

Parallelizing SAT Solver

6.338 Applied Parallel Computing
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- The SAT problem:
- Given a formula made of boolean variables and operators, find an assignment to the variables that makes it true:
 - i.e. $(P \vee Q) \wedge (\sim P \vee R)$
 - A solution: $\{P = \text{false}, Q = \text{true}, R = \text{false}\}$
- SAT is NP-Complete (Hard)

What is SAT?

- Conjunctive Normal Form is a set of clauses, each containing a set of literals
- Clauses are "Anded" with one another
- Literals are "Ored" with one another in a clause
- An efficient SAT solver takes a formula in CNF and returns an assignment as solution or says none exists

SAT Solver

- Naïve Solve:
 - Enumerate assignments and check formula for each assignment
 - For k variables, 2^k assignments!
- DPLL/Davis-Putnam-Logemann-Loveland
 - Back-tracking
 - Unit propagation
 - Pure literal elimination

More on SAT Solver

- function DPLL(Φ , Env):
 - if Φ is empty
 - return Env;
 - if Φ contains an empty clause
 - return null;
 - for every unit clause l in Φ
 - $\Phi = \text{unit-propagate}(\Phi, \text{Env})$;
 - for every literal l that occurs pure in Φ
 - $\Phi = \text{pure-literal-assign}(l, \Phi, \text{Env})$;
 - $l := \text{choose-literal}(\Phi)$;
 - $\text{Env2} = \text{DPLL}(\Phi, \text{assign}(l, \text{Env}))$;
 - If Env2 is null
 - return $\text{DPLL}(\Phi, \text{assign}(\text{not}(l), \text{Env}))$;
 - return Env2
- Env is a map of assignments

DPLL Algorithm

5	3			7				
6			1	9	5			
	9	8					6	
8				6				3
4			8		3			1
7				2				6
	6					2	8	
			4	1	9			5
				8			7	9

A sudoku puzzle...



5	3	4	6	7	8	9	1	2
6	7	2	1	9	5	3	4	8
1	9	8	3	4	2	5	6	7
8	5	9	7	6	1	4	2	3
4	2	6	8	5	3	7	9	1
7	1	3	9	2	4	8	5	6
9	6	1	5	3	7	2	8	4
2	8	7	4	1	9	6	3	5
3	4	5	2	8	6	1	7	9

...and its solution numbers marked in red



Sudoku Puzzle

- Rules of the game
 - No two same digit can appear in a single column
 - No two same digit can appear in a single row
 - No two same digit can appear in a single sub-grid
 - Exactly one number must occupy each cell

Sudoku Puzzle

- For each column, row, sub-grid, cell:
 - An At-least clause
 - An At-most clause
 - Together enforces exactly once behavior
- Example:
 - An At-Least clause: [119 129 139 149 159 169 179 189 199].
 - An associated series of At-Most clause:
 - [-119 -129], [-119 -139], [-119 -149]...[-119 -199]
 - [-129 -139], [-129 -139]...[-129 -199]
 - ...
 - [-189 -199]
 - These 37 clauses together ensure that "9" appears exactly once in row 1.

Reduction

- Unit Propagation:
 - $(X)(X \vee Y \vee Z)(\sim X \vee Y)(Y \vee Z \vee \sim W)(\sim Z \vee W) \rightarrow (Y)(Y \vee Z \vee \sim W)(\sim Z \vee W)$
 - $(Y)(Y \vee Z \vee \sim W)(\sim Z \vee W) \rightarrow (\sim Z \vee W)$
 - Recursively searches the problem for a unit clause then simplify the problem accordingly in serial
- Pure Literal Elimination
 - Tracks whether a literal has become pure at each step of recursion
 - If so, simplify the problem accordingly

Operations

- Spread the problem across nodes
- Have each operations performed on sub-problem in parallel
- Obtain new sub-problem, and recurse

Parallelization

Performance Comparison:

Puzzle Number	Number of Recursion Steps	Performance in milliseconds		
		Java	Matlab Serial	Matlab Star-P
1	145	76	89	4105
2	219	88	101	3809
3	732	148	168	5672
4	2644	465	481	12691
5	4052	858	912	19780
6	8234	1485	1674	24536

Performance

- Sudoku is too small, too easy
 - Computation requires is light
 - Does not justify for communication cost
- Bigger problem more likely to see benefits of parallelization

Conclusion