

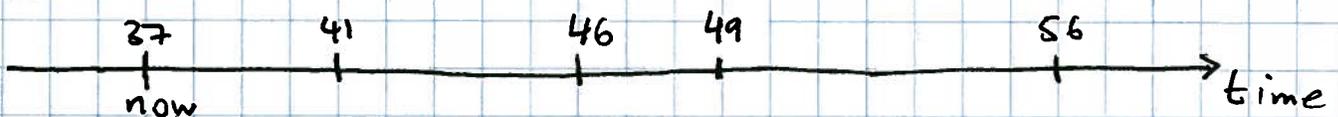
6.006 Lecture 3: Binary Search Trees

- Runway reservation problem
- Binary Search Trees (BSTs), operations
- CLRS Ch 10 + 12.1-3

Runway Reservation System

- Maintains a set R of future landing times
- When a plane lands, remove from the set
- For a request to land at time t :
 - reject request if there are reservations within < 3 minutes of t ,
 - otherwise add t to set

Example



$$R = \{41, 46, 49, 56\}, \quad n = |R| = 4$$

request for time:

44 reject (too close to $46 \in R$)

53 add reservation (just right)

20 already past, not allowed.

Goal: handle requests & landings in $O(\lg n)$ time.

- Keeping R as an (arbitrarily ordered) list
 - Checking a request takes $O(n)$ time
 - Deleting upon landing takes $O(n)$ time

- Keeping R as a sorted list

init: $R = []$

req(t): if $t < \text{now}$ return 'error'

for i in range($\text{len}(R)$):

if $\text{abs}(t - R[i]) < 3$: return 'error' } $O(n)$

$R.$ append(t)

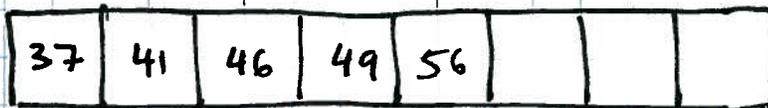
$R.$ sort() $\leftarrow O(n \lg n)$

land: $t = R[0]$

if ($t \neq \text{now}$) return 'error'

$R = R[1:]$ $\leftarrow O(n)$

- Can we do better with a sorted list?



can use binary search to determine conflicts in $O(\lg n)$ but inserting t into R and deleting $R[0]$ still takes $O(n)$
 \Rightarrow better (no certina) but not enough.

- Indicator representations

- Represent a set R such that

$$A[t] == \text{True} \quad \text{iff} \quad t \in R$$

- Python dictionaries
- Python sets
- A Python list of all possible landing times (all possible elements of R)

- Fast insertion, deletion

- Checking a request is fast if landing times are whole minutes, but expensive for hires

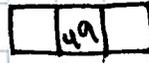
- Not a flexible data structure:

How many planes land on or before t ?

~~Need to iterate over all landing times~~
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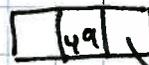
- Binary Search Trees (BSTs)

1. insert(49)



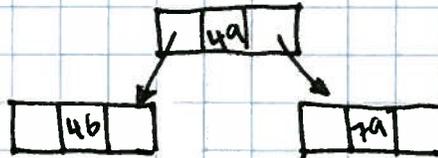
this is the root

2. insert(79)



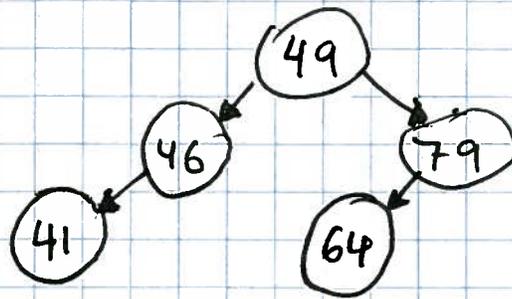
to find elements > 49
go right

3. insert(46)



elements < 49 are
in the left subtree

A more compact ^{graphical} representation:



Same representation
in the program:
every node has
a value and
left/right pointers

find(64): follow left/right pointers until done

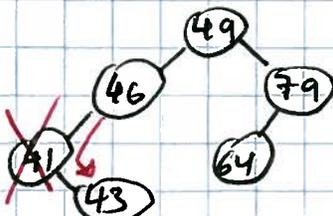
find(48): how do you tell that 48 is not in the set?

findmax(): just go right until you can't go right

findmin(): similar, go left

deletemin(): min & max nodes will have zero

or one child \Rightarrow find, eliminate, & patch

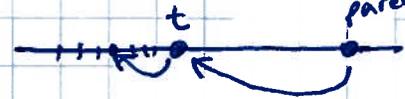


next-larger(t): $v = \text{find}(t)$



if $\text{right}(v) \neq \text{None}$: return $\text{findmin}(\text{right}(v))$

else: $y = \text{parent}(v)$

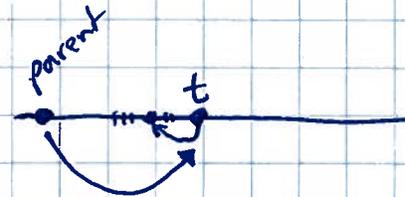


while $y \neq \text{None}$ and $v = \text{right}(y)$:

$v = y$

$y = \text{parent}(v)$

return y



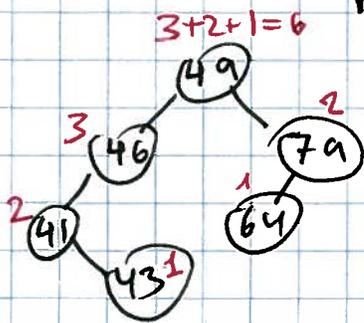
or

- All of these operations on the BST take $O(h)$ time, where h is the height of the tree

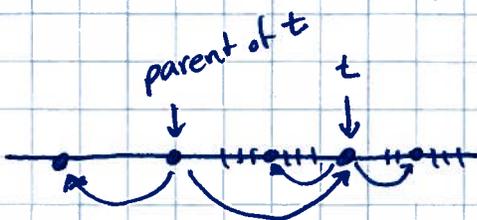
- Other operations that we can do in $O(h)$ time:

rank(t) how many nodes $\leq t$?

must augment each node with size of subtree, update during insertions & deletions.



delete(t) more tricky if t has 2 children, but still $O(h)$.



(delete next-larger(t) than put it in the node that stored t ; see textbook).

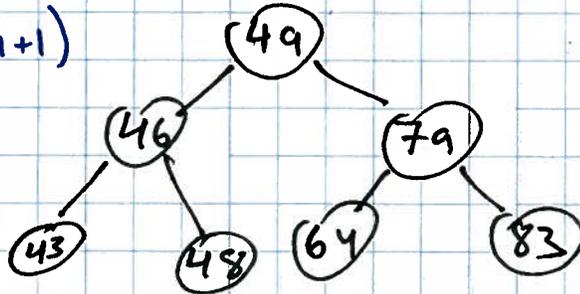
- So how high is a binary tree with n vertices?

- If we're lucky, $h = O(\lg n)$

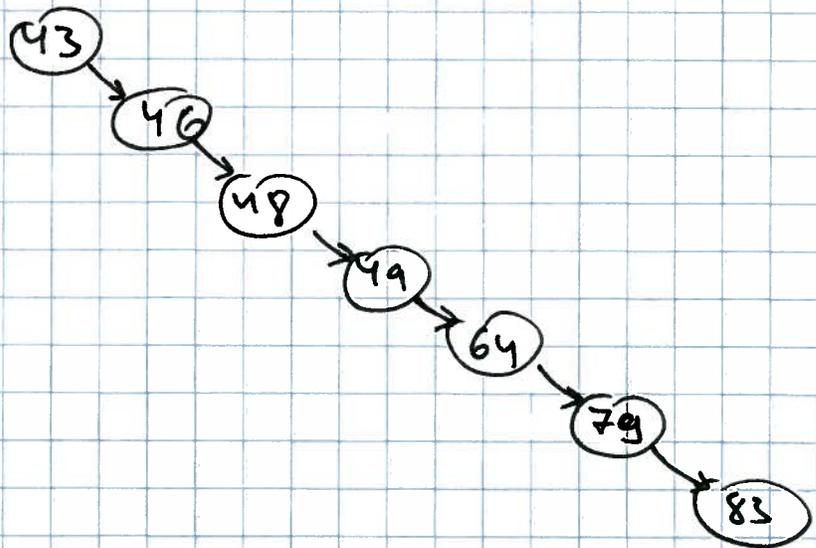
$$h = 3 = \lg(n+1)$$

or

$$n = 2^h - 1$$



- But h can be large, $h = n$



```
# Sivan Toledo, 6.006
```

```
#--- Sorted Lists ---
```

```
def sortedlist_init():  
    return []
```

```
def sortedlist_print(R):  
    print R
```

```
def sortedlist_insert(R,t):  
    R.append(t)  
    R.sort()  
    return R
```

```
def sortedlist_deletemin(R):  
    return R[1:]
```

```
def sortedlist_runway_insert(R,t):  
    for i in range(len(R)):  
        if abs(t-R[i]) < 3:  
            print 'error: ' + str(t) + ' is too close to ' + str(R[i])  
            return R  
    R.append(t)  
    R.sort()  
    return R
```

```
#--- Binary Search Trees ---
```

```
def bst_init():  
    return None
```

```
def bst_insert(T,t):  
    if T == None:  
        return [None, t, None]  
    if t <= T[1]:  
        T[0] = bst_insert(T[0],t)  
        return T  
    else:  
        T[2] = bst_insert(T[2],t)  
        return T
```

```
def bst_str(T):  
    if T == None:  
        return ''  
    return bst_str(T[0]) + ' ' + str(T[1]) + ' ' + bst_str(T[2])
```

```
def bst_print(T):  
    print '[' + bst_str(T) + ' ]'
```