Introduction to Algorithms: 6.006

Massachusetts Institute of Technology March 1st, 2011
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Problem Set 3

This problem set is divided into two parts: Part A problems are theory questions, and Part B problems are programming tasks.

Both Part A and Part B questions are due **Monday, March 7th** at **11:59PM**.
Solutions should be turned in through the course website. Your solution to Part A should be in PDF format using LaTeX. Your solution to Part B should be a valid Python file. A template for writing up solutions in LaTeX is available on the course website. Remember, your goal is to communicate. Full credit will be given only to a correct solution which is described clearly. Convoluted and obtuse descriptions might receive low marks, even when they are correct. Also, aim for concise solutions, as it will save you time spent on write-ups, and also help you conceptualize the key idea of the problem.

Part A

Problem 3-1. [25 points] **Recurrences**

Solve the following recurrences:

  (a) [5 points] $T(n) = 2T(\frac{n}{2}) + n + \log n$

  (b) [5 points] $T(n) = 3T(\frac{n}{4}) + \sqrt{n}$

  (c) [5 points] $T(n) = T(\frac{n}{2}) + T(\frac{n}{4}) + n^2$ (*Hint:* Calculate upper and lower bounds on $T(n)$ by converting the problem into recurrences you can solve using the Master Theorem.)

Merge sort typically divides an array in half, forming two recursive subproblems, recursively sorts the two subproblems, and then merges the two solutions. Consider the ***c-merge-sort*** algorithm, which divides the array into c (roughly) equal-length parts, recursively sorts these c subproblems, and then merges the c solutions. To merge c sorted lists, the algorithm uses a min-heap initialized with the first element of each list, and repeatedly removes the minimum element from the heap, appends that minimum element to the merged list, and then inserts the minimum element's successor (from its source list) into the heap, until all the elements are appended to the merged list.

  (d) [5 points] Write the recurrence relation for c-merge-sort on a list of n integers. Solve it to determine the running time for some constant c.

  (e) [5 points] What happens to the running time if c is not a constant, but rather some function of n?

Problem 3-2. [25 points] ***d*-ary Heaps**

In class, we've seen binary heaps, where each node has at most two children. A d-ary heap is a heap in which each non-leaf node (except perhaps one) has exactly d children. For example, this is a 3-ary (max) heap:
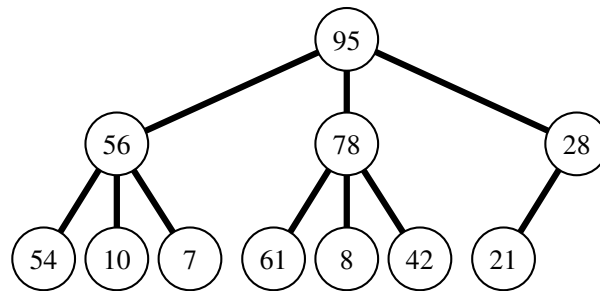
(a) [2 points] Suppose that we implement a d-ary heap using an array A, similarly to how we implement binary heaps. That is, the root is in $A[0]$, its children are in $A[1 \ldots d]$, and so on. How do we implement the PARENT(i) function, which computes the index of the parent of the ith node, for a d-ary heap?

(b) [2 points] Now that there might be more than two children, LEFT and RIGHT are no longer sufficient. How do we implement the CHILD(i, k) function, which computes the index of the kth child of the ith node? ($0 \le k < d$)

(c) [5 points] Express, in asymptotic notation, the height of a d-ary heap containing n elements in terms of n and d.

(d) [5 points] Give the asymptotic running times (in terms of n and d) of the HEAPIFY and INCREASE-KEY operations for a d-ary heap containing n elements.

(e) [6 points] Let's suppose that when we build our d-ary heap, we choose d based on the size of the input array, n. What is the height of the resulting heap (in terms of n) if we choose $d = \Theta(1)$? What if $d = \Theta(\log n)$? What about $d = \Theta(n)$?

(*Hint:* remember that $\log_d n = \frac{\log n}{\log d}$. This might simplify your expressions a little.)

(f) [5 points] What are the running times of HEAPIFY and INCREASE-KEY for the three choices of d above? Do the running times increase or decrease when you increase d? If your program calls HEAPIFY and INCREASE-KEY the same number of times, what would be your choice for d and why?

Part B

Problem 3-3. [50 points] **Pset Scheduling**

Ben Bitdiddle is behind on his problem sets. In fact, he is already late on N different problem sets ($1 \le N \le 100,000$). Fortunately for Ben, all of his classes accept late homework with a grade penalty for each day late (unlike 6.006).

Suppose that problem set i, where $1 \le i \le N$, has a penalty of P_i points per day late, and takes one full day to complete. There is no limit to the number of penalty points Ben can accrue. (Ben's penalty adjusted score can become negative.) Ben is required to finish each problem set.

You should implement a function `best_score(penalties)` which takes as input the list of daily point penalties P_i and returns the minimum number of penalties points it is possible for Ben to be assigned.

As part of your program, you will need to do some sorting. You should write your own implementation of heap sort for this problem. Your implementation of heap sort should have the signature `heap_sort(list)` where `list` is the list to sort. Your function should sort `list` in descending order using the default Python ordering defined by < and >. This means, for instance, that `heap_sort([5,1,4,0])` should return `[5,4,1,0]`. Using this function specification will allow us to better test your code if you have a bug and give you more partial credit.

Sample Input:
```
[1,5,2,3]
```

Sample Output:
```
21
```

Sample Explanation:
In the sample above, Ben works on pset #2 on the first day, pset #4 the next day, pset #3 the third day, and pset #1 the last day. In total, he accrues $5 \cdot 1 + 3 \cdot 2 + 2 \cdot 3 + 1 \cdot 4 = 21$ penalty points.