

Recitation 8

Overhead:

- Psets:
 - look @ solns
 - stats **A** mean: median:
 - B** mean: median:
 - handout @ end
- QUIZ:
 - 10/15 7:30-9:30 PM
 - conflict 10/16 8-10AM
 - email staff w/ reason
 - if previously emailed - confirm.
- Feedback Form

Agenda

- overhead
- exercises
- coding
- (?) shell sort
- Psets

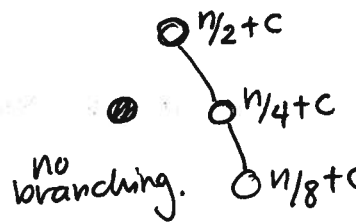
= Questions from lecture / anything you want reviewed.

= Pset

- $n \log n$ bound on $T(n) = T(n/2) + n/2 + c$ is not tight.
 - yes you can prove $O(n \log n)$
 - missing proof of $\Omega(n \log n)$ ← impossible.

real runtime $\Theta(n)$

- why? b/c $n/2$ does not add n per step!



- OOPS - my apologies.

$$\log(f(n)) = \Theta(\lg(g(n))) \Rightarrow f(n) = \Theta(g(n)) \text{ is FALSE}$$

$$\log(2^n) \text{ vs. } \log(3^n)$$

$$n \log(2) \quad n \log(3)$$

yes Θ ...

BUT you need to convert back.

you can't cancel const values b/c

~~they are actually~~ they are actually e^c .

Shell Sort

idea: some sorts are very fast for nearly sorted data.
we can nearly sort data really fast then pass off the data
sort every n^{th} element, n decreases each time.
- lets elements make bigger "jumps" toward correct position


visually: use columns

| | | | |
|---|---|-----|---|
| 1 | 2 | 3 | 4 |
| 5 | 6 | ... | |

 then sort the columns.

gap sequence matters!

1 4 10 23 57 132 301 ...

algorithm on wikipedia 

worst case $\Theta(n^2)$ w/ gap seq 1 2 4 8 ...
best $\Theta(n \log^2 n)$

works well on real data.

Selection Sorts

idea: select an element, remove it, repeat.
if you select in a logical manner (max/min)
you can sort!

variables/options: storage data structure
selection criteria.

| struct | array | array | balanced BST | heap |
|--------------------|----------|----------|-----------------|--------------|
| selection | min | max | min/max | min/max |
| time per select | $O(n)$ | $O(n)$ | $O(\lg n)$ | $O(\lg n)$ |
| total | $O(n^2)$ | $O(n^2)$ | $O(n \lg n)$ | $O(n \lg n)$ |

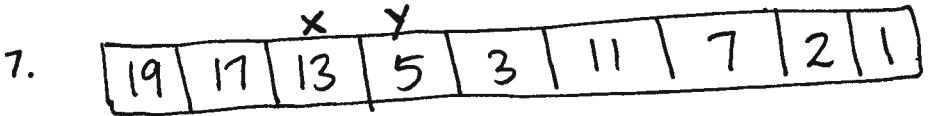
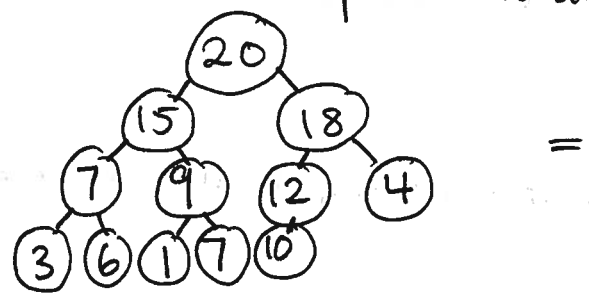
EXERCISES

1. state a situation for which insertion sort would be a good choice of sort.
2. what is the worst possible sort you can think of? each operation must be productive (i.e. no while (1) {3})
3. analyze the runtimes of insertion + merge sort on lists that are:
 - a. sorted
 - b. reverse-sorted
4. how does insertion sort's runtime change if you use a linked list?
5. what is the runtime of the following code? $\Theta(\quad)$
then change the code to improve the runtime.
you may assume all elements are ≥ 0 .

```
def mergesort(A):  
    if len(A) == 1 : return A  
    m = len(A) / 2  
    B = len(A) * [-1]  
    C = len(A) * [-1]  
    B[:m] = mergesort(A[:m])  
    C[m:] = mergesort(A[m:])  
    return merge(A B C)
```

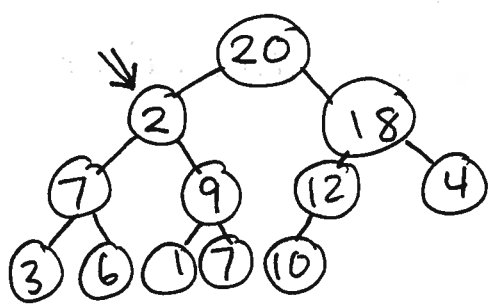
```
def merge(A B C):  
    a = 0  
    b = 0  
    c = 0  
    while b < len(B) and c < len(C):  
        if B[b] < C[c]:  
            A[a] = B[b]  
            b += 1  
        else: A[a] = C[c]  
            c += 1  
    a += 1
```

6. convert the heap into its array representation



- parent of x :
- left (x) :
- right (x) :
- depth of heap :
- parent(y) :
- left (y) :
- right (y) :

8. restore the max-heap property



9. get code heaps.py from website (or take a handout)
- a. implement left(i), right(i) and parent(i)
 - b. complete the missing code in heapify
 - optional: what code would go here for a min-heap?
 - c. modify heapify to take a length argument.
 - d. implement delete-max
 - e. implement heap sort.
 - f. challenge: implement build which builds a heap in place from an arbitrary array.

if you change the method signatures - make sure to alter the tests accordingly!