## Problem Set 10

This problem set is due Wednesday, November 23 at noon.

In this problem set you will develop an algorithm for canceling flow cycles in a given flow assignment. In general graphs this can be done in  $O(m \log n)$  time using Sleator's and Tarjan's dynamic trees. You will use the relation between dual shortest paths and circulations to give a linear time algorithm in planar graphs.

1. Let G be a planar graph with non-negative capacities on its arcs. Let  $\phi$  be shortest path distances from  $f_{\infty}$  in  $G^*$ . Let  $\theta$  be the circulation induced by  $\phi$ . That is,

$$\theta(d) = \phi(\text{head}(d)) - \phi(\text{tail}(d)).$$

Recall that a residual path is a path whose darts all have strictly positive capacities. Show that there are no counterclockwise residual cycles in the residual graph  $G_{\theta}$ .

- 2. What price function  $\phi'$  would you use to get the same property as in (1), but with no clockwise residual cycles?
- 3. Use parts (1) and (2) to give a linear time algorithm that, given a flow assignment  $\gamma$  in G makes  $\gamma$  acyclic by removing all flow cycles in  $\gamma$ . That is, it produces another flow assignment  $\gamma'$  s.t.
  - (a)  $\gamma' \gamma$  is a circulation
  - (b) for every arc  $a, \gamma'((a,1)) \leq \gamma((a,1))$
  - (c) for any cycle C there is a dart  $d \in C$  s.t.  $\gamma'(d) \leq 0$