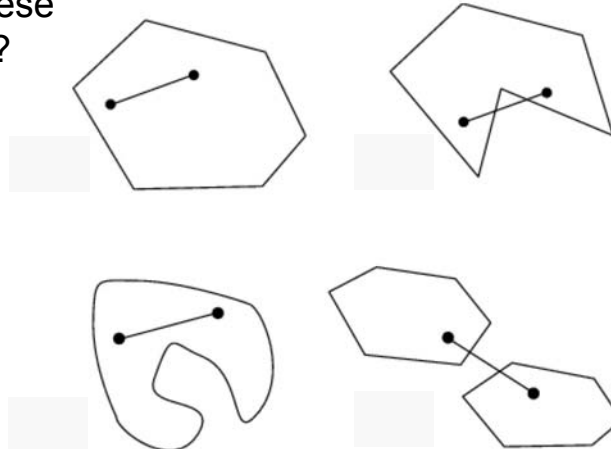


Molecule with n fixed-length bonds



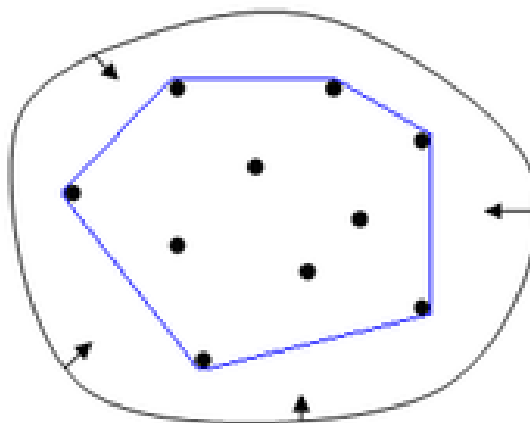
Convexity

- A set S is *convex* if and only if every line segment connecting two points in S is
- Which of these are convex?

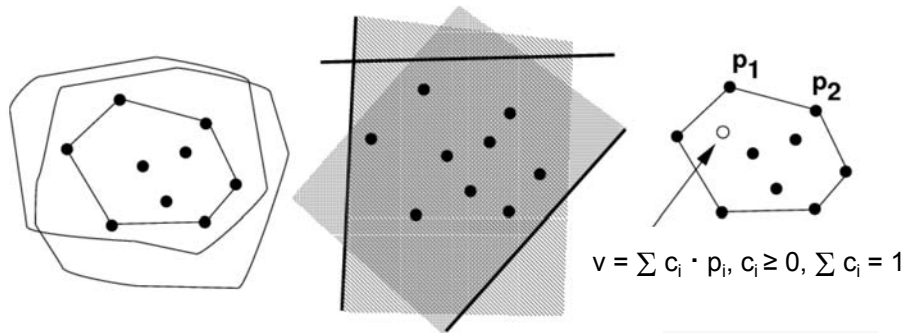


Convex Hull of a Set of Points

- Intuition: stretch a rubber band around point set



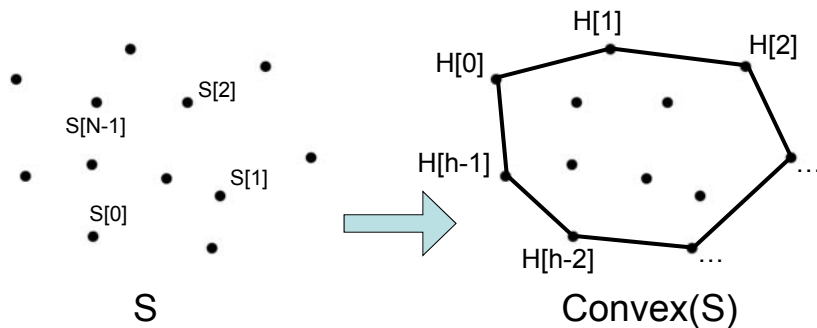
Convex Hull: Formal Definitions



- Which of these are constructive / algorithmic?

Computing 2D Convex Hull

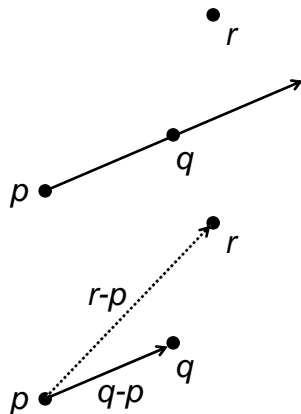
- Input: set S of N points (x_i, y_i) in 2D
- Output: polygonal boundary of convex hull of S



- How can $\text{Convex}(S)$ be computed (efficiently)?

The Leftof Predicate

- Input: three points p, q, r
- Function $\text{Leftof}(p, q, r)$ // argument order matters
- Output: 1 iff r is left of *directed line* \overrightarrow{pq} , otherwise -1



How to implement Leftof()?

1. Compute sign of determinant

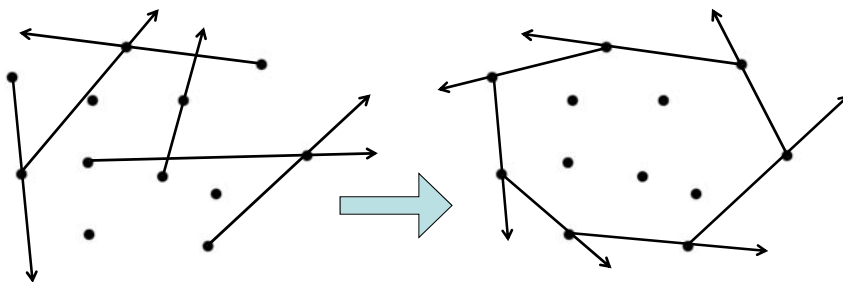
$$\begin{vmatrix} 1 & r_x & r_y \\ 1 & p_x & p_y \\ 1 & q_x & q_y \end{vmatrix}$$

2. Equivalently, find sign of z component of

Brute Force Solution

Identify point pairs that form edges of $\text{Convex}(S)$

I.e. for each pair $p, q \in S$, if $\forall r \in S - \{p, q\}$, r lies left of the directed line \overrightarrow{pq} ,



Running time for input of n points?

Can do better: $O(n^2)$, $O(n \log n)$, $O(nh)$, $O(n \log h)$!

Jarvis March Algorithm

pivot = leftmost point in S ; $i = 0$ // leftmost point must be on convex hull

repeat

$H[i] = \text{pivot}$ // store hull vertices in output point list $H[i]$, $0 \leq i < h$

endpoint = $S[0]$ // check candidate hull edge [pivot .. endpoint]

for j from 1 to $|S|-1$

if (Leftof (pivot, endpoint, $S[j]$))

endpoint = $S[j]$

pivot = endpoint; $i++$

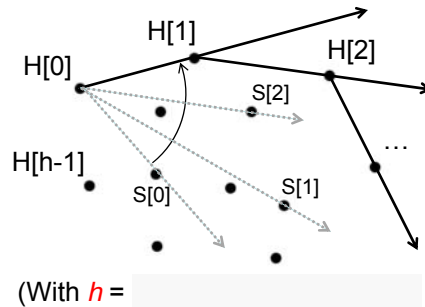
until endpoint == $H[0]$

Outer loop runs times;

inner loop does work

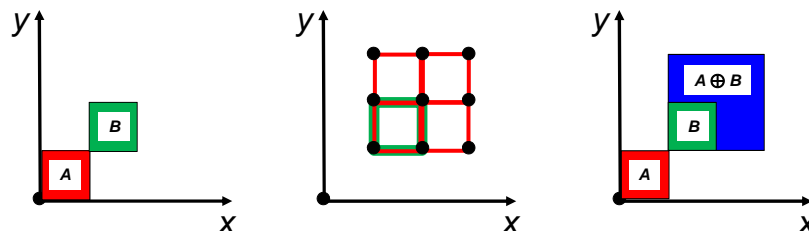
Running time for input

set of n points? “Output-sensitive” algorithm.



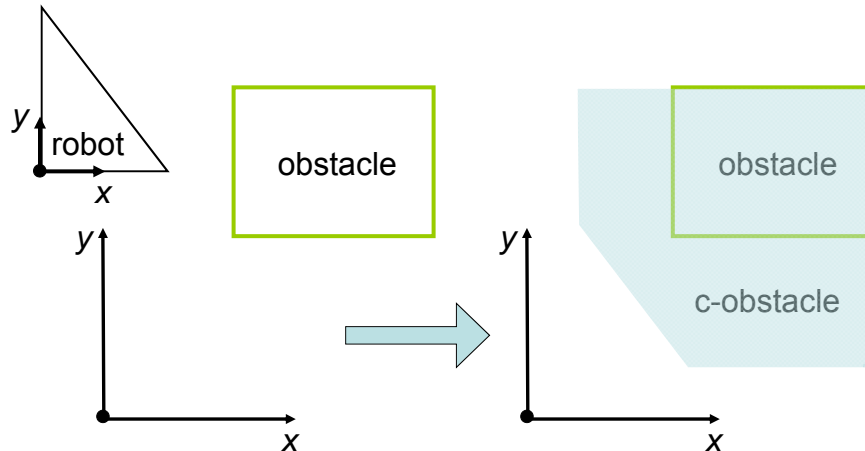
Minkowski Addition

- Given two sets $A, B \in \mathbb{R}^d$, their *Minkowski sum*, denoted $A \oplus B$, is the set $\{a + b \mid a \in A, b \in B\}$
 - Result of *adding* each element of A to each element of B
- If A & B convex, just add vertices & find convex hull:



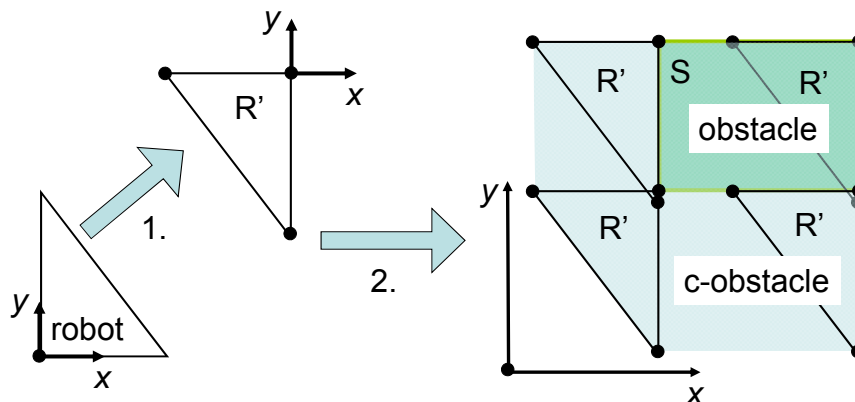
Computation of C-obstacles

- Inputs: robot polygon R and obstacle polygon S
- Output: c-space obstacle $c\text{-obstacle}(S, R)$



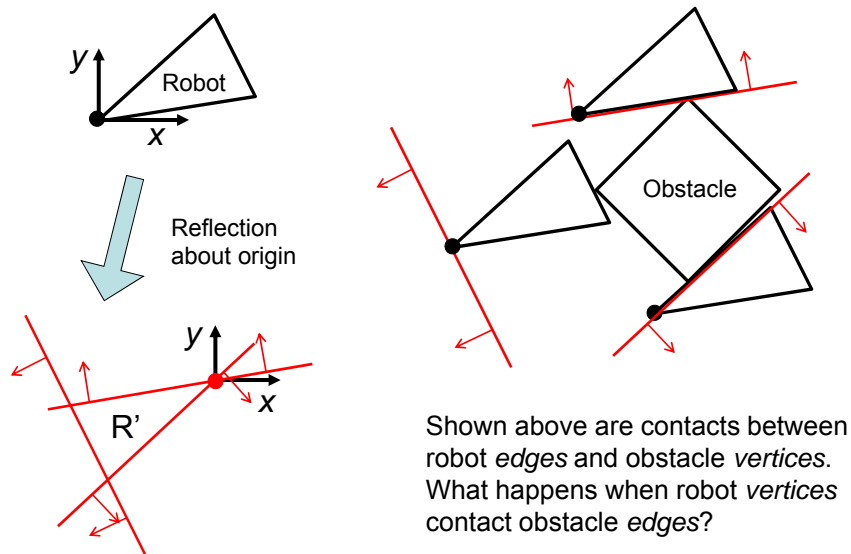
C-obstacle Computation

1. Reflect robot R about its origin to produce R'
2. Compute Minkowski sum of R' and obstacle S

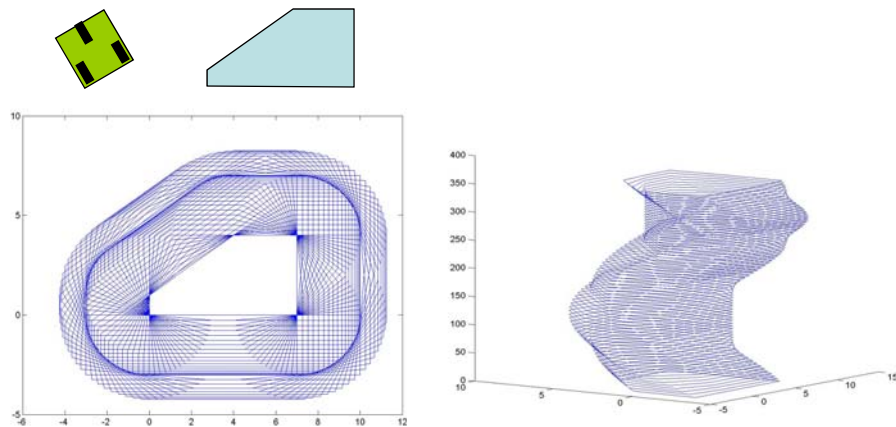


Sanity check: shaded region is infeasible for

Why *Reflect* the Robot?



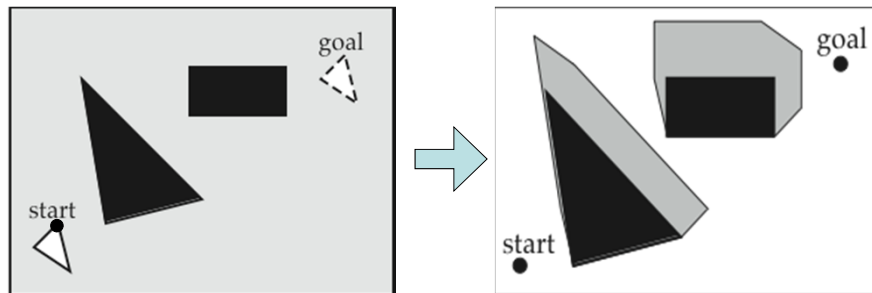
C-obstacles with Rotations



How do we compute this object?

Back to Motion Planning

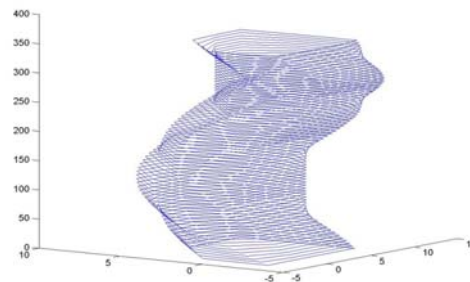
- Given robot and set of obstacles:
 - Compute C-space representation of obstacles
 - Find path from robot start pose to goal pose (point)



- Unfortunately, we have a rather serious problem:
 - We have constructed a representation of the *obstacles*
 - But we need to search a representation of the *freespace*!

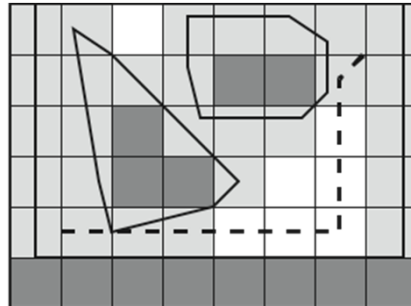
Computational Complexity

- The best deterministic motion planning algorithm known requires *exponential* time in the C-space dimension [Canny 1986]
- D goes up fast – already D=6 for a rigid body in 3-space; articulation adds many more DOFs
- Simple obstacles have complex C-obstacles
- Impractical to compute explicit representation of freespace for robot with many DOFs
- What to do?

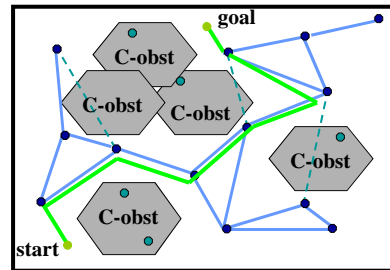


Strategies

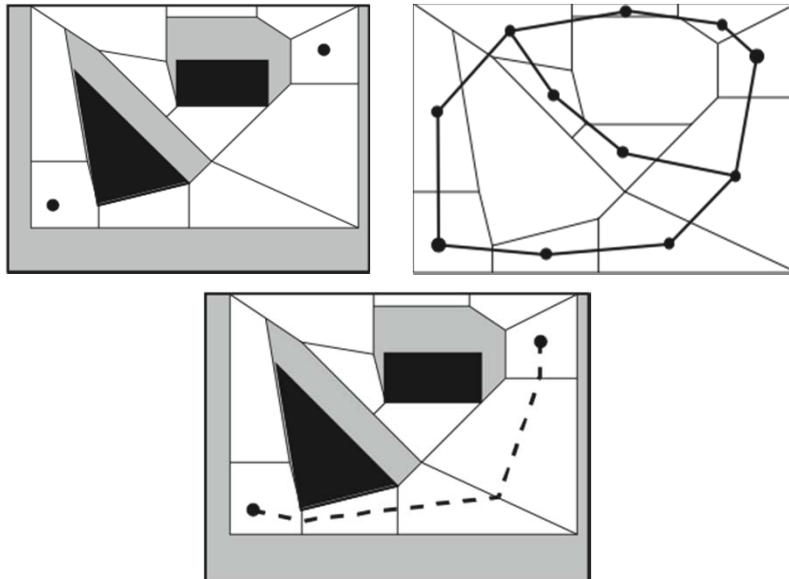
- Approximate: use regular subdivision of freespace
- Randomize: sample and evaluate C-space poses
- Sacrifice for gains in



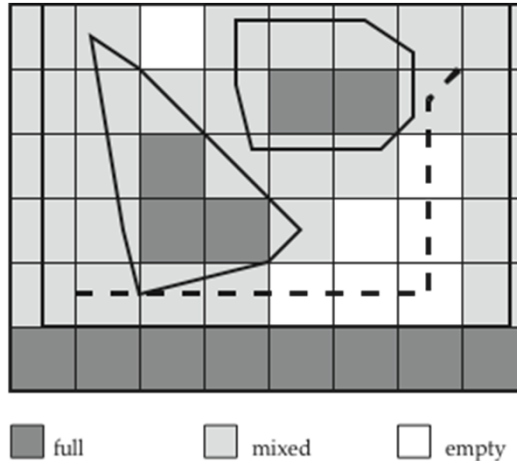
full
 mixed
 empty



Example: Exact Decomposition



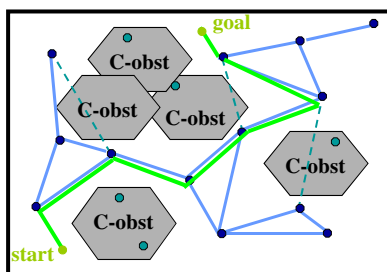
Approximate Cell Decomposition



- Advantage: recasts complex original problem as search within space of many, simpler motion plans

Probabilistic Road Maps for Motion Planning [Kavraki et al. 1996]

C-space



Roadmap Construction (Pre-processing)

1. Randomly generate robot configurations (nodes)
 - Discard invalid nodes (how?)
2. Connect pairs of nodes to form **roadmap edges**
 - Use simple, deterministic *local planner*
 - Discard invalid edges (how?)

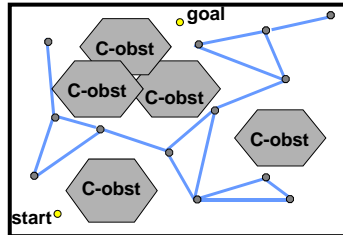
Plan Generation (Query processing)

1. Add *start* and *goal* poses into the roadmap
2. Find path from *start* to *goal* within roadmap
3. Generate a motion plan for each edge used

Requires two primitive operations:

1. Method for sampling C-Space points
2. Method for "validating" C-space points and edges

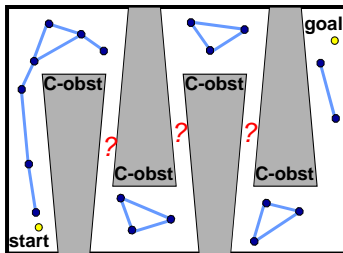
PRMs: Pros and Cons



Advantages

1. *Probabilistically complete*
2. Easily applied to high-dimensional C-spaces
3. Supports fast queries (w/ enough preprocessing)

Many success stories in which PRMs have been applied to problems previously thought intractable



Disadvantages

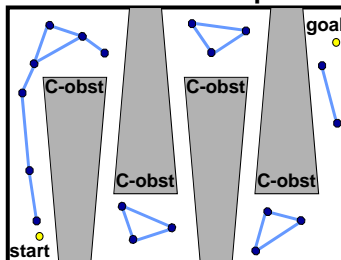
PRMs don't work well for some problems:

-
-

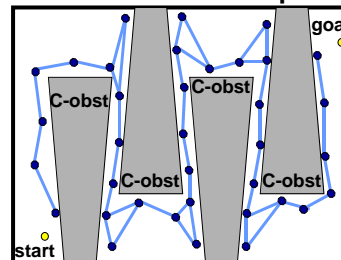
Sampling Around Obstacles: OBPRM [Amato et al. 1998]

To navigate narrow passages we must *sample* inside them
Most PRM nodes placed where planning is easy, not where it's hard

PRM Roadmap



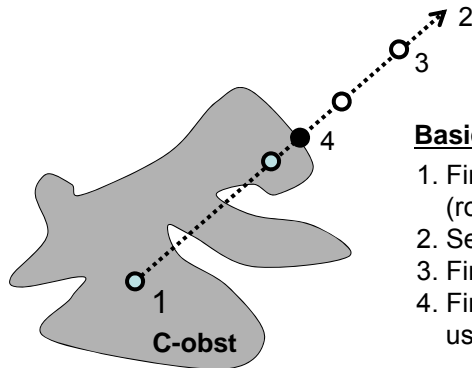
OBPRM Roadmap



Idea: Can we sample nodes near C-obstacle surfaces?

- We cannot explicitly construct the C-obstacles, but...
- We do have models of the (workspace) obstacles!

Finding Points on C-obstacles



Basic Idea (for workspace obstacle S)

1. Find a point in S's C-obstacle (robot placement colliding with S)
2. Select random direction in C-space
3. Find freespace point in that direction
4. Find boundary point between points using binary search (collision checks)

Note: we can use more sophisticated approaches to try to "cover" C-obstacle

Summary

- Introduced drastically simplifying transformation
 - Based on two useful geometric constructions
- Enables use of familiar techniques...
 - Discretization
 - Random sampling
 - Bisection
 - Graph search
- ... To solve high-dimensional motion planning
- We'll use these ideas in Lab 6 (path planning)