

Agenda

overhead

- set recitation on website

- Overhead
- Warmup
- BST
 - defn
 - searching
 - insertion
 - next-largest
- objects in python
- representations of BSTs

Warmup Answers

5 9 4 7 8 2 10 1 6 3

↓

min	4
max	3
3	4
8	2

1 2 3 4 5 6 7 8 9 10

min	1
max	10
3	3
8	8

10 1 9 2 8 3 7 4 6 5

min	1
max	1
3	6
8	8

5 4 3 2 1 6 7 8 9 10

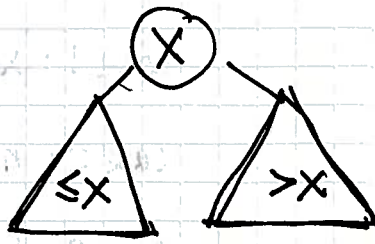
min	5
max	6
3	3
8	9

Concept of Binary Search Tree

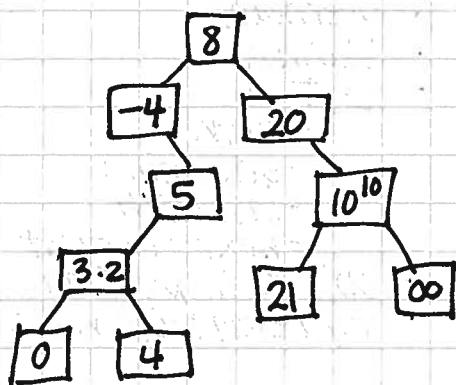
- binary: splits in 2 each time
- search: optimizes search time

specs of data structure

- recursive structure
- node holds a value
- left child tree nodes must be \leq the value
- right child tree nodes must be $>$ the value.



example BST



find 0, -3, 7, ∞

• compare x to root

$x \leq \text{root}$ done
 $x < \text{root}$ goto left child
 $x > \text{root}$ goto right child
 repeat
 (unless child doesn't exist -
 in which case $x \notin \text{tree}$)

why bother?

- can handle so size spaces - not just ints!
- nice insertion properties
- sometimes nice searching properties
- some natural operations easy: min/max

SUX!

- sorted input \rightarrow AWFUL tree (degenerate case)
- deletion is a bitch
- some natural operations difficult: next biggest
- no more random indexing

SEARCHING

- best case time : balanced tree : $\lg(n)$
- worst case time : line : n

INSERTION

- same running times as \uparrow
- same

NEXT LARGER

given x , find $y \in \text{BST}$ s.t. $x < y$ and $\forall z \in \text{BST}$ if $z > x$ $z > y > x$

case 1: x has a right child subtree $y \rightarrow \min(y)$

case 2: no right child $\& \ p(x) > x \Rightarrow p(x)$

case 3: no right child $\& \ p_i(x) < p_{i-1}(x) < \dots < p_1(x) < x \Rightarrow p_{i+1}(x)$
and $p_{i+1}(x) > p_i(x)$

Recitation 3 continued

classes in python

class Name:

class-var = <value>

def class_fcn(self):
do stuff

def fcn2(i, j):
do more

def __init__(self):
auto defined does nothing
can put more code

def __init__(self, x, y):
m

x = Name()

x.class-var

x.class_fcn()

x.fcn2(3, 10)

x = Name(x, y)

example: class counter:

~~count = 0~~

def inc(self):
self.count += 1
def inc(self, a):
self.count += a

def __init__(self):
self.count = 0
def __init__(self, x):
self.count = 10

x = counter()

x.inc()
print x.count
x.inc(10)
print x.count
x.name = "foo"
print x.name
del x.name

need not declare.

BST representation

- in class we used code that doesn't make a new class.

node = [, val,] array of 3 elts
list list

- now use objects to make code more understandable.

[see handout] I implemented next larger + delete for you

exercise: implement printing all elements in order
(or ret. an array of all in order)

⇒

i think I will omit this in the interest of time.
just go over it when looking @ bst code.

array

low child (i) $\rightarrow 2i$
high child (i) $\rightarrow 2i+1$

delete flag

6.006 Recitation 3 Warmup

For each problem:

1. construct the Binary Search Tree formed by inserting the given elements in order
2. determine how many steps the following operations take

min()

max()

find(3)

find(9)

[5 9 4 7 8 2 10 1 6 3]

[5 4 3 2 1 6 7 8 9 10]

[1 2 3 4 5 6 7 8 9 10]

[10 1 9 2 8 3 7 4 6 5]

