

Black-white coloring Let $G$ be a connected, weighted simple graph. Color each vertex of $G$ black or

胃: Connected Graph G Let $G$ be a connected, weighted simple graph. Assume all edges have different weights.

| 6 | 9 | 13 | 7 |
| :---: | :---: | :---: | :---: |
| 12 |  | 10 | 5 |
|  |  |  | 4 |


| 3 |  | 10 |  |
| :---: | :---: | :---: | :---: |
| 3 | 1 | 4 | 14 |
| 15 | 8 | 11 | 2 |

Black-white coloring
Let $G$ be a connected, weighted simple graph. Color each vertex of $G$ black or white, not all same color.

## Gray Edges <br> A gray edge connects vertices with different colors: <br> There must be a gray edge since $G$ connected

##  <br> Theorem: <br> $\mathbb{e}$ is an edge of every min-weight spanning tree (MST)

## Gray Edges <br> A gray edge connects vertices with different colors: <br> Let $e$ be a min-weight gray edge.

Min Gray Edges Sufficient
There is a spanning tree built of min-weight gray edges
-- from previous slides.


```
|\mp@code{cosccc|}
min-weight gray = MST
Corollary: There is a unique
    MST. It consists of all
    min-weight gray edges
    under black-white
    colorings.
```


## Gray Edge Swap Lemma Suppose e not an edge of $\operatorname{css} C$. Then there is an edge $g$ of $C$ : <br> (i) $w t(e)<w t(g)$ <br> (ii) $C-g+e$ is a css

##  <br> So $\mathbb{e}$ is necessarily in any min-weight css.

Gray Edge Swap Lemma
So $C$ is not minimum css because $C-g+e$ has smaller weight.

蹋暗 Proof of Swap Lemma
Say $e=\langle\mathbf{u}-\mathrm{v}\rangle$.
C connected, so have path
$\vec{p}=u-\cdots-v$.
c) (1) (2) Albert R Meyer October 31, 2017





