User’s Guide to Basic Joshua
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1. Introduction

1.1. About the Joshua Language

Joshua is an extensible software product for building and delivering expert system applications. It is implemented on the Symbolics 3600 family, on top of the Symbolics Genera environment. Joshua is optimized for applications where performance and delivered functionality are important.

Joshua is a very compact system, organized around a small number of core functions. Joshua's default structures provide a simple declarative core language with built-in facilities for application development. Programming with the defaults is very straightforward, allowing you to build effective applications quickly. This is due to several features:

- The syntax of Joshua is uniform, statement-oriented, and Lisp-like, so that you need not learn an entirely new language.

- The interface to the database (any database) is simple and uniform, consisting of the three functions, joshua:ask, joshua:tell, and joshua:clear.

- Special Zmacs facilities like bracket matching and special characters ease program development.

User's Guide to Basic Joshua, the first manual in the Joshua documentation set, covers everything you need to know to program using Joshua's built-in facilities.

Among Joshua's strengths is that this system is a coherent, multi-level environment, making advanced features available when you need them. Joshua is built around some 30 core functions, the Protocol of Inference, which are accessible to the user for modification.

This modularity and accessibility offer powerful advanced features: user interfaces, control structures, storage structures can all be customized to reflect what is most natural for the application; external databases can be accessed; existing software tools can be seamlessly integrated into the Joshua application; performance can be fine-tuned. We present all these topics in the companion documentation volume: See the document Joshua Reference Manual.

Joshua dovetails with Genera and Symbolics Common Lisp in much the same way that Lisp and Flavors do. Joshua itself is implemented with Flavors. Joshua is closely integrated with Lisp, and Lisp code can be used within Joshua rules. All of Genera's program development facilities are available to Joshua, namely, the Zmacs editor, Dynamic Windows, formatted output, debugging support, and the User Interface Management System.
1.2. About the Joshua Documentation

Joshua is a powerful and sophisticated tool that can be used at many levels, by people with varying AI programming experience, ranging from the relatively inexperienced to the expert. While the Joshua documentation is designed to help users at any level get the most out of Joshua, it is not an introductory AI text; we assume you have at least a passing acquaintance with AI programming concepts and terms.

Since Joshua is very closely integrated into its Symbolics Common Lisp environment, we also assume that you are familiar with the Genera facilities, Symbolics Common Lisp, and the Zmacs editor. Extensive documentation on these areas is provided elsewhere in the Symbolics documentation set.

The Joshua documentation consists of two manuals,


This division reflects a different task orientation as well as a different stage of familiarity with AI programming.

1.2.1. The User’s Guide to Basic Joshua

*User’s Guide to Basic Joshua* gives you everything you need to know to start developing Joshua applications.

The goal of this manual is to let you develop a feel for Joshua, together with the ability to do all of the most common programming tasks using only Joshua’s built-in facilities. For this reason, we will work with a subset of the Joshua commands, and with only those top-level Protocol functions that are directly usable for standard operations and that do not represent special-purpose functions or subcomponents of the protocol. The manual is designed to answer questions such as "How do I interact with the Joshua database?", "What kind of reasoning operations can I do with Joshua?", and "How does Joshua handle Truth Maintenance?"

*User’s Guide to Basic Joshua* covers the following topics:

- Predications
- Rules and Inference
- Questions
- Unification
- Using Joshua Within Lisp Code
- Justification and Truth Maintenance

*User’s Guide to Basic Joshua* is organized into a conceptual discussion, followed by the "Basic Joshua Dictionary". This is an alphabetized dictionary of reference entries for the basic subset of Joshua symbols that let you program Joshua’s default structures. Each entry provides a complete description of a single Joshua function or command, its syntax, what it returns, examples of its use, and cross-references to related functions or commands.
1.2.2. The Joshua Reference Manual

*Joshua Reference Manual*, the companion volume to this, is a much deeper presentation of the concepts you need to understand in order to put Joshua to the fullest possible use. This manual is for experienced programmers who need to write tailored, optimized applications.

Topics presented earlier are revisited here in more depth, with cross-references to the earlier manual. The focus here is on all levels of the Joshua Protocol functions, so that you can see how each works. With this understanding you can make changes to any component for efficiency. *Modeling*, that is, tailoring any of the Joshua components to your application, is introduced in some detail, including examples of modeling data structures, rule storage, and writing your own TMS.

These are some of the topics covered in

*Joshua Reference Manual*:

- The Database Protocol
- Trigger Indexing Protocol
- The Rule Compiler
- The TMS
- Modeling

As in the earlier volume, the conceptual discussion is followed by the "Joshua Language Dictionary", an alphabetized dictionary of reference entries for all Joshua symbols, methods, and commands. Each entry provides a complete description of a single Joshua function or command, its syntax, what it returns, examples of its use, and cross-references to related functions or commands.

Among the major goals of the documentation is to give you fast access to the information you need, to minimize your reading, and to let you work directly with the information presented. Topics are clearly cross-referenced to point you from basic to advanced coverage, as well as to related topics; important information is summarized visually in tables or figures for quick reference; the majority of the examples in the manual show output, and they can be yanked from Document Examiner into the Lisp Listener so that you can experiment with them yourself.
2. Getting Started with Joshua

2.1. Setting up the Joshua Context and File Attributes

In order to get to a correct Joshua package and to inform the Lisp reader that you are expecting it to deal properly with Joshua syntax, you must set your working context to Joshua.

From the Lisp Listener, enter:
   Set Lisp Context Joshua
This sets the current context to use Joshua syntax, and sets the current package to Joshua. For information on Lisp contexts: See the section "Set Lisp Context Command" in Genera Handbook.

To set your file attribute list, enter the following from a Zmacs buffer:
   * Motion: Create Initial Joshua Attribute List
This creates an attribute list similar to the one below:
   ```
   ;; -* Mode: Joshua; Package: JOSHUA-USER; Syntax: Joshua; Vsp: 0 -*.
   ;; Created 3/18/88 10:24:45 by Covo running on LADY-PEREGRINE at SCRC.
   ```

For information on file attribute lists: See the section "Buffer and File Attributes in Zmacs" in Editing and Mail. You are now ready to use Joshua.
3. Overview of Joshua

Since the core of Joshua is a rule-based inference language, its chief building blocks might already be familiar to you from other AI programming languages. We will briefly review these elements as we present the Joshua context and terminology.

One can think of Joshua as having five major components:

* Predications are ways of expressing knowledge; they are often called statements, facts, or assertions in other languages, and we occasionally use these terms as well.

* The database is a collection of predications and related information that the system remembers;

* Rules are ways of expressing and remembering relationships among predications, as well as procedural knowledge;

* The Protocol of Inference is the mechanism that integrates these components and performs the reasoning;

* The Truth Maintenance System (TMS) keeps track of the reasoning that was used so that it can:

  * Maintain explanations of Joshua’s reasoning
  * Maintain the logical consistency of the system

Basic programming in Joshua consists of supplying predications and rules and determining how to use this knowledge to deduce additional knowledge, or how to use it to answer questions. We cover these topics in the present manual.

Advanced programming in Joshua consists of modeling, that is, tailoring appropriate parts of the Joshua Protocol routines to your particular application; you might want to model for any number of reasons (to increase efficiency, incorporate specialized tools, access external databases, and the like). We cover these topics in the Joshua Reference Manual, the companion volume to this.

We begin with a summary of the Joshua protocol.

3.1. Some Basic Joshua Protocol Functions

The Joshua Protocol of Inference consists of some 30 generic functions broken down into five functional groups:

* **Database Interface**: Manages addition/deletion of facts to the database
* **Truth Maintenance System (TMS) Protocol**: Manages deductive dependencies
* **User Interface Protocol**: Manages interaction with the user
* **Rule Indexing Protocol**: Manages the rule database
* **Rule Customization Protocol**: Manages the rule compiler
Because of the modular nature of the Joshua Protocol, most Protocol functions are broken into subcomponents, each of which performs a specific part of the function’s contract. This modularity lets you fine-tune and localize changes to functionality at any level, letting them remain transparent to the end user.

In this manual we cover some basic top-level protocol functions in the Database Interface, User Interface, and TMS groups. The functions for Rule Indexing and Rule Customization are used only for modeling, and are not of immediate interest here. Similarly, we will not be concerned with the details of modularization in each of the five protocol groups, since this information is only relevant for modeling.

Figure basic-protocol is a depopulated picture of the basic protocol functions we are covering here. For the complete protocol: See the section “A Figure of the Joshua Protocol of Inference”, page A-Figure-of-the-Joshua-Protocol-of-Inference-SECTION.

Figure 1. A Basic Subset of the Joshua Protocol

Here is a summary explanation of each of these functions.
**Database Interface**

- `joshua: tell`: Installs new information
  - `joshua: justify`: Gives the TMS a reason for believing the predication

- `joshua: ask`: Retrieves known or implied data

- `joshua: untell`: Removes a single fact that `joshua: tell` put into the database
  - `joshua: unjustify`: Manages the TMS issues related to removing the fact

- `joshua: clear`: Removes all facts from the database

**TMS Protocol: Manages Deductive Dependencies**

- `joshua: justify`: Installs a new TMS justification
- `joshua: unjustify`: Removes a TMS justification
- `joshua: support`: Finds the set of facts or assumptions that a statement depends on

**User Interface**

`joshua: say`: Prints out the meaning of predications in natural language, or something similar.

For a subset of Joshua symbols introduced in this manual: See the section "List of Basic Joshua Symbols", page List-of-Basic-Joshua-Symbols-SECTION.

### 3.2. List of Basic Joshua Symbols

Here is the list of functions and commands whose basic functionality we present in this manual.

- `joshua: ask`
- `joshua: ask-database-predication`
- `joshua: ask-derivation`
- `joshua: ask-query`
- `joshua: ask-query-truth-value`
- `joshua: clear`
- "Clear Joshua Database Command"
- `joshua: contradictory`
- `joshua: copy-object-if-necessary`
- `joshua: define-predicate`
joshua:defquestion
joshua:defrule

joshua:different-objects
"Disable Joshua Tracing Command"
"Enable Joshua Tracing Command"
"Explain Predication Command"
joshua:explain
joshua:*false*
joshua:graph-query-results
joshua:graph-tms-support
joshua:known
joshua:make-predication
joshua:map-over-database-predications
joshua:predication
joshua:predicationp
joshua:print-query
joshua:print-query-results
joshua:provable
"Reset Joshua Tracing Command"
""\Say\"
joshua:say
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joshua:with-unbound-logic-variables
joshua:with-unification

To begin using Joshua, you need to work with *predications*, learning how to define, store, look up, and delete them. For these and related topics: See the section "Predications and Predicates", page Predications-and-Predicates-SECTION.
4. Joshua Predications

AI and Joshua programming activities involve working with facts that represent knowledge. Programs build and store facts and reason with them, which leads to the building and storing (or removal) of other new facts to reason with.

Joshua facts are called *predications*. This chapter covers the essentials of building and storing predications. We also introduce related topics such as truth values associated with predications, and the use of logic variables to build *predication patterns*. Reasoning with predications is the province of rules and is covered in that context: See the section "Rules and Inference", page Rules-and-Inference-SECTION.

4.1. Predications and Predicates

Macro: 

```
joshua:define-predicate
```

Function: 

```
joshua:undefine-predicate
```

Command: Show Joshua Predicates

Zmacs Command: `n-k` Kill Definition

*Predications* are statements about the world. When predications are stored in a *database*, the knowledge they embody becomes available to the system, to be manipulated by rules.

Terminology note: Occasionally we use the terms *statement*, *assertion*, *fact*, and *belief* synonymously with *predication*. These are broadly equivalent terms that have wide currency in AI programming.

The protocol lets you manipulate predications in numerous ways. You can, for example:

- Insert them into a database
- Look them up in a database
- Infer them by rules
- Supply justifications that are remembered in the database

Bear in mind, also, that predications can be manipulated like other ordinary Lisp-world objects (you can store them in arrays, print them, *joshua::accept* them, and so on). In fact, predications are just flavor instances.

These are examples of predications:

```
[healthy John]
[has-pneumonia John]
[author-of (poems plays) Marlowe]
```

Predications consist of *predicates* such as healthy, has-pneumonia and author-of in the above example, together with their arguments, if any. A predication is always enclosed in brackets [ ].
Predicates thus are names of relationships, and they organize the knowledge in a
predication and express relationships among its parts. A predicate is always the
first item in a predication; it can take zero or more arguments, depending on how
you define it.

You must define a predicate before using it in a predication. Defining the predi-
cate sets up its format and specifies the kinds of information you expect to see af-
after the predicate name (things like the arguments and, optionally, any customized
ways of controlling the predicate’s behavior); At runtime the system checks your
predication’s pattern against the format you set up in the predicate definition, and
notifies you if there is a discrepancy.

Use the macro joshua:define-predicate to define new predicates. As an example,
we’ll define a predicate, healthy, to take a single argument, object. This predicate
uses the default implementations for data storage and representation, so no further
arguments are needed in the definition.

(define-predicate healthy (object))

Once a predicate has been defined, we can use it in any number of predications
with different arguments. This makes it possible to organize related knowledge by
grouping it together.

Note that you can use backquote with the bracket notation for predications. For
example:

`[healthy ,(find-a-healthy-thing)]`

Here are some examples of predications:

[healthy vegetables]
[healthy long-vacations]
[healthy "an apple a day"]

Since there is no one set way of thinking about a world, we can express a given
piece of knowledge in a variety of different predications, depending on what as-
pects of it are important to our problem. We can, for instance, use several diffe-
rent predicates to express the fact that a person is healthy, or that a person is ill:

[healthy John]
[has-health John]
[has-pneumonia John]
[illness John pneumonia]
[pneumonia John]

Or we can relate health and illness conceptually as reflecting different states of
being, and generalize them to a single condition by using a single predicate:

[has-condition John health]
[has-condition John pneumonia]

Or, if John is the only patient whose condition interests us, we can omit explicit
mention of John’s name from the statements altogether.

[has-condition health]
[has-condition pneumonia]
Thus, your choice of predicates reflects your analysis of the problem, the kinds of reasoning you expect your program to do, and the amount of knowledge you need to express explicitly or implicitly. Since predicates are your vocabulary for defining and exploring a knowledge domain, predicate selection is a major problem-solving challenge; choosing the correct set of predicates to model your knowledge can greatly simplify your programming task. Poor choices can lead to clumsy programs with too many rules.

You can "undefine" predicate definitions from either Zmacs or the Lisp Listener. (Undefining a predicate is a rare event that corresponds to a major reorganization of the program.) To remove a predicate definition while you are in Zmacs, place the cursor on the predicate definition and use the extended command \texttt{C-x K} Kill Definition. In the Lisp Listener, use the function \texttt{joshua:undefine-predicate}.

For example:

\begin{verbatim}
(undefine-predicate 'healthy)
\end{verbatim}

Note that while \texttt{joshua:undefine-predicate} removes the \textit{predicate definition}, any \textit{predications} built with this predicate remain in the world until you explicitly remove them. However, almost any attempt to use these predications will cause an error. (See the section "Removing Predications From the Database", page Removing-Predications-From-the-Database-SECTION.)

To find out what predicates are defined in your current world, use the command Show Joshua Predicates. This prints the predicate names and their arguments. There are command options to let you tailor the display. Please consult the dictionary entry for this command.

Free-floating predications such as those we’ve just built are not usable for reasoning operations because the system does not keep track of them. To be on record, predications must be collected in a database that the system recognizes and remembers. Also, once predications are in a database, the system stores additional information about them, such as their truth value. The next section deals with these topics. See the section "Predications, Truth Values, and the Database", page Predications,-Truth-Values,-and-the-Database-SECTION.
Advanced Concepts Note:

In Joshua, predicates are an implementation-organizing as well as a knowledge-organizing tool. Since you can organize related facts in similar fashion by consistent predicate choice, it can also be convenient to implement them in the same way. So predicates, rather than their arguments, have specific protocol implementations associated with them. This means that in Joshua you can, if you wish, deal with each predicate independently, and give it its own unique way of interacting with the protocol; while defining a predicate you can specify such things as where predications using it are to be stored (this could be in any number of places, even on separate computer systems), or whether or not justifications for it are to be remembered. For some predicates you might want to specify algorithmic methods for answering questions about a problem, while for others you might request reasoning methods. In other words, you can model the behavior of each predicate, mixing and matching various built-in models, or adding your own.

The various models you might use are also independent of each other. You can, for example, define a TMS model (mostly) without affecting the storage model; in this way you need only change those portions of the protocol directly affected by your model. Similarly, your modeling choices require no modification of other knowledge-level structures, such as rules.

This flexibility is one of Joshua’s major advantages. However, the extent to which you exercise it is strictly up to you, since it is possible to program effectively in Joshua using only the defaults.

All this is implemented using the generic functions and flavors mechanism of Symbolics Common Lisp.

We cover customized predicate models in Joshua Reference Manual.

4.2. Predications, Truth Values, and the Database

To make a predication available to Joshua, you joshuatell it into a database. A database is an extensible collection of predications together with associated information about each predication. Predications are generated by the programmer or as a result of reasoning done by rules; they are entered into the database by the programmer (or by the system if they are inferred by rules). A database thus contains all the facts Joshua knows about at this point in time, and it changes dynamically with new facts being added or removed as a result of inference or computation.

Predications in the database differ from free-floating predications first because they are remembered by the system, and second, because they have additional information, such as truth values, associated with them. (See the section "Truth Values", page Truth-Values-SECTION.)

As a Joshua user, your main interaction with the database takes place via two functions, joshuatell and joshua:ask. joshuatell stores statements into the
database. (See the section "Entering and Displaying Predications in the Database", page Entering-and-Displaying-Predications-in-the-Database-SECTION.) joshua:ask, among its other functions, looks up knowledge in the database. (See the section "Querying the Database", page Querying-the-Database-SECTION.)

Database operations include the following:

- Inserting predications into the database
- Displaying the database contents (text, or graphic format)
- Looking for predications in the database
- Removing predications from the database

This section covers the basics of the above operations. We also introduce the related topic of truth values that are associated with predications in the database.

**Advanced Concepts Note:**

The Joshua database protocol provides a default mechanism for data representation and storage. (See the section "The Joshua Database Protocol" in Joshua Reference Manual.) If you provide an alternate implementation of the relevant components of joshuatell and joshua:ask, you can customize data indexing for your application. This will, for example, let you access knowledge in any existing database, or combinations of databases in various locations. The Joshua concept of a virtual database expresses this idea. (See the section "Customizing the Data Index" in Joshua Reference Manual.)

### 4.2.1. Entering and Displaying Predications in the Database

**Functions:**

- joshuatell
- joshua:untell
- joshua:clear

**Commands:**

- Show Joshua Database
- Clear Joshua Database

We use the Command Processor command Show Joshua Database to display the contents of the database. An empty database generates the following display:

```
Show Joshua Database
True things
None
False things
None
```

The display headings "True things" and "False things" are for grouping predications that the system knows to be true (those with a truth value of joshua:*true*), and predications that the system knows to be false (those with a truth value of joshua:*false*), respectively. (See the section "Truth Values", page Truth-Values-SECTION.)

Let’s define a new predicate:
(define-predicate has-eye-color (person color))

Now assume you or a rule application generate a predication such as:

[has-eye-color fred green]

This new predication is "free-floating", that is, it does not automatically become part of the database, unless you explicitly enter it there by giving it to joshuatell.

joshuatell takes the "free-floating" predication, and returns a database predication, that is, an object that it stored in the database and that can be accessed and reasoned with. It is important to note that this object, because it is stored with information about itself, is conceptually (and often physically) distinct from the "free-floating" predication you supplied as an argument to joshuatell.

Whenever you need to do something with a stored predication (for example, to remove it from the database), you need to access the actual object that is stored. Later on we’ll discuss ways of retrieving this object.

For clarity we use the terms database predication, whenever it is important to refer to the actual object that is stored in the database rather than to a predication in more general terms.

Using joshuatell is very straightforward; you apply the joshuatell function to the predication, as in this example:

(tell [has-eye-color Fred green])

[HAS-EYE-COLOR FRED GREEN]

T

joshuatell returns the predication object it stored in the database. It also returns a boolean value indicating whether the predication is being inserted for the first time, as in the example above (T), or whether it already existed in the database (nil). This information is of interest primarily to the Truth Maintenance System, so we won’t always reproduce it in the output from our examples; we mention it here only because you’ll be seeing this every time you enter a joshuatell.

Now display the database once more, and you will find that it includes the predication we have just entered with joshuatell. [has-eye-color fred green] appears under "True things", indicating that Joshua knows the statement represented by the predication to be true.

Show Joshua Database
True things

[HAS-EYE-COLOR FRED GREEN]
False things
None

To indicate that a statement is false, prefix the predication representing that statement with the predicate joshua::not. For example, here we tell Joshua that we know Jane doesn’t have green eyes.

(tell [not [has-eye-color Jane green]])

→ [HAS-EYE-COLOR JANE GREEN]

T

Here joshuatell returns the database predication as usual. The Lisp printer pre-
fixes the predication with the Not-sign logic symbol, (¬), denoting that the statement is false.

In the database display the predication appears under the heading "False things" without any prefix.

Show Joshua Database
True things
[HAS-EYE-COLOR FRED GREEN]
False things
[HAS-EYE-COLOR JANE GREEN]

Let's define a few more predications, add them to the database and then display it.

(define-predicate loves (lover beloved))
(tell [loves bonzo jane])
(tell [not [loves jane bonzo]])

Show Joshua Database
True things
[HAS-EYE-COLOR FRED GREEN]
[LOVES BONZO JANE]
False things
[LOVES JANE BONZO]
[HAS-EYE-COLOR JANE GREEN]

The database display always shows the actual predications currently stored there. Click on any item in the display for use in Joshua commands that require database predications.

The Show Joshua Database command by default displays the entire contents of the database. This can become cumbersome with very large databases. Later on, we’ll show how you can limit the database display by requesting only specific patterns or specific truth values. See the section "Predications and Logic Variables", page Predications-and-Logic-Variables-SECTION.

You can remove predications from the database as well as enter them.

See the section "Removing Predications From the Database", page Removing-Predications-From-the-Database-SECTION.

**Advanced Concepts Note:**

The default database is implemented as a discrimination net stored in the value of the global variable `ji:*data-discrimination-net*`. See the section "Joshua's Default Database: the Discrimination Net" in *Joshua Reference Manual*.

### 4.2.2. Removing Predications From the Database

There are several ways of removing predications from the database. Note that predicate definitions are not eliminated, but only the predications built with them.
The function \texttt{joshua:untell} removes a single predication from the database, and frees up some related storage. As its name implies, \texttt{joshua:untell} is the opposite of the function \texttt{joshua:justify}.

Another way to remove a predication from the database conceptually is to use the function \texttt{joshua:justify} if your application includes a TMS model. \texttt{joshua:justify} differs from \texttt{joshua:untell} in important respects. See the section "Retracting Predications with \texttt{joshua:justify}", page Retracting-Predications-with-JOSHUA:UNJUSTIFY-SECTION. The command Clear Joshua Database provides a convenient interface to the \texttt{joshua:untell} function. It asks the database for all predications matching those specified by the arguments, prompts you for confirmation, and \texttt{joshua:untell}s each predicate. (It also allows you to undefine all the Joshua rules, resulting in a fresh Joshua environment.) You can specify predications matching a certain pattern, all predications, or none. You can also specify whether or not to remove predications that match the pattern specified but have the opposite truth value. (Using logic variables to build predication patterns is described elsewhere: See the section "Predications and Logic Variables", page Predications-and-Logic-Variables-SECTION.)

Command: Clear Joshua Database (predications, All, or None [default All]) All Clear all predicates from the database? [default Yes]: Yes

To remove all predications from the database, use the function \texttt{joshua:clear}.

(clear)
Show Joshua Database

True things
None
False things
None

4.2.2.1. Removing a Database Predication with Untell

In the section "Entering and Displaying Predications in the Database", we mentioned that several Joshua operations require the actual predication object that is stored in the database. \texttt{joshua:untell} is one such operation. There are several ways of retrieving a predication object.

One method is to display the database contents and use the mouse to pick up the object you want to operate on. (Position the mouse cursor at the start of the line displaying the object, and click left on the mouse.) This picks up the object and positions it wherever your cursor is. The object so moved displays in capital italics (see example below).

For example, to remove the database predication \texttt{[has-eye-color Jane grey]}, display the database, enter \texttt{(joshua:untell)}, move the cursor one space to the right, and click left on the target predication from the database display. Joshua inserts the predication at the cursor position (after the \texttt{joshua:untell}). Now enter a closing right parenthesis to end the command.
Show Joshua Database (matching pattern [default All]) All

True things

[HAS-EYE-COLOR JANE GREY]
[HAS-EYE-COLOR FRED GREEN]
[LOVES BONZO JANE]
[FAVORITE-MEAL MONKEYS BANANAS]

False things

[LOVES JANE BONZO]

(untell [HAS-EYE-COLOR JANE GREY]) ; enter the predication by clicking left
; on the object in the database display

NIL

Show Joshua Database

True things

[HAS-EYE-COLOR FRED GREEN]
[LOVES BONZO JANE]
[FAVORITE-MEAL MONKEYS BANANAS]

False things

[LOVES JANE BONZO]

After the joshuauntell the database no longer contains the predication [has-eye-color Jane grey].

Another way to get at the predication object is to specifically save the result of the joshuastell. For example:

(setq db-object-1 (tell [loves Narcissus Narcissus]))
[LOVES NARCISSUS NARCISSUS]

Show Joshua Database (matching pattern [default All]) All

True things

[HAS-EYE-COLOR FRED GREEN]
[LOVES BONZO JANE]
[LOVES NARCISSUS NARCISSUS] ; new predication object is added
[FAVORITE-MEAL MONKEYS BANANAS]

False things

[LOVES JANE BONZO]

(untell db-object-1) ; db-object-1 is [loves Narcissus Narcissus]
NIL
Show Joshua Database (matching pattern [default All]) All
True things
[HAS-EYE-COLOR FRED GREEN]
[LOVES BONZO JANE]
[FAVORITE-MEAL MONKEYS BANANAS]
False things
[LOVES JANE BONZO]
The predication [loves Narcissus Narcissus] is no longer in the database.
Other ways of retrieving a database predication involve database queries and using
predication patterns. We'll get to these topics later.
See the section "Querying the Database", page Querying-the-Database-SECTION.
See the section "Predications and Logic Variables", page Predications-and-Logic-
Variables-SECTION.

4.2.3. Truth Values

A truth value is a value denoting what the system currently knows about the truth
state of a database predication. A truth value becomes associated with a predica-
tion when Joshua adds the predication to the database. Predications commonly
change their truth value as knowledge is updated.

A predication can be in any one of four possible truth states:

- joshua:*true* (appears under "True things" in the database display)
- joshua:*false* (appears under "False things" in the database display)
- joshua:*unknown* (does not appear in the database display)
- joshua:*contradictory* (a transient state; does not appear in the database dis-
  play)

The truth value of a predication can be manipulated directly by the user. It can al-
so be manipulated by the TMS (Truth Maintenance System). We cover the TMS in
a separate section, touching on it here only as necessary. (See the section "Justifi-
cation and Truth Maintenance", page Justification-and-Truth-Maintenance-
SECTION.)

Joshua has a three-valued logic. A statement (predication) is joshua:*true* if its
arguments are believed to satisfy the predicate, and joshua:*false* if it is known
that they do not.

When making new predications, joshua::not is prefixed to a fact that is known to be
joshua:*false*, as in [not [loves Jane Bonzo]].

If a fact is neither known to be joshua:*true*, nor known to be joshua:*false*, it is
joshua:*unknown*. In languages such as Prolog, if a fact cannot be proved to
be true, it is assumed to be false. Joshua does not subscribe to this so-called
"Closed World Assumption".
A predication is (or becomes) **joshua:*unknown* when there is no valid reason that supports it. (For example, if a predication was **joshua:*true* but the reason underlying this truth status is removed, the predication becomes **joshua:*unknown*.) The predication does remain physically in the database as an efficiency measure; however, since most reasoning operations look for a truth value of **joshua:*true* or **joshua:*false*, a predication with a truth value of **joshua:*unknown* is conceptually not seen. That is to say, from the point of view of the reasoning process, a predication that has a truth value of **joshua:*unknown* is indistinguishable from one that is not in the database at all. When a predication's truth value changes from **joshua:*unknown* to either **joshua:*true* or **joshua:*false*, the predication is once more “visible” and used in the inferencing process.

The **joshua:*unknown* truth state is used primarily by the TMS as it modifies the database for logical consistency.

As its name implies, the truth value of a predication is **joshua:*contradictory* if there are reasons to believe that the predication is both true and false at the same time. This truth value also is primarily meaningful in conjunction with a TMS: if a TMS is present, it detects contradictory truth states and does not allow them to remain in the system. If no TMS is present, contradictions go undetected, and the most recent truth value is remembered. See the section ‘Revising Program Beliefs’, page Revising-Program-Beliefs-SECTION.

The truth value of a predication can change either because you explicitly modify it, or as a result of the monitoring activities of the TMS.

If you have not included a TMS in your predicate definition, you can change the truth value of a predication from **joshua:*true* to **joshua:*false* by reasserting the predication with **joshuatell**. In this case, the system simply accepts the most recent fact you **joshuatell** it, with no regard to logical consistency. Here, for example, are two **joshuatell** statements.

```prolog
(tell [loves Medea Jason])
[LOVES MEDEA JASON]
T

(tell [loves Medea her-children])
[LOVES MEDEA HER-CHILDREN]
T
```

Here's the database display containing these **joshua:*true* predications.

- Show Joshua Database
- True things
  - [LOVES MEDEA HER-CHILDREN]
  - [LOVES MEDEA JASON]
- False things
  - None

Say we now discover that Medea no longer loves her husband, Jason, and we want to modify the database to reflect this new fact. To change the truth value of
[loves Medea Jason] from joshua:*true* to joshua:*false*, re-enter the predication into the database, this time prefixing it with joshua::not:

(tell [not [loves Medea Jason]])

- [LOVES MEDEA JASON]

NIL

A Truth Maintenance System would insist that Medea cannot both love and not love Jason simultaneously. Since we have no TMS in this example, Joshua accepts this new version of the predication without noting the contradiction. So the latest version entered ([not [loves Medea Jason]]), simply supersedes the old version ([loves Medea Jason]).

Note that the result of these two joshuatell operations about Medea and Jason is the same database entry in each case. The Lisp printer prints the predication with the Not-sign symbol, (¬), the second time, to indicate that it changed its truth value. The second joshuatell also returns the boolean nil, to indicate that the predication already existed in the database.

Show Joshua Database
True things
[LOVES MEDEA HER-CHILDREN]
False things
[LOVES MEDEA JASON]

To change a predication’s truth value from joshua:*false* to joshua:*true*, use joshuatell to reassert the predication into the database, without the joshua:not prefix.

Modifying truth values without using a TMS can introduce inconsistencies into your database. Since the system does not follow up the logical consequences of a change in truth value, any such inconsistencies remain undetected.

Suppose, for example, that the second predication in our database, [loves Medea her-children], was deduced by some forward chaining rule, based on the belief [loves Medea Jason]. That is, the rule concluded that Medea loves her children because of the fact that she loves Jason.

Since we have now told the system that Medea does not love Jason, the conclusion that she loves her children no longer follows. (As we know, the mythical Medea went through some such thought process, killing her children to show her hatred of Jason.)

Without a TMS the system ignores these issues. Specifically, as we see from the last database display, it allows the fact [loves Medea her-children] to remain as a valid belief, even though it has lost its reason for being believed.

The TMS on the other hand follows up the logical consequences of a change in truth value and signals any resulting inconsistencies in the database. In this situation the TMS might have prevented infanticide, since asserting [not [loves Medea Jason]] would immediately cause a warning that Medea's love for her children is no longer a true fact. The TMS would then request you to correct the contradiction by removing one of the inconsistent statements.
As we can see, due to the TMS efforts to maintain a logically consistent database, adding a new predication can result in changed truth values for several existing predications. See the section "Revising Program Beliefs", page Revising-Program-Beliefs-SECTION.

The Joshua predicates joshua:known and joshua:provable help us determine what we know. joshua:provable tells us whether a fact is currently known to have a particular truth value. joshua:known tells us if we have knowledge of a fact (regardless of whether the fact is joshua:*true* or joshua:*false*). Both these predications are more useful in their negative form ( [not [provable [...]]], and [not [known [...]])]. For some examples of their use, please consult the dictionary entries.

Predications combined with the standard predicate calculus connectives and, or, and not can be joshua:*true* or joshua:*false*, according to the rules in the truth tables listing the truth value for each possible combination of argument truth values. See the section "Predications and Logical Connectives", page Predications-and-Logical-Connectives-SECTION.

### 4.2.4. Querying the Database

**Functions:**

- joshua:ask
- joshua:graph-query-results
- joshua:print-query
- joshua:print-query-results
- joshua:say-query

Here are some predicate definitions and joshua:tell statements that create a simple database.

```
(define-predicate has-condition (object condition when))
(define-predicate adult (person))

(clear)

(tell [has-condition Fred measles today])
(tell [adult John])
(tell [adult Fred])
```

We need a way to query this database so that we can find out what Joshua knows. The generic function joshua:ask does this. The main purpose of joshua:ask is to find solutions to queries such as “Is John an adult?” “Who are the adults we know about?” “Do any adults currently have diseases?” It does this by looking in the database and by invoking backward rules (and questions) to derive answers by reasoning.

Because joshua:ask is your main interface to the knowledge and rule structures, many of its typical uses depend on concepts we have not yet covered; so, for the moment, we introduce joshua:ask in its most basic form. Further examples of its use appear throughout the rest of the manual.

joshua:ask works as follows: you give it a query pattern, that is, some pattern like "John is an adult" that you want to validate. The query pattern is a predication that may, or may not, contain logic variables. (Logic variables are covered elsewhere: See the section "Predications and Logic Variables", page Predications-and-Logic-Variables-SECTION.)

joshua:ask looks for answers to the query, collecting information about its search process as it goes along.

As soon as it finds an answer, joshua:ask passes a list containing the answer, together with the related information about how the answer was derived, to a function called a continuation. The continuation can do what it likes with the answer.

The information received from joshua:ask resides in the single continuation argument called backward-support. backward-support is a list whose elements are: the answer to the query, the truth value of the query, and the derivation for the answer - that is, what line of reasoning was followed to determine the answer. Note that although the support is only one element of this list, we refer to the entire list as backward support. Also note that this list is usually stack-consed.

There are two ways of dealing with the backward support information. One way is to use Joshua’s accessor functions in the joshua:ask continuation. These functions break up the backward support into separate elements and let you interpret these elements as you wish. We discuss accessor functions in the dictionary entry for joshua:ask.

An easier way of processing answers to queries is to use Joshua’s convenience functions in the joshua:ask continuation. These functions extract parts of the backward support and interpret it for you. Generally we’ll be using the convenience functions throughout this first manual. Here is a list of these functions. The first two deal with the answer to the query, and the other two interpret the reasons for the answer.

joshua:print-query  Extracts and prints the answer to the query.

joshua:say-query  Extracts the query and displays it in formatted form. See the section "Formatting Predications: the SAY Method", page Formatting-Predications:-the-SAY-Method-SECTION.

joshua:print-query-results  Extracts and prints the support for the answer. The support varies depending on whether the answer was found in the database or derived from rules.

joshuagraph-query-results  Draws a graph of the query support, labeling rules and questions.
An additional convenience function, `joshua:map-over-database-predications` (instead of backward rules) is useful when you want to find answers only in the database and do some operation with those predications that you find. We mention this function here for completeness; for an example of its use: See the section "Predications and Logic Variables", page Predications-and-Logic-Variables-SECTION.

When Joshua is asked a question such as, `(ask [has-condition Fred measles today] ...)`, it tries to derive the answer from the database before looking for backward rules and questions to run. Here we discuss the database lookup.

When trying to answer a query, `joshuask` (conceptually) goes through the whole database, trying to match the predication pattern in the query against all predications in the database. We refer to pattern matching in Joshua as `unification`. A query pattern `unifies` with a predication either because the two patterns are `equivalent` (that is, they look the same and both predicate and arguments match exactly), as in

```
Query pattern: (ask [has-condition Fred measles today] ...)
Database predication: [has-condition Fred measles today]
```

or because appropriate values can be substituted for any variables in the two patterns to make them equivalent. See the section "Pattern Matching in Joshua: Unification", page Pattern-Matching-in-Joshua::Unification-SECTION.

When a query pattern succeeds in unifying with a predication, the query pattern becomes temporarily `instantiated` as that predication. That is, the logic variables in the query pattern are temporarily bound to the corresponding values in the database predication. At this point Joshua calls the `joshua:ask` continuation to process this answer. (The continuation is called once for each time the query can be satisfied.)

Here's the query, using the convenience function `joshua:print-query` in the continuation. Since the query can be answered (by finding a matching predication in the database), the continuation is called to process the answer. `joshua:print-query` displays the query with its variables instantiated, ignoring all the rest of the backward support information.

```
(ask [has-condition Fred measles today] #'print-query)
[HAS-CONDITION FRED MEASLES TODAY]
```

For illustration, let's use the continuation function `joshua:print-query-results` to show how Joshua derived the answer to our query.

```
(ask [has-condition Fred measles today] #'print-query-results)
[HAS-CONDITION FRED MEASLES TODAY] succeeded
[HAS-CONDITION FRED MEASLES TODAY] was true in the database
```

As expected, Joshua tells us that the answer was found in the database, and returns the support for the satisfied query, that is, the actual database object that matched the query. When the answer is derived from rules, `joshua:print-query-results` traces the chain of reasoning that `joshua:ask` followed.
What would a graph of the support look like? Figure ask-graph-1 shows the graph
drawn by joshua:graph-query-results for the above query.

```
(ask [has-condition Fred measles today] ’graph-query-results)
```

Figure 2. Simple example using function graph-query-results

The graph tells us that the predication was in the database, and displays the
database predication inside a rectangle. Rectangles in support graphs denote
queries that were satisfied by the database (rather than by rules, or questions).

If you joshua:ask about a predication that is not in the database, and that cannot
be derived from rules or questions, the system does not call the continuation and
nothing is returned.

```
(ask [adult mary] ’print-query)
```

You can trace joshua:ask and joshua:tell operations in Joshua. See the section
"Tracing Predications", page Tracing-Predications-SECTION.

Asking questions about specific facts is of limited usefulness. It would be prefer-
able to be able to ask the system more general questions, such as asking it to
identify all the adults in our database. In other words, we need to joshua:ask the
system questions using variables for which it will try to find the appropriate val-
ues. Joshua provides logic variables for this purpose. See the section "Predications
and Logic Variables", page Predications-and-Logic-Variables-SECTION.

### 4.3. Predications and Logic Variables

A logic variable is a special type of object used by Joshua and recognized by Sym-
bolics Common Lisp. Building predications with logic variables instead of constant
values lets you create patterns; patterns let you move from the particular object to
the level of generalized queries or statements.

If we are looking for a listing of all the adults the system knows about, we could
joshua:ask Joshua about every single person we can think of:

```
(ask [adult Fred] ’print-query)
(ask [adult John] ’print-query)
```

Or (preferably), we could joshua:ask the question in a more general form such as,
"Name everyone who is an adult". For example:

```
(ask [adult ?x] ’print-query)
```
[adult ?x] is a query pattern containing a predicate followed by a logic variable, ?x. We don't know who ?x might stand for, but Joshua can give us this information. A logic variable stands for something (in this case for a person); Joshua determines what the something is, depending on the query.

Joshua recognizes both the question-mark (?) and the equivalence symbol (=) (typed symbol ~), by itself or followed by a symbol, as a logic variable:

[adult ?]
[adult =]
[adult ?person]
[adult =person]
[adult ?x]
[adult =x]

Regardless of whether you used logic variables or constants to enter your predication in the database, you can use logic variable arguments in an joshuask query. For example:

(define-predicate healthy (object))
(tell [healthy John])

(ask [healthy ?person] #'print-query)
[HEALTHY JOHN]

Choosing descriptive names for logic variables within query patterns is useful for documenting the predicate’s arguments and their order.

Here are some joshuatell statements about adults:

(define-predicate adult (person))
(tell [adult Fred])
(tell [adult John])
(tell [adult Rose])

If we joshuask the system to find all the adults in the database (ask [adult ?person] ...), it does so by first matching the predicate adult, and then successively substituting the value of the predicate’s argument (here a name), for the logic variable ?person in the query. So queries with logic variables are satisfied by finding correct matches for the variables.

Since we want joshuask to look for solutions only in the database, we should prevent it from looking at backward rules or questions as well. We do this by using :do-backward-rules nil. For example:

[ADULT ROSE]
[ADULT JOHN]
[ADULT FRED]

As we can see, the continuation is invoked three times, once for each time the query was satisfied. Let’s walk through this example.
At the outset of the query, the logic variable ?person in the query pattern [adult ?person] does not stand for any particular value; that is, it is as yet uninstantiated. An uninstantiated variable can match any object. So when Joshua searches through the database with an uninstantiated variable as an argument in the query, this variable is allowed to match any argument in the same position in the database predication, provided, of course, the predicates match.

The logic variable ?person in our example can match any first argument in a predication beginning with the predicate adult. When Joshua finds the first predication whose predicate is adult (here [adult Rose]), the argument, Rose, is temporarily substituted for the logic variable ?person in the query, so that ?person stands for Rose (becomes instantiated as Rose). In this way the query pattern matches the database predication. Joshua calls the continuation to print this answer as requested.

Once the continuation has finished executing, Joshua discards its temporary value assignments, resets logic variable ?person to be once more uninstantiated, and continues the search through the database to find the next match for the query, that is, another name that ?person could stand for. Joshua calls the continuation each time ?person is successfully instantiated. This process repeats until no more matching predications are found in the database.

Repeating the query and asking why it succeeded confirms that the query pattern was successfully matched against facts in the database.

[ADULT ROSE] succeeded
[ADULT ROSE] was true in the database
[ADULT JOHN] succeeded
[ADULT JOHN] was true in the database
[ADULT FRED] succeeded
[ADULT FRED] was true in the database

This is a somewhat simplified explanation of the pattern matching process. For more detail: See the section "Pattern Matching in Joshua: Unification", page Pattern-Matching-in-Joshua-Unification-SECTION.

Logic variables can substitute for any argument in the argument list of a query pattern and can be combined with other, non-variable arguments such as constants.

Let’s add some statements about the health of people in our database, and ask various questions. For instance, to determine the current health status of the persons in our sample, the query can combine logic variables with a constant that specifies "today" as the time argument:
(tell [has-condition John measles last-year])
(tell [has-condition Rose pneumonia today])
(tell [has-condition Fred measles today])
(tell [has-condition Skip tendonitis today])
(tell [has-condition Fred cold last-week])
(tell [has-condition Marina sunburn today])
(tell [has-condition Fred infection last-month])
(tell [has-condition Dagwood hunger ?always]); note logic variable in database
 ; predication

[HAS-CONDITION MARINA SUNBURN TODAY]
[HAS-CONDITION SKIP TENDONITIS TODAY]
[HAS-CONDITION ROSE PNEUMONIA TODAY]
[HAS-CONDITION DAGWOOD HUNGER TODAY]; variable in database is instantiated to
 ; constant in query
[HAS-CONDITION FRED MEASLES TODAY]

While matching has-condition predcations, Joshua instantiates ?person to the first argument in the predication and ?condition to the second argument. Since we specified no logic variable for the third argument in the query pattern, Joshua matches it exactly. (Notice that since one database predication, namely [has-condition Dagwood hunger ?always] contains a logic variable, this variable became temporarily instantiated to the constant today in the query pattern.)

Here are some other examples.

[HAS-CONDITION FRED INFECTION LAST-MONTH]
[HAS-CONDITION FRED COLD LAST-WEEK]
[HAS-CONDITION FRED HEALTH YESTERDAY]
[HAS-CONDITION FRED MEASLES TODAY]

Logic variables have lexical scope and can be referred to by any Lisp code embedded in the joshua:ask continuation. The variables have their instantiated values until the continuation for a given answer has finished executing.

We can modify the continuation of our previous query about current health status to display a formatted answer using the logic variables as instantiated for each match:
(ask [has-condition ?person ?condition today]
    #'(lambda (ignore) (format t "-%A is ill with -A." ?person ?condition))
    :do-backward-rules nil)
MARINA is ill with SUNBURN.
SKIP is ill with TENDONITIS.
ROSE is ill with PNEUMONIA.
DAGWOOD is ill with HUNGER.
FRED is ill with MEASLES.

Adding some simple Lisp produces output with a heading:
(progn (format t "-% Today’s Patient Status:"
    (ask [has-condition ?person ?condition today]
        #'(lambda (&rest ignore)
            (format t "-%A is ill with -A." ?person ?condition))
            :do-backward-rules nil)

Today’s Patient Status:
MARINA is ill with SUNBURN.
SKIP is ill with TENDONITIS.
ROSE is ill with PNEUMONIA.
DAGWOOD is ill with HUNGER.
FRED is ill with MEASLES.
NIL

Giving a predication pattern to the convenience function joshua:map-over-database-predications is a very useful way of doing some operation on database predications that match the pattern. For example, suppose you are no longer interested in storing facts about persons with measles, so you want to find and remove these facts all at once. Give the appropriate pattern to joshua:map-over-database-predications. This function uses joshua:ask to search only the database, extracts the matching database predication from the joshua:ask continuation and executes the specified operation on that predication. In this case, we request that it joshua:untell every matching predication it finds.

(map-over-database-predications [has-condition ?person measles ?when] #'untell)

(cp:execute-command "Show Joshua Database")
True things
  [HAS-CONDITION DAGWOOD HUNGER ?ALWAYS]  [HAS-CONDITION FRED INFECTION LAST-MONTH]
  [HAS-CONDITION MARINA SUNBURN TODAY]  [HAS-CONDITION FRED COLD LAST-WEEK]
  [HAS-CONDITION SKIP TENDONITIS TODAY]  [HAS-CONDITION ROSE PNEUMONIA TODAY]
False things
  None

You can, of course, also use the Clear Joshua Database command with a predication pattern to remove matching predications from the database. To illustrate, suppose Fred has left this group of people we were monitoring, and all information about him needs to be removed from the database.
Clear Joshua Database (predications, All, or None [default All])
[has-condition Fred ?x ?when]
Predications being removed:
[HAS-CONDITION FRED INFECTION LAST-MONTH]
[HAS-CONDITION FRED COLD LAST-WEEK]
Untell the above predications? [default Yes]: Yes

Logic variables are also useful for narrowing down database displays to match only specified patterns. The default version of the Show Joshua Database command displays every item in the database.

Show Joshua Database (matching pattern [default All])

If instead of the default you give the Show Joshua Database command a specific pattern to look for, Joshua displays only predications matching that pattern. This lets you isolate those portions of the database that immediately interest you. Additionally, when you specify a pattern to match, you can limit the display to matching predications of only a single truth value, joshua:*true* or joshua:*false*. The default is both. To select one value, answer "No" to the question "(opposite truth-value too?)"

For example, we want to look at only those joshua:*true* predications that relate to hunger and ignore the rest of the database.

Show Joshua Database (matching pattern) [has-condition ?person hunger ?when]
(opposite truth-value too? [default Yes]: No

True things
[HAS-CONDITION DAGWOOD HUNGER ?ALWAYS]

Predication patterns as well as predications can be combined to express related ideas in combination. The next concept in our consideration of predications is that of logical connectives. See the section "Predications and Logical Connectives", page Predications-and-Logical-Connectives-SECTION.

### 4.4. Predications and Logical Connectives

Predicates: joshua::and
joshua::or
joshua::not

So far we have been considering predications in their simplest form: a single predic-  
ation made up of a predicate and its arguments. Such constructs have allowed us to express ideas in isolation, for example:

[has-condition Don-Quixote delusion today]

Knowledge is often more usefully expressed in some logical combination. For ex-  
ample, we might want to relate two ideas: that Don Quixote was suffering from delusions and that he believed windmills to be giants. Such combinations are made possible with the built-in Joshua predicates joshua::and, joshua::or, and
joshua::not. So, after defining the appropriate predicates, we can write a compound predication such as:

```
[and [has-condition Don-Quixote delusion today]
 [believes Don-Quixote windmills giants]]
```

Nesting predications inside each other is another way of building more complex predications out of single ones.

You can use logical connectives in joshuatell and joshuask statements (they are particularly useful, though, in rules and questions).

Compound joshuatell statements using joshua::and just save some of the labor of entering multiple predications into the database. Once entered, the predications don't remain yoked together, that is, each of the component predications is inserted individually. You cannot use or with joshuatell statements in the current release (joshuatelling a disjunction is a statement of partial knowledge).

joshuask queries can define more complex goals by using logical connectives.

To write compound joshuatell and joshuask forms, wrap the predications with the appropriate connective. For example:

```
(tell [and [has-condition Don-Quixote delusion today]
 [believes Don-Quixote windmills giants]]
 [AND [HAS-CONDITION DON-QUIXOTE DELUSION TODAY]
 [BELIEVES DON-QUIXOTE WINDMILLS GIANTS]]
```

```
(ask [or [has-condition ?person measles now]
 [has-condition ?person pneumonia now]] '#print-query)
 [HAS-CONDITION FRED MEASLES NOW]
 [HAS-CONDITION ROSE PNEUMONIA NOW]
```

Note that when you joshuatell compound forms, joshuatell does not return a boolean value.

For the next example, we build part of a library database of authors, book titles, and author information. We'll embed several Joshua operations inside a Lisp function that clears the database, enters the compound joshuatell predications, and displays the newly created database. To display the database within Lisp code, we use the function cp:execute-command "Show Joshua Database".

First, the predicate definitions.

```
(define-predicate author-of (work author))
(define-predicate lived (object when))
(define-predicate profession-of (person profession))
```

Next we build up part of the library.
(defun library ()
  (clear)
  (tell [and [author-of Tempest Shakespeare]
            [lived Shakespeare 16th-century]
            [profession-of Shakespeare actor]
            [author-of "Art of Love" Ovid]
            [lived Ovid 1st-century-BC]
            [profession-of Ovid poet]
            [author-of poems Henry-VIII]
            [lived Henry-VIII 16th-century]
            [profession-of Henry-VIII king]
            [author-of "Art of Love" Capellanus]
            [lived Capellanus 12th-century]
            [profession-of Capellanus chaplain]
            [author-of "Art of Love" Fromm]
            [lived Fromm 20th-century]
            [profession-of Fromm sociologist]
            [author-of poems Ralegh]
            [lived Ralegh 16th-century]
            [profession-of Ralegh soldier]])
  (cp:execute-command "Show Joshua Database"))

LIBRARY

(library)

True things

[PROFESSION-OF RALEGH SOLDIER]  [LIVED HENRY-VIII 16TH-CENTURY]
[PROFESSION-OF FROMM SOCIOLOGIST]  [LIVED OVID 1ST-CENTURY-BC]
[PROFESSION-OF CAPELLANUS CHAPLAIN]  [LIVED SHAKESPEARE 16TH-CENTURY]
[PROFESSION-OF HENRY-VIII KING]  [AUTHOR-OF POEMS RALEGH]
[PROFESSION-OF OVID POET]  [AUTHOR-OF POEMS HENRY-VIII]
[PROFESSION-OF SHAKESPEARE ACTOR]  [AUTHOR-OF "Art of Love" FROMM]
[LIVED RALEGH 16TH-CENTURY]  [AUTHOR-OF "Art of Love" CAPELLANUS]
[LIVED FROMM 20TH-CENTURY]  [AUTHOR-OF "Art of Love" OVID]
[LIVED CAPELLANUS 12TH-CENTURY]  [AUTHOR-OF TEMPEST SHAKESPEARE]

False things

None

Now we use a compound query to joshua:ask who wrote the book "Art of Love," when the author lived, and what the author's profession was; the continuation formats the answers into discursive English.

(ask [and [author-of "Art of Love" ?author]
      [lived ?author ?century]
      [profession-of ?author ?profession]]
  #'(lambda (ignore) (format t """"%-%-%-%A, a %-A %-A, wrote a version

FROMM, a 20TH-CENTURY SOCIOLOGIST, wrote a version of The Art of Love.

CAPELLANUS, a 12TH-CENTURY CHAPLAIN, wrote a version of The Art of Love.

OVID, a 1ST-CENTURY–BC POET, wrote a version of The Art of Love.

Logical connectives focus and refine queries and cut down the system’s search time. In a compound joshua:ask using joshua:and, the logic variables that are common to all component predications must be instantiated to a common object for the query to succeed.

In our example, the logic variable ?author is instantiated to a name matching the first query pattern, [author–of "Art of Love" ?author], namely Fromm.

?author remains instantiated to Fromm and the search for predications to match the rest of the queries above ([lived ?author ?century] and [profession–of ?author ?profession]), is narrowed down to only those predications matching [lived Fromm ?century] and [profession–of Fromm ?profession]. Only when all the predications in the compound query have been satisfied is the continuation called.

After the continuation executes, ?author, ?century, and ?profession become uninstantiated in the first two predication patterns. The last predication pattern still has ?author instantiated. Joshua looks for the next solution to [profession–of Fromm ?profession]. It finds none, so it backs up and tries solutions to [lived Fromm ?century]. It finds none of these either, so it backs up and looks for solutions to the first query pattern, [author–of "Art of Love" ?author]. If we had joshua:asked the three questions separately,

(ask [author–of "Art of Love" ?author] ...)
(ask [lived ?author ?century] ...)

the system would have produced three separate lists containing a great deal of extraneous information, especially if we had a very large database: we would have gotten a list of authors of "The Art of Love"; a list of every author in the database and his century; and a list of every author in the database and his profession.

The logical connective joshua:or instantiates its logic variables separately for each predication in a query. The effect is the same as asking separate queries.

Logical connectives can, of course, be combined, as in this example. Here we look for all persons who are not children, and who are currently ill with either measles or mumps. (Assume the appropriate joshua:tell statements have been entered.)

(define-predicate child (person))
(ask [and [not [child ?person]]
  [or [has-condition ?person measles today]
   [has-condition ?person mumps today]]]
  #'print-query :do-backward-rules nil)
[AND [NOT [CHILD HELEN]] [OR [HAS-CONDITION HELEN MEASLES TODAY]
  [HAS-CONDITION HELEN MUMPS TODAY]]]

Now that we've seen how to build predications and predication patterns, it would be useful to display the meaning of predications in natural language. See the section "Formatting Predications: the SAY Method", page Formatting-Predications-the-SAY-Method-SECTION.

4.5. Formatting Predications: the SAY Method

Function: joshua:say
Formatt directive: "\SAY\"
Macro: joshua:define-predicate-method

To print the meaning of a predication in natural language (or some other alternative, such as graphics), as opposed to the predicate calculus notation in which programs are written, you can use the format function, as we have been doing in some of our joshua:ask continuations; or we can embed formatting directives into a special method available for that purpose, the joshua:say method. joshua:say is actually a multi-purpose "hook" for using natural language, or graphics, or any other predicate-dependent approach you wish. Judicious use of joshua:say methods can make it easier to generate user interfaces quickly.

You define a joshua:say method by using the general method definition function, joshua:define-predicate-method. (This is the same function you would use to modify other parts of the Joshua protocol, and its advanced use is covered in the Joshua Reference Manual.)

The arguments to joshua:define-predicate-method are:

• The function-spec of the protocol method you are defining, for example, (say foo).
• The argument list. In the case of joshua:say methods, this should almost always be (&optional (stream *standard-output*)).
• The body of the method.

The body of the method contains any Lisp code, presumably doing output to the stream mentioned in the argument list.

For example, suppose you want to write a joshua:say method for the predicate [favorite-meal ...]. This predicate takes two arguments, namely, eater and food.
You’ll want to bind Lisp variables to the arguments of the predication so that you can use the arguments within the predicate method. To do this, wrap the macro `joshua:with-statement-destructured` around the body. The predicate arguments then become lexically available, and can be referred to by the Lisp forms. (If you have already defined the predicate with its instance variables destructured, you don’t need to use `joshua:with-statement-destructured`. For more on methods of making predicate arguments lexically available: See the macro `joshua:define-predicate`, page `joshua:define-predicate-FUNCTION`.)

A `joshua:say` method for the predicate `[favorite-meal ...]` might look something like this:

```
(define-predicate favorite-meal (eater food))

(define-predicate-method (say favorite-meal) (&optional (stream *standard-output*))
  (with-statement-destructured (eater food) self
    (format stream
      "- & You can please "A by giving them "A to eat" eater food)))
```

Having defined this `joshua:say` method, you can use it in a variety of ways, as shown here. You can call the method directly, giving it an instantiated predication argument:

```
(say [favorite-meal monkeys bananas])
  You can please MONKEYS by giving them BANANAS to eat
NIL
```

Or you can use a `joshua:format` directive included in Joshua to call `joshua:say` instead of `joshua:prin1` or `joshua:princ`:

```
(format t "-% Is it true that: ~\"\" [favorite-meal monkeys bananas]
  Is it true that:
    You can please MONKEYS by giving them BANANAS to eat?
NIL
```

Joshua provides the convenience function, `joshua:say-query` to display a satisfied query using either the default `joshua:say` method, or your own `joshua:say` method, if any. (The default `joshua:say` method simply prints out the answer to the query.)

Let’s fatten up our database of foods and eaters before using the `joshua:say` method.
(defun eat-it ()
  (clear)
  (tell [and [favorite-meal bears honey]
              [favorite-meal mosquitoes people]
              [favorite-meal spiders flies]
              [favorite-meal Joshuas predications]
              [favorite-meal monkeys bananas]])
  (cp:execute-command "Show Joshua Database"))

(eat-it)
True things
[FAVORITE-MEAL MONKEYS BANANAS] [FAVORITE-MEAL MOSQUITOES PEOPLE]
[FAVORITE-MEAL JOSHUAS PREDICATIONS] [FAVORITE-MEAL BEARS HONEY]
[FAVORITE-MEAL SPIDERS FLIES]
False things
None

You can please MONKEYS by giving them BANANAS to eat
You can please JOSHUAS by giving them PREDICATIONS to eat
You can please SPIDERS by giving them FLIES to eat
You can please MOSQUITOES by giving them PEOPLE to eat
You can please BEARS by giving them HONEY to eat

Note: Technically our joshua:ask method is correctly defined and it certainly works. But we purposely misphrased it, in order to make a point: the phrasing makes an implicit, new (and unwarranted) connection between the idea "favorite-meal" and the idea "to please by giving to eat". Much confusion in programming comes from such casual redefinition. A better phrasing would have been something like "The favorite meal of ... is ...".

4.6. Tracing Predications

Commands:    Enable Joshua Tracing
             Disable Joshua Tracing

Because database operations like joshuatell and joshua:ask are so fundamental to the operation of Joshua programs, we often want to watch as things are stored into or retrieved from the database. You can trace the basic operations on predications by using the command:

Enable Joshua Tracing (type of tracing) Predications

This turns on the tracing of joshua:ask and joshuatell. Each time your program calls joshua:ask or joshuatell, the tracing facility prints a message saying which operation is being done and on what predication.

Example:
(defun tell-and-ask-foods ()
  (clear)
  (tell [and [favorite-meal bears honey]
            [favorite-meal mosquitoes people]
            [favorite-meal spiders flies]
            [favorite-meal Joshuas predcations]
            [favorite-meal monkeys bananas]]

Command: Enable Joshua Tracing (Type of tracing) Predications

Predication tracing is on
  Tracing All predicates
  Traced events: Ask and Tell

Command: (tell-and-ask-foods)

► Telling predication [AND [FAVORITE-MEAL BEARS HONEY]
                      [FAVORITE-MEAL MOSQUITOES PEOPLE]
                      [FAVORITE-MEAL SPIDERS FLIES]
                      [FAVORITE-MEAL JOSHUAS PREDICATIONS]
                      [FAVORITE-MEAL MONKEYS BANANAS]]
► Telling predication [FAVORITE-MEAL BEARS HONEY]
► Telling predication [FAVORITE-MEAL MOSQUITOES PEOPLE]
► Telling predication [FAVORITE-MEAL SPIDERS FLIES]
► Telling predication [FAVORITE-MEAL JOSHUAS PREDICATIONS]
► Telling predication [FAVORITE-MEAL MONKEYS BANANAS]
► Asking predication [FAVORITE-MEAL =EATER =FOOD]
  You can please MONKEYS by giving them BANANAS to eat
  You can please JOSHUAS by giving them PREDICATIONS to eat
  You can please SPIDERS by giving them FLIES to eat
  You can please MOSQUITOES by giving them PEOPLE to eat
  You can please BEARS by giving them HONEY to eat

In order to disable the tracing of predications use the command:

  Disable Joshua Tracing (type of tracing) Predications

This turns off all tracing of predications.

You can get a greater degree of control over tracing by using the menu option to the Enable Joshua Tracing command:

  Enable Joshua Tracing (type of tracing) Predications :Menu Yes

This shows a menu of all of the tracing options available for predications, letting you trace only predications matching a particular pattern, only predications using a particular predicate (or type of predicate), and to specify at which events you would like to see trace information.

In addition to the above commands you can set some tracing options by mousing right on a predication. This gives you a menu of the available options.
4.7. Miscellaneous Predication Facilities

Several predication facilities are available for use within Lisp code. We have already used one, namely, *joshua:with-statement-destructured*. Here are some others. We suggest you look up their dictionary entries for more detail.

**joshua:different-objects**  If the arguments are *joshua:eql*, or if either argument is an uninstantiated logic variable, *joshua:different-objects* returns **nil**. Otherwise it returns **t**.

**joshua:make-predication**  Constructs a predication (does not enter it into the database). The [ ] syntax is a reader macro that expands into this.

**joshua:predication**  The base flavor for all predications in Joshua. This works well with functions like *joshua::typep* and *joshua::typecase*.

**joshua:predicationp**  Checks whether an object is built on *joshua:predication*.

**joshua:with-statement-destructured**

Lets you bind Lisp variables to predication arguments.

This chapter has shown you how to generate information in Joshua by building predications, storing them in a database, and searching the database to answer queries. The next major concept is using Joshua to derive information by reasoning about predications with rules and questions. See the section "Rules and Inference", page Rules-and-Inference-SECTION.
5. Joshua Rules and Inference

Macro: \texttt{joshua:defrule}
Function: \texttt{joshua:undefrue}
Command: Show Joshua Rules
Zmacs command \texttt{\textasciitilde } Kill definition
Concepts: Forward Chaining (\texttt{joshuatell})
Backward Chaining (\texttt{joshua:ask})

Often the information one looks for by searching the database is not there explicitly, but must be inferred by reasoning about the knowledge that is already there.

While predications define the relationships between objects and supply us with information about the domain, rules define the \textit{reasoning} that is performed about predications, and control how we deduce knowledge from existing knowledge.

A rule is an independent piece of declarative and procedural information that determines how Joshua responds to a specified set of circumstances. Unlike conventional programming constructs, rules automatically execute in the proper order whenever the appropriate circumstances occur, rather than executing when a program’s control structure reaches a specific portion of the code. In other words, rules can execute at any time, regardless of the order in which they are written.

The reasoning done by rules involves either \textit{forward} or \textit{backward} chaining. Joshua programs can use either or both of these rule types.

Forward chaining is \textit{data-directed inference}, that is, reasoning from \textit{known} facts to some conclusion. This form of reasoning says: "I now know fact X. What can I conclude from this?" For example, given the facts that birds can fly and that Tweety is a bird, a forward chaining rule can deduce the new fact that Tweety can fly. Thus, forward chaining is instrumental in adding to the database.

Forward chaining is activated by \texttt{joshuatell}. That is, whenever you \texttt{joshuatell} Joshua a new predication, the system looks for forward chaining rules that it can use (combined with knowledge already in the database) to draw conclusions from the new knowledge you gave it.

Backward chaining is \textit{goal-directed inference}, that is, reasoning to satisfy some desired conclusion. This form of reasoning says: "I want to know fact X. How do I determine its validity?" For example, given the goal of determining whether Tweety can fly, a backward chaining rule would look for the facts it needs in order to support this goal, asking, for example, whether birds can fly and whether Tweety is a bird.

Backward chaining is activated by an \texttt{joshua:ask}. The \textit{query pattern} becomes the \textit{goal}, and the system then looks for the facts and rules and questions that might substantiate this goal. Backward chaining is thus useful in helping you determine the validity of a conclusion or goal.
Forward and backward rules look the same, except for the keyword :forward or :backward that indicates the rule’s control structure (its inference method). If the rule doesn’t use Lisp code, it should work equally well in either direction. Either method of reasoning accomplishes the same result; the choice of inference method depends on the problem being solved. Some problems are much more efficiently approached with one control structure than with the other.

This chapter summarizes basic information about forward and backward rules in Joshua.

---

**Advanced Concepts Note:**

The Joshua protocol has functions that determine how rules are stored, deleted, and looked up. See the section “The Joshua Rule Indexing Protocol” in *Joshua Reference Manual*. If you provide a consistent alternate implementation of these generic functions, you can customize rule management for your application. See the section “Customizing the Rule Index” in *Joshua Reference Manual*.

---

**5.1. Defining Joshua Rules**

A rule is defined with joshua:defrule. A rule has a name, a keyword argument specifying its control structure (whether it is a forward or a backward chaining rule), and a combination of patterns divided into an if-part and a then-part.

The control structure argument :importance lets you prioritize rules (the higher the value you give to :importance, the higher the priority). Higher priority rules run before lower priority rules.

Another control structure argument, :documentation, lets you add a documentation string explaining what the rule does. Use the Lisp function documentation to retrieve the rule’s documentation string.

For example, this forward chaining rule describes some facts that let you deduce an identity for an unknown creature.

```lisp
(defrule dragon-id-kit (:forward :documentation "Identifies dragons")
  if [and [huge ?creature]
       [breathe ?creature fire]
       [or [guard ?creature gold]
            [guard ?creature maiden]]]
  then [dragon ?creature])

(documentation 'dragon-id-kit)
"Identifies dragons"
```

This backward chaining rule describes how to compute the grandfather relationship along paternal lines.
(defrule paternal-grandfather (:backward)
  if [and [father ?person ?dad]
          [father ?dad ?gramps]]
  then [grandfather ?person ?gramps])

Except for their control structure, both rules look similar.

A rule’s if-then parts are its logical structure. The if-part of a rule logically describes conditions under which the rule is applicable. The then-part of a rule logically describes the rule’s conclusions.

The if- and then- clauses in a rule can occur in any order. That is, both

if [...] then [...] and

then [...] if [...] are valid.

While the form of the if-then parts is identical for forward and backward rules, procedurally the parts differ for each rule type. This is because the logical structure maps into an imperative structure that reflects the inferencing method being used (forward or backward chaining). Forward and backward chaining rules differ in their mapping of logical to imperative parts.

A rule’s imperative parts consist of a trigger part and an action part. The trigger part determines if a rule is applicable to a given situation. The action part determines what operations the rule performs when it executes. Forward and backward rules have different trigger and action parts, so we discuss each rule type separately.

The command Show Joshua Rules displays the currently defined rules. Various options let you tailor the display. Please consult the dictionary entry for this command.

5.1.1. How Forward Rules Work

In a forward rule, the if-part is the trigger, and the then-part is the action. See figure dragon.

![Diagram of Forward Rule Trigger and Action Parts](image-url)

Figure 3. Forward Rule Trigger and Action Parts
In a forward rule the trigger is a single (possibly a compound) predication pattern stating conditions that must be satisfied for the rule to fire. The trigger can also contain Lisp code which is discussed below.

Data-directed inference is triggered by the addition of new facts into the database with joshuatell. That is, new facts cause Joshua to look for forward rule triggers that can be satisfied by these facts.

What are new facts? A fact is new only when you joshuatell it for the first time. A fact that you joshuatell more than once is no longer new knowledge. Similarly, a fact that you joshua:unjustify and then joshuatell again is not new knowledge, since it was never removed from the database. In contrast, a fact that you remove from the database (joshua:untell) and then joshuatell again does represent new knowledge for Joshua.

There is no special order in which forward rule patterns are triggered. This is determined entirely by the order in which the new facts are entered.

To satisfy part of a rule’s trigger, the pattern of that trigger must match that of the newly added database predication. This matching, called unification in Joshua, happens either because the two patterns are already equivalent, or because Joshua succeeded in finding substitutions for logic variables so that the two patterns become equivalent. To understand the basics of forward rule operation we can postpone looking at the details of unification. We cover this topic in the section “Pattern Matching in Joshua: Unification” along with the scoping rules that determine when two logic variables with the same name are the same and when they differ.

If a rule’s trigger parts are joined by and, all conditions must be satisfied for the rule to trigger. If the trigger parts are joined by or, satisfying any of the conditions triggers the rule.

Note that backward rules are not automatically invoked while trying to satisfy a forward trigger. Since you can use Lisp code in forward triggers, your forward rules can explicitly call (ask [foo ?x] ...) in order to use a backward rule to satisfy foo.

When all of a forward rule’s conditions (if-parts) are satisfied, the rule is fully triggered; it then fires and executes the action(s) in its action (then) parts. The action part can specify any action, including the new facts to be deduced directly from the current facts. If the action involves deducing a new fact, that fact is automatically added to the database (that is, a joshuatell is implicit for forward rule deductions).

When the firing of forward chaining rules results in the addition of new facts to the database, this in turn can cause more forward rules to be triggered and fired, generating chains of conclusions until no more new facts can be generated.

Can you define a forward rule after adding a new fact that will trigger the rule into the database? The answer is that you can enter new rules and new facts in any order, because Joshua ensures that the right forward rules always get triggered.

Here’s an example of forward chaining in Joshua. Rule danger-sign says that if a person is known to be a smoker and to have hypertension, then we can deduce
that this person is in a high-risk category. We also want the system to notify us if it makes such a deduction so that we can take some appropriate action.

(define-predicate smoker (person))
(define-predicate has-condition (person condition when))
(define-predicate at-risk (person))

(defrule danger-sign (:forward)
  if [and [smoker ?person]
           [has-condition ?person hypertension today]]
  then [at-risk ?person]
       (format t "Suggest to ~S that smoking is dangerous to hypertensive persons" ?person))

Our database already contains the predication [smoker Ashley]. When that predication was first added, it satisfied the first if-part of rule danger-sign, namely [smoker ?person]. Now we joshuatell the system that Ashley is suffering from hypertension. This satisfies the rule’s second if-part and triggers the rule. The rule fires and executes its then-part (the action part). This causes it to display a message and to add the newly deduced predication, [at-risk Ashley], to the database.

(tell [has-condition Ashley hypertension today])
Suggest to Ashley that smoking is dangerous to hypertensive persons
[HAS-CONDITION ASHLEY HYPERTENSION TODAY]

Show Joshua Database (matching pattern) [at-risk ?person]
True things
[AT-RISK ASHLEY]
False things
None

Assume we have another forward rule, preventive-care, as follows:

(define-predicate needs-checkup (person when))

(defrule preventive-care (:forward)
  if [at-risk ?person]
  then [needs-checkup ?person monthly])

The addition of the new fact, [at-risk Ashley], to the database now satisfies the rule trigger of the above rule, causing it to fire in its turn, and to generate yet another new fact, namely, [needs-checkup Ashley monthly].

Show Joshua Database (matching pattern) [needs-checkup ?person ?frequency]
True things
[NEEDS-CHECKUP ASHLEY MONTHLY]
False things
None

This chain of inferences continues as long as there are forward rules that can be fully triggered by newly generated facts.
Forward chaining rules can contain Lisp code in both their if (trigger) and then-(action) parts. Such Lisp code can refer to any logic variables that appear inside the body of the rule.

Our earlier rule, danger-sign, for example, uses Lisp in its action part. Lisp code in the then-part of a forward rule is just put into the rule body, to be run when the rule fires.

Here is an example using Lisp code in a forward rule trigger. We refer to such code as a procedural node. Procedural nodes can introduce new variables, have side-effects, call joshua:succeed, and so on.) Lisp code in the if-part of a forward rule returns non-nil if it wants the match to continue.

```
(defrule acceptable-price-rule (:forward)
    if [and [price-ceiling ?available-cash]
        [todays-price ?cost]
        (≤ ?cost ?available-cash)]
    then [acceptable-price ?cost])
```

You can watch forward rule execution by enabling Joshua tracing. See the section "Tracing Rules", page Tracing-Rules-SECTION.

### 5.1.2. How Backward Rules Work

Since backward rules are goal-directed inference, it is the rule’s then-part that specifies the desired goal. Thus, backward rule inferencing is triggered by the rule’s then-part, while the if-part is the action part. See figure gramps.

Note that the order of the if-then clauses does not matter. Those who like to place the then-part of backward rules first, so that the trigger always comes first, can safely do so.

```
(defrule paternal-grandfather (:backward)
    if [and [father =person =dad]
        [father =dad =gramps]]
    then [grandfather =person =gramps])
```

Figure 4. Backward Rule Trigger and Action Parts

Currently, backward rule inferencing is triggered by a single goal, that is, a backward trigger must be a single predication pattern. This restriction may be removed in the future.

A backward rule can contain Lisp code in its action (if-) part, but not in its trigger (then-) part. For example:

```
(defrule good-condition-rule (:backward)
    if (not (eql ?condition 'rusted))
    then [good-condition ?condition])
```
In the section "Querying the Database", we introduced \texttt{joshua:ask} for the limited purpose of looking up items in the database. Typically, the main purpose of \texttt{joshua:ask} is to find solutions through reasoning (finding backward rules and questions to run) as well as through database lookup.

To satisfy a query, \texttt{joshua:ask} always looks in the database first, trying to match the query pattern with a database predication. Next, \texttt{joshua:ask} by default searches for backward chaining rules whose trigger pattern (\textit{then}-part) matches the query pattern. When such a match occurs, the backward rule is triggered. In Joshua, this pattern matching is called \textit{unification}. For more on the mechanics of unification and the scoping of variables: See the section "Pattern Matching in Joshua: Unification", page Pattern-Matching-in-Joshua:-Unification-SECTION.

Once a backward rule is triggered, it processes its \textit{if}-parts (the action parts) from the top in the order given, searching for facts to validate the desired conclusion. (If the \textit{if}-part is a compound predication joined by \texttt{and}, each component becomes a subgoal that is the subject of an implicit \texttt{joshua:ask}.)

The rule succeeds when all its \textit{if}-parts (subgoals) have been satisfied. Satisfied means that there is a fact in the database that can serve as the bottom of the support structure for each subgoal.

\texttt{joshua:ask} collects information describing the solution process; when it finds a solution, \texttt{joshua:ask} calls its continuation, passing it the answer, along with the information collected about it. The continuation then executes on this result. We’ll look at some typical continuation requests shortly.

(For the basics of \texttt{joshua:ask}: See the section "Querying the Database", page Querying-the-Database-SECTION.)

The information resulting from the success of a backward chaining rule is not automatically added to the database, unless you explicitly do this in your program (in the continuation of the \texttt{joshua:ask}, for example).

Note that for special cases where you don’t want \texttt{joshua:ask} to use backward rules, you can disable this feature by specifying \texttt{:do-backward-rules nil} to \texttt{joshua:ask}.

Here’s a backward chaining example. Assuming your goal is to find a treatment for patients with a given condition, you might formulate a backward chaining rule something like this:

```
(define-predicate has-condition (person condition when))
(define-predicate effective-treatment (drug condition))
(define-predicate allergic (person drug))
(define-predicate appropriate-treatment (drug person condition))

(defrule find-cure (:backward)
  if [and [has-condition ?person ?condition today]
        [effective-treatment ?drug ?condition]
        [not [allergic ?person ?drug]]]
  then [appropriate-treatment ?drug ?person ?condition])
```

If you then joshua:ask what drug is appropriate for a certain condition, in a certain patient, Joshua first searches the database, after which the query triggers rule find-cure by matching its then-part. When the rule fires, Joshua works backwards from this goal to satisfy each of the component parts of the rule.

- Checks that the patient you specified has that condition (or, if you gave a variable for the patient, it iterates over all patients with that illness);
- Tries to find an effective drug for that condition; if successful, checks whether the patient is allergic to this drug.
- If the patient is allergic to this drug, or if Joshua simply doesn’t know, it discards the drug and repeats the process of searching for another effective drug and testing for patient allergy, until it finds a drug to which the patient is not allergic, or fails trying.

For example, assume the database contains the following predications:

```
(tell [and [has-condition primadonna sore-throat today]
  [effective-treatment gargle-with-ammonia sore-throat]
  [not [allergic primadonna gargle-with-ammonia]])
```

When you ask the system for a treatment for sore throat for all patients with that affliction, the backward rule find-cure triggers, satisfies its subgoals, and executes the continuation which in this case uses the convenience function joshua:print-query to print out the answer.

```
[APPROPRIATE-TREATMENT GARGLE-WITH-AMMONIA PRIMADONNA SORE-THROAT]
```

5.1.2.1. Explaining Backward Chaining Support

As mentioned earlier, joshua:ask passes its continuation a list containing the answer it found (or derived) together with information tracing the reasoning process that was followed (for example, whether the answer was found in the database or came from backward rules, and if so, which one(s), and so on).

Joshua provides two convenience functions that extract and interpret the support information for you. We’ve introduced these earlier. (See the section “Querying the Database”, page Querying-the-Database-SECTION.) We mention them again here, since their usefulness becomes far more apparent in a rule context.

joshua:print-query-results extracts and displays the successful query and tells you why it succeeded.

This, for example, is the reasoning Joshua went through to answer our query about the right treatment for sore throat.
[APPROPRIATE-TREATMENT GARGLE-WITH-AMMONIA PRIMADONNA SORE-THROAT] succeeded
It was derived from rule FIND-CURE
[HAS-CONDITION PRIMADONNA SORE-THROAT TODAY] succeeded
[HAS-CONDITION PRIMADONNA SORE-THROAT TODAY] was true in the database
[EFFECTIVE-TREATMENT GARGLE-WITH-AMMONIA SORE-THROAT] succeeded
[EFFECTIVE-TREATMENT GARGLE-WITH-AMMONIA SORE-THROAT] was true in the database
[not [ALLERGIC PRIMADONNA GARGLE-WITH-AMMONIA]] succeeded
[ALLERGIC PRIMADONNA GARGLE-WITH-AMMONIA] was false in the database

This tells us that the query pattern triggered rule find-cure, and that each of this rule’s three subgoals was satisfied from a database lookup. Since backward chaining stops when it reaches database predications, that is the end of the support information; there is no attempt to trace why those database predications are valid.

If more than one backward rule had been invoked to find the answer, joshua:print-query-results would have interpreted that information as well, tracing the support through each rule to the predications used to satisfy parts of the rule.

The convenience function joshua:graph-query-results gives you the same information as joshua:print-query-results but in graph form.

Figure graph-rule shows what this graph looks like for our previous query.

The top of the graph shows the satisfied query and names the rule that satisfied it. Ovals denote queries that were not satisfied in the database. Rectangles denote queries (in this case subgoals) that were satisfied from the database. The arrows point from the support to the object being supported.

In our example, the graph shows the database predications that satisfied the rule’s subgoals. Note that the Database heading indicates when the truth value for a given predication is joshua:*false*.

\[
\text{ASK} \quad \text{[appropriate-treatment?drug?person sore-throat] } \#’\text{print-query-results} \\
\text{[APPROPRIATE-TREATMENT GARGLE-WITH-AMMONIA PRIMADONNA SORE-THROAT]} \text{ succeeded} \\
\text{It was derived from rule FIND-CURE} \\
\text{[HAS-CONDITION PRIMADONNA SORE-THROAT TODAY]} \text{ succeeded} \\
\text{[HAS-CONDITION PRIMADONNA SORE-THROAT TODAY]} \text{ was true in the database} \\
\text{[EFFECTIVE-TREATMENT GARGLE-WITH-AMMONIA SORE-THROAT]} \text{ succeeded} \\
\text{[EFFECTIVE-TREATMENT GARGLE-WITH-AMMONIA SORE-THROAT]} \text{ was true in the database} \\
\text{[not [ALLERGIC PRIMADONNA GARGLE-WITH-AMMONIA]] succeeded} \\
\text{[ALLERGIC PRIMADONNA GARGLE-WITH-AMMONIA] was false in the database} \\
\]

\[
\text{Figure 5. Graphing query support from backward rule} \\
\]

If you prefer to extract and interpret the support information yourself rather than have it interpreted for you, Joshua provides accessor functions to let you do this. Consult the dictionary entry for joshua:ask to see how these accessor functions work.

(If you have a TMS, you can get an explanation of forward support, also. That is, with a TMS Joshua can tell you the reasons why a database predication is in the database. The function joshua:explain displays the support in a manner similar to
the output of \texttt{joshua:print-query-results}; \texttt{joshua:graph-tms-support} graphs the explanation in a manner similar to \texttt{joshua:graph-query-results}.

You can trace backward rule operation by using the tracing facility. See the section “Tracing Rules”, page Tracing-Rules-SECTION.

\section*{5.2. Removing Joshua Rule Definitions}

You can remove rule definitions from the system either individually or collectively.

The function \texttt{joshua:undefrule} removes a single rule definition.

For example:

\begin{verbatim}
(undefrule 'danger-sign)
\end{verbatim}

You can do this same operation from a Zmacs editor buffer with the extended Zmacs command $\texttt{K}$ Kill Definition. For a sample interaction with this command: See the macro \texttt{joshua:undefine-predicate}, page \texttt{joshua:undefine-predicate-FUNCTION}. Note that any predications previously deduced by a rule still remain in the database after you remove the rule's definition.

To remove all rule definitions at once, use the function \texttt{joshuaclear}. We have been using this function to clear the database, but it has a second optional argument to let you clear all rule definitions.

The full argument list of \texttt{joshuaclear} is:

\begin{verbatim}
CLEAR: (&OPTIONAL (CLEAR-DATABASE T) UNDEFRULE-RULES)
\end{verbatim}

Thus, typing (clear t t) clears the database and at the same time removes all rule definitions. Be aware that this is a rather drastic step, as applications depending on these rules will no longer work, until you reload these rules.

The command Clear Joshua Database has an option to let you clear all Joshua rules. The cautions just mentioned apply here as well.

\section*{5.3. Tracing Rules}

Commands: Enable Joshua Tracing

Disable Joshua Tracing

While debugging Joshua programs you often want to be able to watch the rules execute. You can trace the behavior of rules by using the command

\begin{verbatim}
Enable Joshua tracing (type of tracing) Forward Rules
\end{verbatim}

for forward rules, or the command

\begin{verbatim}
Enable Joshua Tracing (type of tracing) Backward Rules
\end{verbatim}

for backward rules.

This causes a trace message to be printed every time a rule runs. With forward rules the message appears each time the \texttt{if} part of the rule (the trigger) is suc-
cessfully completed, and the then-part (the action) is about to be executed. Here's a simple example of tracing forward rules.

Example:

(define-predicate higher-in-food-chain (eater lower-in-food-chain))
(define-predicate favorite-meal (eater food))

(defrule basic-food-chain (:forward)
  if [favorite-meal ?eater ?eatee]
  then [higher-in-food-chain ?eater ?eatee])

(defrule transitive-food-chain (:forward)
  if [and [favorite-meal ?eater ?eatee]
       [higher-in-food-chain ?eatee ?food]]
  then [higher-in-food-chain ?eater ?food])

(defun meals ()
  (clear)
  (tell [and [favorite-meal red-herring worm]
           [favorite-meal worm algae]])
  (tell [favorite-meal Miss-Marple red-herring])
  (cp:execute-command "Show Joshua Database"))

Command: Enable Joshua Tracing (Type of tracing) Forward Rules

Forward Chaining tracing is on
Tracing All forward rules
Traced events: Fire and Queue

Command: (meals)
  ▶ Firing forward rule BASIC-FOOD-CHAIN (1 trigger)
  ▶ Firing forward rule BASIC-FOOD-CHAIN (1 trigger)
  ▶ Firing forward rule TRANSITIVE-FOOD-CHAIN (2 triggers)
  ▶ Firing forward rule TRANSITIVE-FOOD-CHAIN (2 triggers)
  ▶ Firing forward rule TRANSITIVE-FOOD-CHAIN (2 triggers)
  ▶ Firing forward rule BASIC-FOOD-CHAIN (1 trigger)

True things
[HIGHER-IN-FOOD-CHAIN MISS-MARPLE RED-HERRING]
[HIGHER-IN-FOOD-CHAIN MISS-MARPLE WORM]
[HIGHER-IN-FOOD-CHAIN MISS-MARPLE ALGAE]
[HIGHER-IN-FOOD-CHAIN WORM ALGAE]
[HIGHER-IN-FOOD-CHAIN RED-HERRING ALGAE]
[HIGHER-IN-FOOD-CHAIN RED-HERRING WORM]
[FAVORITE-MEAL MISS-MARPLE RED-HERRING]
[FAVORITE-MEAL WORM ALGAE]
[FAVORITE-MEAL RED-HERRING WORM]

False things
None

Notice that as one rule firing joshua:tells a predication that causes another rule
to fire, the tracing facility indents another level. This shows you the dependency between the rules. Also, various items in the trace display are mouse sensitive and can provide more information about the program execution.

Tracing backward rules is a little more complicated, as they can be used to generate multiple solutions to a query. There are four events associated with the running of a backward rule.

First, when we try to match the trigger of the rule (the *then*-part of a backward rule), the trace message says that we are **Trying** the rule. When we successfully complete the rule action (the goals in the *if*-part), the message says that we are **Succeeding** from the rule. As we try to find another way to satisfy the rule, the message says that we are **Retrying** the rule. And lastly, when there are no more solutions for the rule, the trace message says that the rule is **Failing**.

The backward rule tracing facility uses this terminology to let you follow the execution of the rules as Joshua tries to satisfy a query. In order to demonstrate the tracing of backward rules we will first enable the tracing of both predications and backward rules. Tracing predications along with rules shows you the *joshua:asks* or *joshuatells* corresponding to each rule firing. Here’s a simple example of backward rule tracing.

**Example:**

```
(define-predicate has-condition (person condition when))
(define-predicate effective-treatment (drug condition))
(define-predicate allergic (person drug))
(define-predicate appropriate-treatment (drug person condition))

(defun find-cure (:backward)
  (if [and [has-condition ?person ?condition today]
          [effective-treatment ?drug ?condition]
          [not [allergic ?person ?drug]]
        then [appropriate-treatment ?drug ?person ?condition])

(tell [and [has-condition primadonna sore-throat today]
          [effective-treatment gargle-with-ammonia sore-throat]
          [not [allergic primadonna gargle-with-ammonia]]])
```
Command: Enable Joshua Tracing (Type of tracing) Predications

**Predicate** tracing is on
  Tracing **all** predicates
  Traced events: Ask and Tell

Command: Enable Joshua Tracing (Type of tracing) Backward Rules

**Backward Chaining** tracing is on
  Tracing **all** backward rules
  Traced events: Try, Fail, Retry, and Succeed

Command: (ask [appropriate-treatment =drug =person sore-throat] #’print-query)
  ▶ Asking predication [**APPROPRIATE-TREATMENT** =**DRUG**
  =**PERSON** **SORE-THROAT**]
  ▶ Trying backward rule **FIND-CURE** (Goal...)
  ▶ Asking predication [**HAS-CONDITION** =**PERSON**
  **SORE-THROAT** TODAY]
  ▶ Asking predication [**EFFECTIVE-TREATMENT** =**DRUG**
  **SORE-THROAT**]
  ▶ Asking predication [**NOT** (**ALLERGIC** **PRIMADONNA**
  **GARGLE-WITH-AMMONIA**)]
  ▶ Asking predication [not (**ALLERGIC** **PRIMADONNA**
  **GARGLE-WITH-AMMONIA**)]
  ▶ Succeeding backward rule **FIND-CURE**

  [**APPROPRIATE-TREATMENT** **GARGLE-WITH-AMMONIA** **PRIMADONNA** **SORE-THROAT**]
  ▶ Retrying backward rule **FIND-CURE** (Goal...)
  ▶ Failing backward rule **FIND-CURE**

Notice that even after the rule derives the only possible answer and prints the unified query, backward rule tracing displays the **Retrying** and **Failing** events as the rule tries and fails to find another answer.

You can get a greater degree of control over rule tracing by providing the **menu** keyword to Enable Joshua Tracing. This lets you specify particular rules to trace, trace only rules triggered by certain predications, and trace rule importance queuing and dequeuing.

To disable rule tracing use the command

  Disable Joshua Tracing (type of tracing) Forward Rules

for forward rules, or

  Disable Joshua Tracing (type of tracing) Backward Rules

for backward rules.

You can also adjust rule tracing options by mousing on rule names and predications. Mouse right on the object to show the possible options.
5.4. Joshua Rule Basics At a Glance

Figure rule-summary, page rule-summary summarizes the basic rule information we've just covered.

<table>
<thead>
<tr>
<th>Rule Type</th>
<th>Triggered by</th>
<th>Trigger Part</th>
<th>Predications In Trigger</th>
<th>Rule Fires When</th>
<th>Action Part</th>
<th>Rule Success (all if-parts satisfied)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Forward</td>
<td>TELL</td>
<td>if-part</td>
<td>Compound</td>
<td>Triggers Satisfied</td>
<td>then-part</td>
<td>Rule executes action part</td>
</tr>
<tr>
<td>Backward</td>
<td>ASK</td>
<td>then-part</td>
<td>Single</td>
<td>Trigger Asked</td>
<td>if-part</td>
<td>ASK calls continuation</td>
</tr>
</tbody>
</table>

Figure 6. Summary of Joshua Rule Operation

Sometimes the knowledge needed to satisfy a rule trigger can be elicited from the user through the question facility. This works somewhat similarly to backward rules, and we discuss it in the chapter that follows. See the section "Asking the User Questions", page Asking-the-User-Questions-SECTION.
6. Asking the User Questions

We have seen that knowledge stored in the database can be extended by reasoning about it with rules. Questions behave similarly to rules, and are another way of extending knowledge by seeking out information from the user. "User" in this context is a very general term denoting any person, process, or device that a question can interact with.

Questions are like backward chaining rules. A question has a name, a trigger pattern and a body, and you use it in goal-directed inference to satisfy backward chaining goals.

This chapter assumes that you are familiar with the basics of rule operation presented in the chapter "Rules and Inference".

**joshua:ask** does not automatically invoke questions, unless you specify :do-questions t. (In contrast, rules are used by default, unless you specify :do-backward-rules nil.)

Joshua searches for applicable questions after searching the database and running all appropriate backward rules. When a question trigger unifies with the query pattern, the question body runs, and, if successful, calls the **joshua:ask** continuation.

Joshua supplies a default question facility that you can customize if you wish.

6.1. Adding and Removing Joshua Question Definitions

**joshua:defquestion** defines a question. The basic arguments to **joshua:defquestion** are a rule name, a control-structure that specifies forward or backward questions (currently can only be :backward), and a pattern specifying the question’s trigger.

Like rules, questions have control structure arguments :importance and :documentation. The former lets you prioritize the order in which questions are run. The latter lets you add a documentation string to explain what your question does.

In the custom version you supply a question body, using the keyword argument :code before the custom code.

Please consult the dictionary entry for **joshua:defquestion** for a full description of this macro.

The function **joshua:undefquestion** removes a single question definition. To remove a question from a Zmacs buffer, use *-K Kill Definition.
6.2. Default Joshua Questions

The basic components are a name, a :backward control structure, optional control structure arguments, and a trigger pattern. For example:

```
(defquestion question1 (:backward
  [author-of Winnie-the-Pooh Milne])

(defquestion question2 (:backward
  :documentation "Get information about authorship")
  [author-of ?work ?author])
```

Here's what happens if the question is invoked:

- If the query contains no logic variables at run time, a Yes or No question is generated once. (A No answer means the proposition is false or unknown.)

- If the query does contain logic variables at run time, the question loops, presenting iterations of an AVV (Accept Variable Values) menu, each looking for values of the variables that would make the triggers true.

Either the default joshuasay method or a user-written joshuasay method if available is used in formatting the question.

Here are some examples:

```
(define-predicate lost (person object))
(define-predicate in-possession-of (object suspect))
(define-predicate suspects (person spouse))
(define-predicate unfaithful (person to-spouse))

(defquestion where-is-it? (:backward
  [lost ?person ?object])
```

Now we do an joshua:ask, specifying that we want questions to be run.

Example 1: No variables in the query that triggers the question

```
(ask [lost Desdemona handkerchief] #'print-query :do-questions t)
```

Is it true that "[LOST DESDEMONA HANDKERCHIEF]? [default No]: Yes

[LOST DESDEMONA HANDKERCHIEF]

Let's add a user-defined joshuasay method to the above example and do another joshua:ask.

Example 2:

```
(define-predicate-method (say lost) (&optional (stream *standard-output*))
  (with-statement-destructured (person object) self
    (format stream "'A lost the "A" person object)))
```


(ask [lost Desdemona handkerchief] #'print-query :do-questions t)
Is it true that "DESEMOMA lost the HANDKERCHIEF"? [default No]: Yes
[LOST DESDEMOMA HANDKERCHIEF]

In the next example, the query that triggers the question at run time contains variables. The question loops until we indicate that no more solutions exist by clicking on No and on End.
Example 3:

(defrule fidelity-test (:backward)
  (if [and [lost ?person ?object]
     [in-possession-of ?object ?a-suspect]
     [suspects ?spouse ?a-suspect]]
  then [unfaithful ?person ?spouse])

(defun desdemoniad ()
  ;; see if Desdemona is unfaithful to Othello
  (clear)
  (tell [and [in-possession-of handkerchief Cassius]
         [suspects Othello Cassius]])
  (ask [unfaithful Desdemona Othello] #'graph-query-results :do-questions t))
6.3. Writing Custom Questions

To write a custom version of a question, use the basic question format adding the keyword, `code`, and its arguments before you write the question body.

```lisp
(defquestion <question-name> (:backward)
    <trigger-pattern>
    :code
    ((query truth-value continuation &optional query-context)
     <body>)
```

`query` is the query for `joshua:ask`, unified with the trigger pattern of the question. If `truth-value` is `joshua:*true*`, the system is asking whether the statement in true, as opposed to being false or unknown. If `truth-value` is `joshua:*false*`, the system is asking whether the statement is `joshua:*false*`, as opposed to being true or unknown.

The `query-context` argument can almost always be ignored.

The question `body` can be Lisp forms or Joshua commands, and it works like Lisp code in the body of a backward rule. If the value of `body` is `nil`, the query that triggered the question fails. If the value of `body` is non-`nil`, the query succeeds. Calling the `joshua:succeed` function explicitly within the `body` allows the query to succeed many times. This is how questions with variables at run-time can loop, eliciting all possible bindings from the user.

`body` can do anything, including an `joshua:ask` or `joshua:tell`. It should, however, do the following:

- If there are no logic variables in the query, decide somehow (perhaps by asking the user a question), if the query is true. If so, call the continuation.

- If there are logic variables present, solicit sets of bindings for them from somewhere (for example, the user). For each such set, call `joshua:succeed`. 

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Asking the User Questions

April 2003
The question in this example tries to find one or more languages to which a given word belongs. Standard choices are offered in a languages menu. The user can click on one or more of these choices, as well as enter any others. We also ask the function joshua:succeed to return a user-id telling us who answered the question.

To see the results and how they came about, we use the convenience function, joshua:print-query-results in the joshua:ask continuation.

(define-predicate valid-word (word language))

;;; The customized defquestion
(defquestion check-if-valid-word (:backward)
  [valid-word ?word ?language]
  :code
  ((query truth-value continuation &optional ignore)
   (unless (eql truth-value *true*)
     (error "I don't know how to ask if ~S is false." query))
   (typecase ?word
     (unbound-logic-variable
      (error "I don't know how to ask questions about every possible word: ~S" query))
     (otherwise
      (typecase ?language
        (unbound-logic-variable
         (let ((list-of-languages)
             (dw:accepting-values
              (*query-io* :label "Languages" :own-window t)
              (format *query-io*
                "Languages in which ~A is a word" ?word)
              (append
               (accept '(subset english french german swahili sanskrit)
                 :prompt nil)
               (accept '(sequence symbol) :prompt "Others":default nil))))))
           (loop for the-language in list-of-languages
              do (with-unification
                (unify ?language the-language)
                (succeed sys:user-id)))))))
     (otherwise
      (when (dw:accepting-values
        (*query-io* :label "Languages" :own-window t)
        (format *query-io*
          "Is ~A a word in ~A? " ?word ?language)
        (accept 'boolean :prompt nil :prompt-mode :raw :default t))
        (succeed sys:user-id)))))))))

(ask [valid-word boutique ?language] #'print-query-results :do-questions t)
Click on your selections, and they are highlighted. As the next screen shows, we selected two languages from the default, and added another one from the "Others" option. After we click on Done, the joshua:ask continuation executes, printing out the answers to the query and the reason why each answer succeeded. Since we asked joshuasucceed to return a user id telling us who answered the question, we get this information as well.

```
(ask [valid-word boutique ?language] #'print-query-results :do-questions t)
[valid-word boutique french] succeeded
  it was derived from question check-if-valid-word
  "pinhead"
[valid-word boutique english] succeeded
  it was derived from question check-if-valid-word
  "pinhead"
[valid-word boutique esperanto] succeeded
  it was derived from question check-if-valid-word
  "pinhead"
```
7. Pattern Matching in Joshua: Unification

Functions

joshua:unify
joshua:variant
joshua:succeed

Macro

joshua:with-unification

As described in the section "Rules and Inference", pattern matching underlies all inferencing operations in Joshua. In forward chaining, Joshua matches rule trigger patterns with database predications. In backward chaining, it matches goals with database predications and with rule trigger patterns. The type of pattern matching used is called unification.

7.1. Unification Rules

Two predications that contain no variables match (or unify), if they are structurally equivalent, that is, if they "look the same". This is much the same test as the Lisp equal function. For example:

\[
\begin{align*}
\text{[fact a b c]} & \text{ matches } \text{[fact a b c]} \\
\text{[fact a b c]} & \text{ does not match } \text{[fact a b d]}
\end{align*}
\]

The more interesting case for pattern matching is when the predication patterns contain logic variables. Then predications unify if there is a way of substituting values for the variables so that both predications become structurally equivalent. In the simple case, an uninstantiated logic variable matches any object in the equivalent position, becoming instantiated as that object. In the pattern pairs below, for example, these are the matches and substitutions:

Pattern 1: [fact a b c]   
Pattern 2: [fact a b ?x]   

\[
\begin{align*}
\text{\texttt{fact} = fact (fact matches fact)} \\
\text{a = a} & \quad (a \text{ matches } a) \\
\text{b = b} & \quad (b \text{ matches } b) \\
\text{?x --\rightarrow c} & \quad (c \text{ is substituted for } ?x)
\end{align*}
\]

Pattern 1: [fact a b]   
Pattern 2: [fact ?x ?y]   

\[
\begin{align*}
\text{\texttt{fact} = fact (fact matches fact)} \\
\text{?x --\rightarrow a;} \\
\text{?y --\rightarrow b;}
\end{align*}
\]
Advanced Concepts Note:

The "occur-check" states that it is not valid to substitute for a variable a structure containing that variable. One reason for this is that such a substitution might cause the system to draw logically unsound inferences.

The following do not unify:

\[ f \ ?x \ ?x \]

\[ f \ (g \ ?y) \ ?y \]

For a detailed discussion of the "occur-check": See the function \texttt{joshua:unify}, page \texttt{joshua:unify-FUNCTION}.

To determine the correct match in the next example, one needs to know when variables with the same name are identical and when they differ. For instance, why does a trigger pattern like this:

\[ f \ ?x \ ?x \]

match the first predication pattern below, but not the second one.

\[ f \ 	ext{goo} \ 	ext{goo} \quad ;\text{matches trigger pattern} \]

\[ f \ 	ext{silly} \ 	ext{putty} \quad ;\text{does not match trigger pattern} \]

Section "Variables and Scoping in Joshua" covers these scoping rules.

7.2. Variables and Scoping in Joshua

In general terms, the scope of a variable is the area within which it is visible and can be referred to. For more on scoping within Symbolics Common Lisp: See the section "Scoping" in Symbolics Common Lisp Language Concepts.

In Joshua, since logic variables typically have names (?x, ?item, and so on), valid matching is based on understanding when variables with the same name are identical and when they differ. The scoping rules for Joshua determine this.

Logic variables are lexically scoped within rule bodies. That is, logic variables with the same name within a rule body are identical and must match the same object.

- Trigger pattern using ?x: \[ f \ ?x \ 1 \]; this ?x is different from \n- Database predication using ?x: \[ f \ 2 \ ?x \]; this ?x

Matches and substitutions:

- In trigger \[ f \ ?x \ 1 \] \text{ ?x --> 2 becomes } \[ f \ 2 \ 1 \]
- In predication \[ f \ 2 \ ?x \] \text{ ?x --> 1 becomes } \[ f \ 2 \ 1 \]

Used by itself, the question-mark or the equivalence symbol (?) is an anonymous logic variable. Each one you type is different, imposing no scoping constraints, and can therefore be used for "no care" slots. For example:
[foo ? ?] matches [foo 1 2]
and [foo = =] matches [foo 1 2]
whereas [foo ?x ?y] does not match

Logic variable names are new each time a rule is triggered. Conceptually the system makes a copy of the triggered rule. The variable names in the "copy" are the same names as those in the original, but are a different version of these variables. During rule execution this "copy" of the rule is successively modified as new unifications occur. Once rule execution terminates, all bindings are undone and the rule "copy" is discarded.

### 7.3. Some Examples of Joshua Unification

Here are some simple examples:

- [father ?x ?y] unifies with [father john john-jr]
- [father ?x john] unifies with [father george ?y]
- [father ?x ?y] unifies with [father ?z ?z]
- [father ?x john] does not unify with [father ?y george]

Here's a forward rule example with a format statement to help us see what is happening:

```lisp
(define-predicate plays-instrument (person instrument))
(define-predicate owns (person thing))
(define-predicate invite-to-audition (person))

(tell [plays-instrument Jane tuba])

(defrule test-eligibility (:forward)
   if [and [plays-instrument ?person ?instrument]
        [owns ?person ?instrument]]
   then [and [invite-to-audition ?person]
            [format t "% In the rule body ?person is bound to ~A, and ?instrument is bound to ~A.\n" ?person ?instrument]]
)
```

When this rule's trigger pattern ([plays-instrument ?person ?instrument]) is matched against the database, it unifies with [plays-instrument Jane tuba]. That is, ?person is unified with Jane, and ?instrument is unified with tuba.

If we add another joshuatell statement such as [owns Jane tuba], this unifies with the second part of the rule's trigger, and the rule fires. The action part then adds the rule's inference to the database with the bindings that have been established by unification. The format message we added to the action part confirms these bindings:
(tell [owns Jane tuba])
   In the rule body ?person is bound to JANE, and ?instrument
   is bound to TUBA.
[OWNS JANE TUBA]

The next example uses backward chaining and compound subgoals.

Backward rule ownership-1 states that a person owns whatever object s/he has paid
for. Backward rule ownership-2 states that if Fred and Bob own the same item
then it is a desirable item. We then joshua:ask a query whose goal is to deter-
mine what items are desirable. The query succeeds. What items unified?

The query [desirable item] unifies with the trigger of backward rule ownership-2;
that rule's subgoal [owns ?person ?object], then unifies with the trigger of rule
ownership-1. The latter's subgoal [paid-for ?person ?object] unifies with the
database predication [paid-for Fred stereo]. The goal of rule ownership-1 is satis-
fied. Next Joshua tries to satisfy the second part of rule ownership-2's subgoal,
[owns Bob ?item] and so on, until the goal succeeds.

Here are the definitions and joshuatell statements for this example.

;;; Define some additional predicates
(define-predicate paid-for (person thing))
(define-predicate desirable (thing))

(tell [and [paid-for Fred stereo]
       [paid-for Bob stereo]])

(defrule ownership-1 (:backward)
   if [paid-for ?person ?object]
   then [owns ?person ?object])

(defrule ownership-2 (:backward)
   if [and [owns Fred ?item]
          [owns Bob ?item]]
      then [desirable ?item])

If we could look inside the unifier during execution of the query, this is what we
would see:

(ask [desirable ?x] #'print-query)
7.4. Basic Unification Facilities

Joshua provides some unification facilities for use within Lisp code. The dictionary entry for each facility details its use and provides examples.

The function **joshua:unify** unifies expressions within Lisp code embedded in the `if` part of rules, or in the body of a **joshua:defquestion**, or wherever you find it convenient to call it yourself.

The macro **joshua:with-unification** establishes the scope of unifications done within its body, and establishes a place to be thrown to if a unification in its body fails. That is:

- If a unification cannot be done, **joshua:unify** throws to the dynamically innermost enclosing **joshua:with-unification**, and

- The extent of the unification is the dynamic extent of the dynamically innermost enclosing **joshua:with-unification**.

The function **joshua:succeed** will, based on its context, find the continuation and call it accordingly.

The function **joshua:variant** is related to **joshua:unify**, but shouldn’t be confused with it. Whereas **joshua:unify** tries to see if two objects can be *made* the same, **joshua:variant** checks whether two objects *are* the same. Predications that differ only in the names of the logic variables they contain are equivalent, and are *variants* of each other. (**joshua:assert** uses **joshua:variant** to check whether the predication it is adding to the database is already there.)
**joshua:variant** is based on the notion that it should not matter what the names of the logic variables are, so long as the structures are the same. This is a much stronger condition than **joshua:unify**. Pairs that satisfy **joshua:unify** are not necessarily variants, but every pair that satisfies **joshua:variant** also satisfies **joshua:unify**.
8. Using Joshua Within Lisp Code

Here are several basic functions that you can use inside Lisp code. These functions can be grouped under predication facilities and unification facilities. The tables below summarize each function's use. Please consult the respective dictionary entries for more detail and examples.

Predication Facilities:

- **joshua:different-objects** If the arguments are **joshua:eql**, or if either argument is an uninstantiated logic variable, **joshua:different-objects** returns **nil**. Otherwise it returns **t**.

- **joshua:make-predication** Constructs a predication (does not enter it into the database). The [ ] syntax is a reader macro that expands into this.

- **joshua:predication** The base flavor for all predications in Joshua. This works well with functions like **joshua::typep** and **joshua::typecase**.

- **joshua:predicationp** Checks whether an object is built on **joshua:predication**.

- **joshua:with-statement-destructured** Lets you bind Lisp variables to predication arguments.

Unification Facilities:

- **joshua:succeed** Based on context, finds and calls continuation

- **joshua:unify** Unifies two Joshua patterns, while side-effecting any logic variables for the extent of the enclosing **joshua:with-unification**

- **joshua:variant** Checks whether two predications are equivalent, differing only in the names of the logic variables they contain. This does not need to be done under a **joshua:with-unification**; it merely returns **nil** if it fails.

- **joshua:with-unification** Establishes the scope within which substitutions specified by **joshua:unify** take effect; establishes a place to be thrown to if a unification in its body fails.
9. Advanced Features of Joshua Rules

This section summarizes the full syntax of both forward and backward chaining rules.

Both forward and backward rules allow various keywords to be attached to the patterns of the If-part of the rule. Both Forward and Backward rules allow the Keyword :support followed by a logic-variable:

```
(defun rule foobar (:forward)
  (if [and [foo ?x ?y] :support ?f1
    Then (format t "&I won with F1 = "s and F2 = "s" ?f1 ?f2))

(defun rule foobar (:backward)
  (if [and [bar ?x ?y] :support ?f1
    Then [foo ?x ?z])
```

This indicates that the logic-variable should be bound to the "support" for this pattern. In the case of a forward rule, the support is simply the fact which matched the corresponding pattern. Thus

```
(defun rule foobar (:forward)
  (if [and [foo ?x ?y] :support ?f1
    Then [foo ?x ?z])

(defun rule foobar (:backward)
  (if [and [bar ?x ?y] :support ?f1
    Then [foo ?x ?z])
```

will cause the first rule above to print:

```
I won with F1 = [FOO 1 2] and F2 = [BAR 2 3]
```

Backward rules turn their If-part into a series of nested joshuaask's. When the first joshuaask finds a match, it calls a continuation which performs the next joshuaask. The argument to this continuation is a "backward-support" structure, see the section "Continuation Argument", page Continuation-Argument-FRAGMENT.

The support keyword in a backward rule binds the logic-variable to the backward support corresponding to its query.

Thus with the following rule and data:

```
(defun rule foobar (:backward)
  (if [and [bar ?x ?y] :support ?f1
       (progn (format t "&I won with F1 = "s and F2 = "s" ?f1 ?f2)
             (succeed))
    Then [foo ?x ?z])

(defun rule foobar (:backward)
  (if [and [bar ?x ?y] :support ?f1
       (progn (format t "&I won with F1 = "s and F2 = "s" ?f1 ?f2)
             (succeed))
    Then [foo ?x ?z])
```

will cause the first rule above to print:

```
I won with F1 = [FOO 1 2] and F2 = [BAR 2 3]
```

Backward rules turn their If-part into a series of nested joshuaask's. When the first joshuaask finds a match, it calls a continuation which performs the next joshuaask. The argument to this continuation is a "backward-support" structure, see the section "Continuation Argument", page Continuation-Argument-FRAGMENT.

The support keyword in a backward rule binds the logic-variable to the backward support corresponding to its query.

Thus with the following rule and data:

```
(defun rule foobar (:backward)
  (if [and [bar ?x ?y] :support ?f1
       (progn (format t "&I won with F1 = "s and F2 = "s" ?f1 ?f2)
             (succeed))
    Then [foo ?x ?z])

(defun rule foobar (:backward)
  (if [and [bar ?x ?y] :support ?f1
       (progn (format t "&I won with F1 = "s and F2 = "s" ?f1 ?f2)
             (succeed))
    Then [foo ?x ?z])
```
The query:

(ask [foo 1 3] '#'print-query)

will cause the following output:

I won with F1 = ([BAR 1 2] 1 [BAR 1 2]) and F2 = ([BAR 2 3] 1 [BAR 2 3])
[F00 1 3]

This backward support may be used to provide a justification (for a TMS) when a backward rule caches the results of its work, as follows:

(defun foobar (:backward)
  (if [and [bar ?x ?y] :support ?f1
       (progn
         (tell [foo ?x ?z]
          :justification '(foobar
                          ,(ask-database-predication ?f1)
                          ,(ask-database-predication ?f2)))
          (succeed))
       ]
  Then [foo ?x ?z])

Backward rules also support two other keywords :do-backward-rules and :do-questions. These can be used to control the behavior of the joshua:ask corresponding to a backward action. If the :do-backward-rules keyword is present then the value following it should evaluate to either joshua::t or joshua::nil; if it is joshua::nil, then this query will not attempt to use rules to satisfy the query, otherwise rules will be used. Similarly, the :do-questions questions controls whether backward questions will be invoked to query the user. The default value is that backward rules are used and that questions will be attempted if the query which caused this rule to be invoked allowed questions to be used.
10. Justification and Truth Maintenance

Functions:  joshua:explain
            joshua:unjustify
            joshua:graph-tms-support

Keyword:   :justification

Symbols:   :premise
           :assumption
           :none

A Truth Maintenance System (TMS) is a tool used by deductive systems to keep
track of interdependencies among statements in a database.

A TMS has two main functions:

- Recording and maintaining the reasoning that supports the current set of predi-
cations in the database
- Maintaining the logical consistency of these predications

Since very often database predications logically depend on each other as deter-
dined by rules, keeping track of these dependency relationships lets the system ex-
plain its reasoning process. Moreover, knowing the dependencies for each predica-
tion lets the system work out the consequences of changes to the truth values of
predications and modify the database to keep it current and consistent. Database
modifications involve retracting (removing) facts. This means that a justification is
a bidirectional link, from fact to conclusion and back.

Since a TMS may not always be necessary, its inclusion is optional in Joshua.
Without a TMS, Joshua records the truth value of each predications, and changes
this value if new joshuatell statements assert a different truth value. See the sec-
tion “Truth Values”, page Truth-Values-SECTION.

However, without a TMS, the system has no knowledge of the reasons supporting
its beliefs, and hence no awareness of logical contradictions. Once the database be-
comes inconsistent it remains so unless you modify it under program control. On
the other hand, you don’t pay either the space or time penalties of having a TMS.
The tradeoff is up to you.

An inconsistent database may or may not be acceptable depending on the problem.
If you do want to use a TMS, Joshua currently supports a clausal TMS (LTMS)
that you can mix into your predicate definitions. This TMS is based on David
McAllester’s three-valued TMS (Massachusetts Institute of Technology, Artificial
Intelligence Lab, A.I. Memo 473, 5 May, 1978).

This chapter is an overview of how Joshua works with the LTMS.

To include a TMS in your application, you specify the particular TMS model you
are using (LTMS, JTMS, any other) as an argument to your predicate definition.
For instance:
(define-predicate temperature-of (object temperature)
(ltms:ltms-predicate-model))

Here, as in the following examples, we use the supplied LTMS model, which resides in its own package.

**Advanced Concepts Note:**

You can incorporate any TMS of your choice into your Joshua application. As with any other tool you want to build into Joshua this is straightforward, since Joshua talks to Truth Maintenance Systems via a generic protocol; thus you need only write protocol methods for the TMS generic functions. Joshua’s facilities for supporting external TMS systems are discussed in the *Joshua Reference Manual*.

With a TMS model Joshua can *justify*, *explain*, and *revise* its beliefs. We explore each of these activities in the following sections.

### 10.1. Justification

Joshua *believes* a fact to be valid if it is in the database and has a truth value of *joshua:*true* or *joshua:*false*. (A truth value of *joshua:*unknown* denotes a fact whose validity is currently not known.) Without a TMS the system records truth values, but not the reasons for them.

See the section "Truth Values", page Truth-Values-SECTION.

When a TMS is present, a believed predication must have a *justification*. This means there must be at least one currently valid reason supporting the predication’s truth status. When *joshuatell* enters a new predication it provides a justification for it. The TMS records this justification as part of the information it maintains about the database predication. Truth values and their justification status can change as a result of new information and resulting action taken by the TMS, or as a result of program action.

Because the LTMS ensures data consistency, it does not accept logically contradictory facts; thus it cannot believe simultaneously predication [P ...] and predication [NOT [P ...]] (that is, a predication cannot be *joshua:*true* and *joshua:*false* at the same time). If this happens, one of these statements must be retracted by removing its justification (*joshua:unjustify*).

TMS justifications can be either *primitive*, or *compound*. The terms *primitive* and *compound* refer to the degree of reasoning used in generating the justification.

#### 10.1.1. Primitive Justifications

Primitive justifications are primitive with respect to the reasoning process. (The system remembers no further explanation.) Primitive justifications are usually specified by the programmer or by system defaults in a *joshuatell* used outside a rule. This type of justification which depends on nothing except itself is the basic component of the belief structure.
There are three types of primitive justification:

- Assumptions
- Premises
- None

Assumptions justify predications whose retraction you are willing to leave to the TMS.

Premises justify predications whose retraction you want to control yourself, rather than leave them to the TMS.

As we shall see, when the TMS finds a logical contradiction that can be traced back to a single assumption, it retracts that assumption automatically. If the inconsistent predication depends only on premises, or if there is more than one assumption in the support, the LTMS signals a condition and requests you to take care of correcting the inconsistency. That usually means using the Debugger; programs can handle the condition to do otherwise. For more on condition-handling: See the section "Conditions" in Symbolics Common Lisp Programming Constructs.

None justifications are justifications that don’t really cause the predication to be believed. That is,

```
(tell [foo ?x] :justification :none)
```

puts the predication in the database, but keeps its truth value at joshua:\texttt{*unknown*}. (You might want to do this because you are using joshuataell to canonicalize the predication, but you don’t yet want the system to believe it.)

When the LTMS is given a justification for a predication in the database, the LTMS builds a clause consisting of the predication and its support.

If a new predication is being added by a joshuataell outside a rule context and if you supply no explicit justification, the default justification is premise. To see this, enable the tracing of TMS operations for the next examples.

Enable Joshua Tracing (Type of tracing) TMS Operations

---

TMS tracing is on
Tracing All TMS predicates
Traced events: Contradiction, Justify, and Unjustify

```
(define-predicate temperature-of (object temperature)
  (ltms:ltms-predicate-model))
```

Now joshuataell a fact about the temperature of some object. The trace shows that the new predication by default is being established in the database as a premise.

```scheme
(tell [temperature-of skin cool])
```

```
► Justifying: [TEMPERATURE-OF SKIN COOL] <- PREMISE
[TEMPERATURE-OF SKIN COOL]
T
```
To specify a justification yourself, use the keyword `justification` within a joshuatell statement, followed by the justification (a symbol or a list). If the justification is a symbol, a primitive justification is built whose mnemonic is the symbol. `assumption` is the symbol for the primitive justification assumption. `premise` is the symbol for the primitive justification premise, and `none` is the symbol for the primitive justification none.

```
(tell [temperature-of surface very-hot]
   :justification :assumption)

► Justifying: [TEMPERATURE-OF SURFACE VERY-HOT] <-- ASSUMPTION
  [TEMPERATURE-OF SURFACE VERY-HOT]
```

10.1.2. Compound Justifications

There are several types of compound justification. Here we deal with the most basic type, namely, deduction. We discuss others in the Joshua Reference Manual.

When Joshua joshuatells the database a new predication inferred from a forward chaining rule, the default justification is deduction; this states that the fact depends on the executing rule and the predications that triggered it. These predications may in turn be deductions from other rules; ultimately the chain of dependencies leads to primitive support (premises and/or assumptions).

Joshua can extract the primitive and compound justification for a database predication and interpret it for you. See the section "Explaining Program Beliefs", page Explaining-Program-Beliefs-SECTION.

10.1.3. Database Predications Can Have Multiple Justifications

Every time you (or the program) joshuatell a statement, it is given a new justification. So although there is only one version of the statement in the database, it can have multiple justifications. For example, earlier we added the predication [temperature-of surface very-hot] to the database as an assumption. Now we add a joshuatell statement that triggers the forward rule determine-temp1.

```
(define-predicate condition-of (object condition) (ltms:ltms-predicate-model))

(tell [condition-of surface melting])

(defrule determine-temp1 (:forward)
  if [condition-of ?object melting]
  then [temperature-of ?object very-hot])
```

When this rule fires, it infers [temperature-of surface very-hot]. Since this predication is already in the database, the system does not reinsert it. Rather, it adds a new justification for it that mentions the rule determine-temp1. Note, however, that the "support" is always the current justification; the system only uses the new justification if the current one for some reason becomes obsolete. For example, if you remove the current justification, the next justification, if any, becomes the predi-
cation’s support. The predication’s truth value does not become joshua:*unknown* until every one of its justifications has been removed. For more on removing a predication’s support:

See the section "Retracting Predications with joshua:unjustify", page Retracting-Predications-with-JOSHUA:UNJUSTIFY-SECTION.

10.2. Explaining Program Beliefs

The function joshua:explain provides information about the justification(s) for a database predication. joshua:explain is the forward chaining dual to joshua:ask continuation functions such as joshua:graph-query-results and joshua:print-query-results that trace backward chaining support for a satisfied query. (The LTMS does keep some notes about the results of backward chaining, namely the facts in the database from which you have chained.)

Here’s a forward rule, love-medieval-style, to help us deduce who might properly engage in a courtly love affair. One of the facts that triggers this rule, the predication [is-attached ...], is deduced by the forward rule engaged.

```
(define-predicate is-engaged (lady husband) (ltms:ltms-predicate-model))
(define-predicate is-attached (lady husband) (ltms:ltms-predicate-model))

(defun courtly-setup ()
  (clear)
  (tell [and [noble-lady Isolde]
            [outranks Isolde Tristan]
            [in-proximity Isolde Tristan]
            [is-engaged Isolde Mark]])
```

(define love-medieval-style (:forward)
  (if [and [noble-lady ?lady]
          [outranks ?lady ?gentleman]
          [in-proximity ?lady ?gentleman]
          [is-attached ?lady ?husband]
          (different-objects ?gentleman ?husband)]
    [good-to-accept-as-lover ?lady ?gentleman])

Now we ask what the system knows about acceptable lover candidates, and why. 
joshuaexplain must be given the actual database predication in order to extract
its justification. One simple way to do this is to use the convenience function, 
joshuamap-over-database-predications.

(map-over-database-predications [good-to-accept-as-lover ?x ?y] #'explain)

[GOOD-TO-ACCEPT-AS-LOVER ISOLDE TRISTAN] is true
  It was derived from rule LOVE-MEDIEVAL-STYLE
  It is a :PREMISE
[NOBLE-LADY ISOLDE] is true
  It is a :PREMISE
[OUTRANKS ISOLDE TRISTAN] is true
  It is a :PREMISE
[IN-PROXIMITY ISOLDE TRISTAN] is true
  It is a :PREMISE
[IS-ATTACHED ISOLDE MARK] is true
  It was derived from rule ENGAGED
  It is a :PREMISE
[IS-ENGAGED ISOLDE MARK] is true
  It is a :PREMISE

joshuaexplain extracts and displays the chain of support through forward rules to
primitive support (premises and assumptions) together with their truth values.

To see the same information in graph form, use the function joshua:graph-tms-
support. (This is analogous to the function joshua:graph-query-results that
graphs backward support.)

Figure graph-support uses the previous example and shows the graph of the justi-
fication.

(define love-medieval-style (:forward)
  (if [and [noble-lady ?lady]
          [outranks ?lady ?gentleman]
          [in-proximity ?lady ?gentleman]
          [is-attached ?lady ?husband]
          (different-objects ?gentleman ?husband)]
    [good-to-accept-as-lover ?lady ?gentleman])

Figure 7. Sample Graph of TMS Support
Since all objects in the graph are in the database, all are drawn inside rectangles. The top of the graph shows the database predication whose justification is being traced. The graph traces the justification through rules to underlying premises. The arrows move up from the primitive support through intermediate support to the object being explained.

Justifying and explaining the reasons underlying current knowledge are not static activities, since the flow of new information makes the database subject to constant modification. The next section surveys the role played by the TMS in keeping the knowledge base updated and logically consistent. See the section "Revising Program Beliefs", page Revising-Program-Beliefs-SECTION.

10.3. Revising Program Beliefs

To reason effectively, a program must be able to revise beliefs when new knowledge contradicts them. A contradiction comes about when a valid justification for accepting a new belief conflicts with the justification for believing an already existing predication. A truth value of joshua:*contradictory* is temporarily assigned to the new predication; this alerts the TMS to correct the contradiction. Since contradictory facts cannot be simultaneously accepted, one of them must be retracted by having its justification(s) removed (joshua:unjustify). The TMS keeps current information consistent with new information by allowing the updating of current knowledge in accord with new reasons.

You can let the TMS handle contradictions, or you can bind into your program your own code to handle these conditions. For more on the signalling and handling of conditions: See the section "Conditions" in Symbolics Common Lisp Programming Constructs.

This is what happens in the absence of bound condition handlers: when the TMS detects a contradictory predication (in the case of an LTMS, a clause that cannot be satisfied) it traces backward through the reasons for the conflicting belief and finds the primitive support underlying it.

If a single assumption is causing the contradiction, the TMS automatically does an joshua:unjustify operation on it, and continues processing. joshua:unjustify removes the current support for the predication. If the predication had a single support, joshua:unjustify changes its truth value to joshua:*unknown*. If there were auxiliary justifications, they may change the truth value from joshua:*unknown* back to either joshua:*true* or joshua:*false*. (See the section "Retracting Predications with joshua:unjustify", page Retracting-Predications-with-JOSHUA:UNJUSTIFY-SECTION.)

If more than one assumption underlies the contradiction, or if the contradiction rests on premises rather than assumptions, the TMS signals an appropriate condition; if no program has bound a handler for this condition the TMS invokes the Debugger, listing all the assumptions and premises on which the contradiction depends.
The Debugger offers you the choice of retracting some justifications or of aborting out. After you make your retraction(s), processing continues until the current operation succeeds, or the TMS finds another contradiction.

10.3.1. An LTMS Example

Following is a sequence of examples showing the changes to a database as the processing of new joshuatell statements gives rise to contradictions and subsequent retractions. (For simplicity, we assume a single justification for each predication.) The first example, shown in figure tms-1, illustrates a contradiction resulting from the attempt to enter a new fact into the database. Circled numbers in the figure indicate the sequence of events.

We’ve added step-by-step comments as well as Trace and Debugger displays to show interaction with the system. The comments are numbered to correspond to the figure numbers.

The definitions and joshuatell statements for this example appear at the end of this section. See the section “Definitions for the LTMS Example”, page Definitions-for-the-LTMS-Example-SECTION.

1. We begin with an initial database and two forward chaining rules, as shown in figure 8. To see what happens, we turn on the Trace facility for TMS operations before proceeding.

   Enable Joshua Tracing (Type of tracing) TMS Operations

   TMS tracing is on
   Tracing All TMS predicates
   Traced events: Contradiction, Justify, and Unjustify

2. A new joshuatell statement is entered and the system attempts to insert it into the database.

3. The TMS detects a contradiction between the new joshuatell statement, [not [has-condition John pneumonia today]], and its opposite assertion, [has-condition John pneumonia today], currently residing in the database. The contradiction rests on premises and since we have not specified how to handle the condition, the Debugger is invoked. (Without a TMS, Joshua would simply have changed the truth value of the old statement to joshua:*false*.) Figure tms-system1 shows the Trace display and the user interaction with the system after the TMS intervenes.

4. We decide to joshuaunjustify the belief that John does have pneumonia. Note that we do so by pressing the RESUME key and clicking on those leaves of the support graph that we want to joshuaunjustify. The Trace display in figure 9 shows that the truth value of the statement [has-condition John pneumonia today] is now changed from joshua:*true* to joshua:*unknown*.
Figure 8. TMS Example -- New Fact Contradicts Existing Fact

Figure 9. TMS Example -- Trace and Debugger Displays, and First Retraction
5. The removal of the original statement allows the system to insert the new
\textit{joshuastell} statement, \texttt{[not [has-condition John pneumonia today]]}, into the
database.

6. The database with one former statement retracted, and one new statement
added under the heading "False things".

Continuing from this situation, the next example illustrates a contradiction
generated by the firing of a forward chaining rule. Figure tms-2 shows what
happens.

7. A compound \textit{joshuastell} statement of two predications joined by \textit{joshua::and}
is added to the database. These statements trigger the forward rule, \textit{allergy-test}.

8. Rule \textit{allergy-test} fires and attempts to add a newly deduced fact, \texttt{[not [al-
lergic John penicillin]]}, to the database. Once again, there is a contradic-
tion; the database has a belief that John is allergic to penicillin, while the
rule deduces that he is not. (Evidently, someone performed the experiment of
giving John penicillin despite John's known allergy to it and John is still
alive, thus posing a problem for the TMS.)

9. Bowing to the facts, we ask the system to \textit{joshua::unjustify} the belief that
\texttt{[allergic John penicillin]}.

Figure tms-system2 shows the Trace display as the new statements are added
to the database, the user interaction with the system after the TMS inter-
venes, and the Trace display showing the retraction. [We have artificially bro-
ken up the system displays into separate figures to make each step easier to
follow. In real life, the entire sequence happens in one continuous event until
all contradictions arising from a \textit{joshuastell} statement are resolved.]

As the Trace display at the bottom of the figure shows, the truth value of the
predication \texttt{[allergic John penicillin]} that we have retracted is changed
from \texttt{joshua:*true*} to \texttt{joshua:*unknown*}.

10. The belief deduced by rule \textit{allergy-test}, namely, \texttt{[not [allergic John peni-
cillin]]}, is now inserted into the database.

11. This is the modified database.

Continuing from this point, the final sequence shows how retractions can
cause inconsistencies in the database, and how the TMS automatically cor-
rects these. Figure tms-3 illustrates.

12. One predication in the database, namely, \texttt{[schedule-test-for-allergies John
now]} depended for its support on a predication \texttt{[allergic John penicillin]}
that we retracted in step 9. Without support, a predication can no longer be
believed.
Figure 10. TMS Example, Continued -- Deduced Fact Contradicts Existing Fact

13. The TMS knows about all dependencies, and automatically retracts the unsupported predication [schedule-test-for-allergies John now]. The Trace display in Figure tms-system3, shows this retraction.

14. This shows the final, updated, and consistent database.
10.3.1.1. Definitions for the LTMS Example

;;; Predicate definitions
(define-predicate allergic (person drug)
  (ltms:ltms-predicate-model)))

(define-predicate had-treatment (person condition drug when)
  (ltms:ltms-predicate-model)))

(define-predicate test-for (problem person when-done)
  (ltms:ltms-predicate-model)))

(define-predicate schedule-test-for (problem person when)
  (ltms:ltms-predicate-model)))

(define-predicate has-condition (person condition when)
  (ltms:ltms-predicate-model)))

;;; Rule definitions
(defrule schedule-tests (:forward)
  if [and [allergic ?person ?drug]
    [test-for ?problem ?person never-done]]
  then [schedule-test-for allergies ?person now])

(defrule allergy-test (:forward)
  if [and [had-treatment ?person ?condition ?drug yesterday]
    [not [has-condition ?person ?condition today]]
    [has-condition ?person health today]]
  then [not [allergic ?person ?drug]])
Database

True things
- [HAD-TREATMENT JOHN PNEUMONIA PENICILLIN YESTERDAY]
- [HAS-CONDITION JOHN HEALTH TODAY]
- [ALLERGIC JOHN PENICILLIN]
- [TEST-FOR ALLERGIES JOHN NEVER-DONE]
- [SCHEDULE-TEST-FOR ALLERGIES JOHN NOW]

False things
- [ALLERGIC JOHN PENICILLIN]
- [HAS-CONDITION JOHN PNEUMONIA TODAY]

Rules

(defrule schedule-tests (:forward)
  if [and [allergic =person =drug]
       [test-for =problem =person never-done]
  then [schedule-test-for allergies =person now])

(defrule allergy-test (:forward)
  if [and [had-treatment =person =condition =drug yesterday]
       [has-condition =person health today]
  then [not [allergic =person =drug]])

Changes to Database

True things
- [HAD-TREATMENT JOHN PNEUMONIA PENICILLIN YESTERDAY]
- [HAS-CONDITION JOHN HEALTH TODAY]
- [TEST-FOR ALLERGIES JOHN NEVER-DONE]
- [SCHEDULE-TEST-FOR ALLERGIES JOHN NOW]

False things
- [ALLERGIC JOHN PENICILLIN]
- [HAS-CONDITION JOHN PNEUMONIA TODAY]

Final, Updated Database

True things
- [HAD-TREATMENT JOHN PNEUMONIA PENICILLIN YESTERDAY]
- [HAS-CONDITION JOHN HEALTH TODAY]
- [TEST-FOR ALLERGIES JOHN NEVER-DONE]

False things
- [ALLERGIC JOHN PENICILLIN]
- [HAS-CONDITION JOHN PNEUMONIA TODAY]

Figure 12. TMS Example, Concluded -- TMS Automatically Retracts Unsupported Fact

Figure 13. TMS Example, Concluded -- Automatic Retraction by the TMS
10.3.2. Retracting Predications with joshua:unjustify

You can use the function joshua:unjustify to remove the support for specific predications from the database independently of action by the TMS. (Without a TMS, joshua:unjustify just sets the truth value of the predication to joshua:*unknown*.)

Note that when more than one justification supports a predication, joshua:unjustify must be called once for each justification. (Like joshua:explain, joshua:unjustify needs the actual database predication that you want to operate on, not a copy of it that you type in.)

In an earlier example we used a predication with two justifications. See the section "Database Predications Can Have Multiple Justifications", page Database-Predications-Can-Have-Multiple-Justifications-SECTION. First we inserted this predication, [temperature-of surface very-hot], into the database as an assumption. Later it was deduced from forward rule determine-temp1 shown below, thus receiving an additional justification as a deduction.

```
(defun setup-initial-conditions ()
  (clear)
  (tell [and [allergic john penicillin]
           [has-condition john pneumonia today]
           [test-for allergies john never-done]])

(defun create-contradiction ()
  (tell [not [has-condition john pneumonia today]])

(defun create-contradiction2 ()
  (tell [and [had-treatment John pneumonia penicillin yesterday]
         [has-condition john health today]])
```

Let’s enable the tracing of TMS operations to see what happens when we try to joshua:unjustify this predication. (We get at the predication by clicking on it in the database display.)
Show Joshua Database
True things
[TEMPERATURE-OF SURFACE VERY-HOT]
[CONDITION-OF SURFACE MELTING]
False things
None

(unjustify [TEMPERATURE-OF SURFACE VERY-HOT])
  ▶ Unjustifying: [TEMPERATURE-OF SURFACE VERY-HOT]
  ▶ Justifying: [TEMPERATURE-OF SURFACE VERY-HOT] ← Rule:
    DETERMINE-TEMP1
    NIL

Although joshua:unjustify removed the current support for the predication, the Trace display shows that the predication is still joshua:*true*, because it was supported by an additional justification which has now become its current support. The support says that the predication was deduced from a forward rule. This deduction remains valid as long as the predications supporting the rule’s conclusion are valid. Thus, trying to joshua:unjustify this deduced predication now would cause an error (with the LTMS you can only joshua:unjustify primitive support).

(unjustify [TEMPERATURE-OF SURFACE VERY-HOT])
Error: Predication [TEMPERATURE-OF SURFACE VERY-HOT] can’t be unjustified,
  its support is ¥<DETERMINE-TEMP1
  [CONDITION-OF SURFACE MELTING] → [TEMPERATURE-OF SURFACE VERY-HOT]>

This section concludes our survey of basic Joshua concepts. You now have a working knowledge of Joshua that enables you to build applications using the default Joshua facilities.

The Joshua functions and commands covered in this manual appear in the "Basic Joshua Dictionary" immediately following this section.

Advanced Joshua concepts are covered in the companion volume to this. See the document Joshua Reference Manual.
11. Dictionary Notes: Basic Joshua Dictionary

The entries in this dictionary are a subset of those in the "Joshua Language Dictionary". We have included here only those functions and commands that you will find useful for getting started and whose operation has been discussed in the conceptual portions of this manual. Advanced functions and functions needed for modeling are omitted.

Here is the alphabetized list of Joshua language objects included in this dictionary.
11.1. List of Entries in the Basic Joshua Dictionary

joshua:ask
joshua:ask-database-predication
joshua:ask-derivation
joshua:ask-query
joshua:ask-query-truth-value
joshua:clear
"Clear Joshua Database Command"
joshua:*contradictory*
joshua:copy-object-if-necessary
joshua:define-predicate
joshua:defquestion
joshua:defrule

joshua:different-objects
"Disable Joshua Tracing Command"
"Enable Joshua Tracing Command"
"Explain Predication Command"
joshua:explain
joshua:*false*
joshua:graph-query-results
joshua:graph-tms-support
joshua:known
joshua:make-predication
joshua:map-over-database-predications
joshua:predication
joshua:predicationp
joshua:print-query
joshua:print-query-results
joshua:provable
"Reset Joshua Tracing Command"
""\"Say\"
"
joshua:say
joshua:say-query
"Show Joshua Predicates Command"
"Show Joshua Rules Command"
"Show Joshua Tracing Command"
"Show Rule Definition Command"
joshua:succeed
joshua:tell
joshua:*true*
joshua:undefine-predicate
joshua:undefquestion
joshua:undefrule
joshua:unify
joshua:unjustify
joshua:*unknown*
joshua:untell
joshua:variant
joshua:with-statement-destructured
joshua:with-unbound-logic-variables
joshua:with-unification
12. Basic Joshua Dictionary

`joshua:ask` query continuation &key (:do-backward-rules \( t \)) :do-questions

Queries the virtual database and backward rules and questions.

**Note:** `joshua:ask` is a macro, and as such it cannot be used as an argument to the function `funcall`.

- **query** Should be a predication.
- **continuation** Should be a function of one argument, describing what you want done with the answers to the query.

Note that the argument given to `continuation` might be ephemeral in one of two ways: it could be stack-consed, and it could contain logic variables whose bindings will be undone when you exit this frame. Instantiated queries almost always need to be copied with `joshuacopy-object-if-necessary`, because the variable bindings are ephemeral. See example 6 below.

If, on the other hand, you are collecting database predications, they are not ephemeral, and you don’t want to copy them. (Copying a database predication causes loss of the database information associated with the predication.)

**Keywords:**

- **:do-backward-rules** If this keyword has a non-nil value, backward chaining rules are checked for solutions. The default is \( t \). Use :do-backward-rules nil to check out just the database solution.

- **:do-questions** If this keyword has a non-nil value, any questions that claim to answer query are run to solicit more solutions from the user. The default is nil.

`joshua:ask` uses the database, backward rules, and questions to satisfy the query predication. Each time `joshua:ask` finds a solution to query it calls the continuation, passing it a list that contains the answer and information about how the answer was derived.

`joshua:ask` doesn’t return an interesting value. Normally the continuation performs some action with each solution. You can collect values in the continuation, or return a value to some caller of `joshua:ask` using throw, return-from, or some similar Lisp form. Such uses of throw and return-from are like the Prolog cut feature. See examples 6 through 9.
Any logic variables used in *query* can be referred to as though they were lexical Lisp variables within *continuation*; *joshua:ask* establishes a binding contour for the logic variables. (See example 1 below.) In this sense, *joshua:ask* is like *let* combined with *mapc*. Like *let*, *joshua:ask* establishes lexical binding contours for the logic variables in the query. Like *mapc*, it iteratively calls the continuation on the answer. For a discussion of scoping rules: See the section "Variables and Scoping in Joshua", page Variables-and-Scoping-in-Joshua-SECTION.

*joshua:ask* calls the continuation function with a single argument, *backward-support*, a list containing information about the solution process. The list contains the instantiated query, its truth value, and the support for the query; the form of the support varies, depending on how the query was satisfied.

Typically you'll want to deal only with part of the information provided in *backward-support* rather than with the entire list. For instance, you might want to see only the answer, or only the database predication that matched the answer, or only the support for the answer.

Joshua supplies *accessor functions* to extract various elements of the list in *backward-support*, making it available to you for interpretation.

In addition, Joshua provides *convenience functions* that extract some element of the list in *backward-support* and interpret it for you. These functions let you postpone dealing with the details of *backward-support* and accessor functions until you need them for more advanced work. So before reading on you might want to skip ahead to the section "Streamlining Typical Continuation Requests with Convenience Functions" and see if these functions meet your current needs.

**Continuation Argument**

*backward-support*  A list of the following form:

- The first element is always the unified query, that is, the query that was passed to *joshua:ask*, with appropriate variables instantiated as a side-effect of unification.

- The second element is the truth value of the query. This corresponds to the truth value of the matching predication in the database at the time *joshua:ask* looked at it.

- The rest of the elements are the support for the instantiated query. The support can take several forms, depending on how the query was satisfied.
  - When the query is satisfied by matching a predication in the database, the support is that database object.
° When the query answer comes from a conjunction (and), the support is the symbol and, followed by the backward support for each of the compound predications.

° When the query answer comes from a disjunction (or), the support is the symbol or, followed by the support for the single predication from the or that succeeded.

° When the query answer is derived from a backward rule, the support has the format

    ((rule rule-name) . rule-support)

where

- rule is the symbol rule
- rule-name is the name of the rule used to satisfy the query
- rule-support is a list containing (recursively) the backward support used to satisfy parts of the rule body.

° When the query answer comes from a question, the support is like that for rules, except that it uses the question name instead of the rule name.

° When the query answer comes from the predicates joshuaknown or joshuaprovable, the support is the respective symbol name (joshuaknown or joshuaprovable), followed by the support for the predication that served as the symbol’s argument.

° When the query originates from an joshua:ask or an joshua:ask-data method, the support is whatever the writer of that method provided. See the section “Customizing the Data Index” in Joshua Reference Manual.

In schematic form, the backward-support list looks as follows:

The backward-support list:

    (<unified query> <truth-value> . <derivation>)

    (<(unified) query>)

    (<t/f>)
(derivation) Possibilities for these elements are:
(database predication)
(AND <conjunct1 derivation> <conjunct2 derivation> ...)
(OR <successful disjunct derivation>)
((RULE <rule name>) <conjunct1 t/f derivation> <conjunct2 t/f derivation> ...)
((QUESTION <question name>) <succeed argument>)
(KNOWN <derivation>)
(PROVABLE <derivation>)

Extracting Parts of the Continuation with Accessor Functions

Joshua provides four accessor functions to extract specific portions of backward-support. Use these functions if you want to interpret the answer yourself. Use the convenience functions described below if you want the system to interpret the information for you.

joshua:ask-query Extracts the instantiated query (the first element) from backward-support. For example:
(ask [...] #'(lambda (backward-support)
               (print (ask-query backward-support))))

joshua:ask-query-truth-value
Extracts the truth value of the instantiated query (the second element) from backward-support. For example:
(ask [...] #'(lambda (backward-support)
               (print
                (ask-query-truth-value backward-support))))

joshua:ask-database-predication
Extracts the database object that matched query. If the backward support is a rule, displays the rule name (see example 4). Use this function only when you know the support is a database object (that is, with :do-backward-rules nil). For example:
(ask [...] #'(lambda (backward-support)
               (print (ask-database-predication backward-support)))
               :do-backward-rules nil)

joshua:ask-derivation
Extracts the support information in backward-support. Makes fewer assumptions than joshua:ask-database-predication about where the support came from. For example:
(ask [...] #'(lambda (backward-support)
               (print (ask-derivation backward-support)))
               :do-backward-rules nil)
Streamlining Typical Continuation Requests with Convenience Functions

When an joshua:ask query succeeds, there are some standard things you might want to do with the answer, such as: printing or formatting the unified query, operating on the database predication supporting the query, or interpreting all of the backward support.

Joshua provides five convenience functions that extract an appropriate part of the answer and interpret it in some specific way. The first four are joshua:ask continuation functions. The fifth is a special-purpose function that lets you do database lookup only, and interpret the answer in some way. joshua:map-over-database-predications uses joshua:ask to search the database and extract the predication(s) matching its argument pattern.

These functions are:

- joshua:print-query Extracts and displays the unified query. For example:
  
  (ask [...] #'print-query)

- joshua:say-query Extracts the unified query and displays it in formatted form.

- joshua:print-query-results Takes the information in backward-support and displays it with annotations.

- joshua:graph-query-results The above in graph form.

- joshua:map-over-database-predications For special cases of the solution process, where you look only in the database for an answer, extracts all database predications that unify with a predication pattern and applies some function to each. For example:

  (map-over-database-predications [foo ?x] #'untell)

  joshua:map-over-database-predications is equivalent to:

  (ask query #'(lambda (x) (funcall continuation
                      (ask-database-predication x)))
       :do-backward-rules nil)

We use some of the convenience functions in the examples to joshua:ask. For more on each function, please consult its dictionary entry.
Examples of Using joshua:ask

Let’s define some predicates, enter them into the database, then add a backward rule and a backward question. The rule determines what is an eater’s favorite food. The question elicits information to satisfy the rule’s subgoal.

\[
\begin{align*}
\text{(define-predicate favorite-meal (eater food))} \\
\text{(define-predicate guzzles (eater food))}
\end{align*}
\]

\[
\begin{align*}
\text{(defun eat-it ()} \\
\text{\quad (clear)} \\
\text{\quad (tell [and [favorite-meal bears honey]}} \\\n\text{\quad \quad [favorite-meal mosquitoes people]}} \\\n\text{\quad \quad [favorite-meal spiders flies]}} \\\n\text{\quad \quad [favorite-meal monkeys bananas]}} \\\n\text{\quad \quad [guzzles ted ice-cream]])} \\\n\text{\quad (cp:execute-command "Show Joshua Database"))}
\end{align*}
\]

Show Joshua Database

True things

[FAVORITE-MEAL BEARS HONEY]
[FAVORITE-MEAL MOSQUITOES PEOPLE]
[FAVORITE-MEAL SPIDERS FLIES]
[FAVORITE-MEAL MONKEYS BANANAS]
[GUZZLES TED ICE-CREAM]

False things

None

\[
\begin{align*}
\text{(defrule not-finicky (:backward)} \\
\text{\quad if [guzzles ?eater ?food]}} \\\n\text{\quad then [favorite-meal ?eater ?food]}
\end{align*}
\]

\[
\begin{align*}
\text{(defquestion guzzler? (:backward)} \\
\text{\quad [guzzles ?eater ?food])}
\end{align*}
\]

Next we joshua:ask what Joshua knows about everybody’s favorite meals. Example 1 uses the variables in the unified query to print an English-like sentence (not fussy about number agreement between subject and verb) about everybody’s meals. It ignores the backward-support argument and uses a format directive. It looks in the database and rules, but not in questions.
Example 1.
(ask [favorite-meal ?eater ?food]
  #'(lambda (ignore)
      (format t "-"S is the preferred food of -S." ?food ?eater)))
BANANAS is the preferred food of MONKEYS.
FLIES is the preferred food of SPIDERS.
PEOPLE is the preferred food of MOSQUITOES.
HONEY is the preferred food of BEARS.
ICE-CREAM is the preferred food of TED.

Example 2 prints the instantiated query for everybody’s meals, using the convenience function, joshua:print-query. It uses the database only, ignoring both rules and questions.

Example 2.
;print just those in the database
[FAVORITE-MEAL MONKEYS BANANAS]
[FAVORITE-MEAL SPIDERS FLIES]
[FAVORITE-MEAL MOSQUITOES PEOPLE]
[FAVORITE-MEAL BEARS HONEY]

Example 3 prints the instantiated query for the meals of bears, using the convenience function, joshua:print-query. It looks in the database and backward rules, but not in questions.

Example 3.
(ask [favorite-meal bears ?food] #'print-query)
;print out bears’ favorite-meal foods
[FAVORITE-MEAL BEARS HONEY]

Example 4 prints the predication object that satisfied the query for everybody’s meals using the accessor function joshua:ask-database-predication. It looks in the database and backward rules, but not in questions. Notice that when the query is satisfied from a rule, the rule name is printed, not a predication object. It is best to use joshua:ask-database-predication with :do-backward-rules nil, that is, when you know the support is only in the database.

Example 4.
(ask [favorite-meal ?eater ?food]
  #'(lambda (backward-support)
      (print (ask-database-predication backward-support)))))
[FAVORITE-MEAL MONKEYS BANANAS]
[FAVORITE-MEAL SPIDERS FLIES]
[FAVORITE-MEAL MOSQUITOES PEOPLE]
[FAVORITE-MEAL BEARS HONEY]
(RULE NOT-FINICKY)

Example 5 prints the instantiated query for everybody’s meals. It uses the database, backward rules, and questions. Note that we supplied just one answer interactively to the question, although we could have supplied more.
Example 5.
; look for backward questions as well

For what values of #EATER and #FOOD is it true that "[GUZZLES #EATER #FOOD]"?
Some solution exists: Yes No
Value for #EATER: CHRISTOPHER
Value for #FOOD: BANANA-PIE
 aborts, #do uses these values

[FAVORITE-MEAL CHRISTOPHER BANANA-PIE]
What are some more values of #EATER and #FOOD such that "[GUZZLES #EATER #FOOD]"?
Some solution exists: Yes No
 aborts, #do uses these values
NIL

[FAVORITE-MEAL MONKEYS BANANAS]
[FAVORITE-MEAL SPIDERS FLIES]
[FAVORITE-MEAL MOSQUITOES PEOPLE]
[FAVORITE-MEAL BEARS HONEY]
[FAVORITE-MEAL TED ICE-CREAM]
[FAVORITE-MEAL CHRISTOPHER BANANA-PIE]

Example 6 collects a list of patterns that describe everybody's meals. It uses the database and rules, but not questions. Note the use of joshuacopy-object-if-necessary. This is because the bindings in the query are undone upon exit from the continuation, so we must make a copy in which to preserve them.

Note that the resulting list is not a list of things that are in the database, but rather a list of free-floating predications that are copies of the query. If you want the latter, use joshua:ask-database-predication with :do-backward-rules nil and don't copy it. See example 7.

Example 6.
(defun collect-answers ()
 (let ((answers nil))
   (ask [favorite-meal ?eater ?food]
     "'(lambda (backward-support)
       (push (copy-object-if-necessary
               (ask-query backward-support)) answers))
     answers))
   COLLECT-ANSWERS
   (collect-answers)
   ([FAVORITE-MEAL TED ICE-CREAM] [FAVORITE-MEAL BEARS HONEY]
    [FAVORITE-MEAL MOSQUITOES PEOPLE]
    [FAVORITE-MEAL SPIDERS FLIES] [FAVORITE-MEAL MONKEYS BANANAS])

Example 7 is identical to example 6, except that here we collect database predications instead of instantiated queries, and the former don't need to be copied. Since we are only looking in the database we specify :do-backward-
rules nil.

(defun collect-answers-database-predications ()
  (let ((answers nil))
    (ask [favorite-meal ?eater ?food]
      #'(lambda (backward-support)
          (push (ask-database-predication backward-support)
               answers)
      :do-backward-rules nil))
   answers))

COLLECT-ANSWERS-DATABASE-PREDICATIONS

(collect-answers-database-predications)

([FAVORITE-MEAL BEARS HONEY]
 [FAVORITE-MEAL MOSQUITOES PEOPLE]
 [FAVORITE-MEAL SPIDERS FLIES] [FAVORITE-MEAL MONKEYS BANANAS])

Better style for the above example would be:

(defun collect-answers-database-predications2 ()
  (let ((answers nil))
    (map-over-database-predications [favorite-meal ?eater ?food]
      #'(lambda (db-predication)
           (push db-predication answers)))
   answers))

Often you're interested in whether there is a solution, but not any particular solution. Example 8 illustrates the use of return-from in a continuation to return when the first solution is found.

Example 8.

(defun solution-exists-p ()
  (ask [favorite-meal ?eater ?food]
      #'(lambda (ignore)
         (return-from solution-exists-p t))
    ;; return nil if nothing succeeded
    nil))

(solution-exists-p)

T

Example 9 is like the example above, but it returns a copy of the query, instead of a boolean. This is useful if you want to know something about the solution, in addition to its existence. (However, if you want to use database-related properties, such as TMS-relation, use joshua:ask-database-predication and don't copy it).
Example 9.
(defun first-solution ()
  (block find-a-solution
    (ask [favorite-meal ?eater ?food]
      #'(lambda (backward-support)
          (return-from find-a-solution
            (copy-object-if-necessary (ask-query backward-support))))
        ;; return nil if nothing succeeded
        nil))

(first-solution)
[FAVORITE-MEAL MONKEYS BANANAS]

**Modeling Note:**
Chances are that you seldom want to define a method that takes over the entire functionality of `joshua:ask`. It's more likely you want to define a method for one of the generic functions it calls, such as `joshua:fetch`, `joshua:ask-data`, `joshua:ask-rules`, `joshua:ask-questions`, or `joshua:map-over-forward-rule-triggers`.

Also, there is a `sys:downward-funarg` declaration on *continuation*, so your implementations of `joshua:ask` should not use *continuation* in other than stack-like ways.

Related Functions:
- `joshua:teLL`
- `joshua:clear`
- `joshua:copy-object-if-necessary`
- `joshua:map-over-database-predications`

See the section "Querying the Database", page Querying-the-Database-SECTION. See the section "The Joshua Database Protocol" in *Joshua Reference Manual*. See the section "Customizing the Data Index" in *Joshua Reference Manual*.

**joshua:ask-database-predication** `backward-support`

Function
An accessor function for use in an `joshua:ask` continuation. It extracts the database predication that matched the query from the continuation argument, `backward-support`, that contains information about the satisfied query.

We describe this continuation argument fully in the dictionary entry for `joshua:ask`.

Note that if the backward support did not come from the database, `joshua:ask-database-predication` gives a bogus answer; in some cases, such as user-written models, it may even cause a trip to the debugger. Thus, you should use `joshua:ask-database-predication` only with `:do-backward-rules nil`.

Examples:
We build a library database using joshuatell statements as well as a forward rule that says the library owns any work authored by Shakespeare. We also include an LTMS in our predicate definitions so that we can later apply joshua:explain to the database predications we find.

(define-predicate author-of (work author) (ltms:ltms-predicate-model))
(define-predicate owns-library (work) (ltms:ltms-predicate-model))

(defrule Shakespeare-holdings (:forward)
  if [author-of ?work Shakespeare]
  then [owns-library ?work])

(defun build-author-title-index2 ()
  (clear)
  (tell [and [author-of "King Lear" Shakespeare]
            [author-of "Hedda Gabler" Ibsen]
            [owns-library "Trumpeting Joshua"]
            [author-of "A Doll’s House" Ibsen]])
  (cp:execute-command "Show Joshua Database")
  BUILD-AUTHOR-TITLE-INDEX2)

(build-author-title-index2)
True things
[OWNS-LIBRARY "Trumpeting Joshua"] [AUTHOR-OF "Hedda Gabler" IBSEN]
[OWNS-LIBRARY "King Lear"] [AUTHOR-OF "King Lear" SHAKESPEARE]
[AUTHOR-OF "A Doll’s House" IBSEN]
False things
None

Now we ask Joshua to find and joshua:explain the database predications that tell what the library owns.

(ask [owns-library ?work]
  #"(lambda (backward-support)
       (explain (ask-database-predication backward-support)))))
[OWNS-LIBRARY "Trumpeting Joshua"] is *True*.
It’s a :Premise.
[OWNS-LIBRARY "King Lear"] is *True*.
It’s derived from the rule Shakespear-Holdings, using:
[AUTHOR-OF "King Lear" SHAKESPEARE]

Usually you can use the convenience function joshua:map-over-database-predications instead of joshua:ask-database-predication.

For comparison we use the same library example for both functions.

For more on these and related functions: See the function joshua:ask, page joshua:ask-FUNCTION.

joshua:ask-derivation backward-support  Function
An accessor function for use in an joshua:ask continuation. It extracts the support information about the satisfied query from the continuation argument backward-support.
Note that the accessor function `joshua:ask-database-predication` makes more assumptions about the support than `joshua:ask-derivation` does.

Here is a schematic representation of the contents of `backward-support`. `joshua:ask-derivation` extracts only the derivation portion. For more detail please consult the dictionary entry for `joshua:ask`.

The backward-support list:

```
(<unified query> <truth-value> . <derivation>)

((unified) query>)

(t/f)
```

(<derivation>) Possibilities for these elements are:

```
(<database predication>)
(AND <conjunct1 derivation> <conjunct2 derivation> ...)
(OR <successful disjunct derivation>)
((RULE <rule name>) <conjunct1 t/f derivation> <conjunct2 t/f derivation> ...)
((QUESTION <question name>) <succeed argument>)
(KNOW <derivation>)
(PROVABLE <derivation>)
```

Like the other accessor functions, `joshua:ask-derivation` does not interpret the information it extracts. Generally you won't need to use it very often.

Note that the convenience functions `joshua:print-query-results` and `joshua:graph-query-results`, respectively, display and graph an annotated version of the support information, so that you don’t have to interpret it yourself.

For comparison we’ll use the same examples to illustrate all three of these functions.

Examples:

The first example shows the support for a query satisfied by database lookup — the database predication that satisfied the query is printed.

```
(define-predicate type-of (object type))

(tell [type-of Iliad epic])

Example 1.
(ask [type-of ?book epic]
    #'(lambda (backward-support)
        (print (ask-derivation backward-support)))
    ([TYPE-OF ILIAD EPIC])
```

The next example shows the support for a query that is satisfied from rules. We have a rule, `dessert?`, that determines if a given food is a dessert. Each of this rule’s subgoals is derived from other rules. Here are
the definitions.

; Example 2. Query is derived from backward rules
; Define the predicates
(define-predicate edible (object))
(define-predicate is-food (object))
(define-predicate contains (object substance))
(define-predicate sweet (object))

; Define the rules
(defrule food? (:backward)
  if [edible ?object]
  then [is-food ?object])

(defrule sweet? (:backward)
  if [or [contains ?object chocolate]
       [contains ?object sugar]
       [contains ?object honey]]
  then [sweet ?object])

(defrule dessert? (:backward)
  if [and [is-food ?object]
         [sweet ?object]]
  then [type-of ?object dessert])

; tell some sticky facts
(tell [edible chocolate-coated-ants])
(tell [contains chocolate-coated-ants honey])

Now we 
(joshua:ask what foods qualify as desserts and why. The display is
a list starting with rule dessert? that satisfied the query; next is the first
subgoal that was satisfied, together with its truth value, and the name of
the rule which satisfied it (rule food?). That rule's first subgoal is then
listed with its truth value and the database predication that satisfied it,
and so on, through all the backward support.

(ask [type-of ?object dessert]
  #'(lambda (backward-support)
     (print (ask-derivation backward-support))))
(RULE DESSERT?)
  ([IS-FOOD CHOCOLATE-COATED-ANTS] 1 (RULE FOOD?)
   ([EDIBLE CHOCOLATE-COATED-ANTS] 1 [EDIBLE CHOCOLATE-COATED-ANTS]))
  ([SWEET CHOCOLATE-COATED-ANTS] 1 (RULE SWEET?)
   ([CONAINS CHOCOLATE-COATED-ANTS HONEY] 1
    [CONAINS CHOCOLATE-COATED-ANTS HONEY]]))

For more on these and related functions: See the function 
(joshua:ask, page
joshua:ask-FUNCTION.
**joshua:ask-query** backward-support  
*Function*  
An accessor function for use inside an **joshua:ask** continuation. It extracts the instantiated query from the continuation argument backward-support.

backward-support is fully described in the dictionary entry for **joshua:ask**.

Example:

Here we collect and save all the answers from a query. (See example 6 in the dictionary entry for **joshua:ask**.)

```lisp
(defun collect-answers ()
  (let ((answers nil))
    (ask [favorite-meal ?eater ?food]
      #(lambda (backward-support)
          (push (copy-object-if-necessary
                  (ask-query backward-support))
               answers)))
  answers))
```

To extract and print out the instantiated query, use the convenience function **joshua:print-query**.

For more on these and related functions: See the function **joshua:ask**, page joshua:ask-FUNCTION.

**joshua:ask-query-truth-value** backward-support  
*Function*  
An accessor function for use inside an **joshua:ask** continuation. It extracts the truth value of the instantiated query from the continuation argument backward-support.

backward-support is fully described in the dictionary entry for **joshua:ask**.

The truth value is a number, as follows:

- **0** Truth value of **joshua:*unknown*  
- **1** Truth value of **joshua:*true*  
- **2** Truth value of **joshua:*false*  
- **3** Truth value of **joshua:*contradictory*  

The **joshua:truth-value** presentation type translates these numbers into symbols naming a truth value.

Most of the time you know the query’s truth value from the query pattern itself, so that you have little need of this function. The truth value information is mostly there for system use, to let the system interpret the query.

Examples:

```lisp
(define-predicate status-of (object status))
(tell [status-of smoke-alarm off])
```
; Example 1.
(ask [status-of ?indicator off]
  #'(lambda (backward-support)
      (print (ask-query-truth-value backward-support))))
1

; Example 2. Use truth-value-name to translate the number
(ask [status-of ?indicator off]
  #'(lambda (backward-support)
      (print (truth-value-name (ask-query-truth-value backward-support))))
TRUE

For more on this and related functions: See the function joshua:ask, page joshua:ask-FUNCTION.

joshuaclear &optional (clear-database t) (undefrule-rules nil)
Function
With arguments t t, empties the database and "undoes" all rule definitions.

clear-database Specifies whether or not to clear the database. Default is t.

undefrule-rules Specifies whether or not to delete all rule definitions. Default is nil.

Clearing the database is equivalent to joshua:untelling each fact in the database.

Note that undefining all rule definitions is a drastic thing to do, as it clears out all rules in your world. Any application depending on these rules will no longer work. Clear out all rules only if you want a "clean" environment, for example, if you need to get rid of a runaway rule that you cannot stop by other means.

Examples:

Show Joshua Database
True things
[FAVORITE-MEAL BEARS HONEY]
[FAVORITE-MEAL MOSQUITOES PEOPLE]
[FAVORITE-MEAL SPIDERS FLIES]
[FAVORITE-MEAL MONKEYS BANANAS]
False things
None

(clear)
Show Joshua Database
True things
None
False things
None

Related Command:
"Clear Joshua Database Command"

See the section "Removing Predications From the Database", page Removing-Predications-From-the-Database-SECTION.


See the section "Customizing the Data Index" in Joshua Reference Manual.

Clear Joshua Database Command

Clears predications from the Joshua Database.

Predications Which predications to remove from the database. Clear Joshua Database asks the database for all predications matching those specified in the Predications argument and joshua:untells them from the database. The value of Predications can also be All or None.

:Other Truth Values Too
Whether or not to clear the predications in the database which match those specified by the Predications argument, but have the opposite truth value. This argument defaults to Yes.

:Query Whether to ask you before making changes to the database. By default, the command stops and asks before removing any predications or rules.

:Undefine Rules If Undefine Rules is Yes, the command will undefine all of the Joshua Rules. This argument defaults to No.

:Verbose Whether to print information about what the command is doing.

Clear Joshua Database provides a convenient interface to the joshua:untell function. It asks the database for all predications matching those specified by the arguments, prompts you for confirmation, and joshua:untells each predicate. It also allows you to undefine all the Joshua rules, resulting in a fresh Joshua environment.

Note that undefining all rule definitions is a drastic thing to do, as it clears out all rules in your world. Any application depending on these rules will no longer work. Clear out all rules only if you want a "clean" environment, for example, if you need to get rid of a runaway rule that you cannot stop by other means.
Related Functions:
   joshua:clear
   joshua:untell

**joshua:contradictory**

A named constant used by Joshua to denote an interim state of computation wherein a predication is believed to be both joshua:*true* and joshua:*false*. When this occurs, Joshua invokes the appropriate Truth Maintenance System to resolve the contradictory state.

joshua:*contradictory* is not meaningful unless a TMS is present. However, not all Truth Maintenance Systems are required to use this value.

Related Topics:
   joshua:*true*
   joshua:*false*
   joshua:*unknown*
   joshua:truth-value
   joshua:predication-truth-value

See the section "Truth Values", page Truth-Values-SECTION. See the section "Justification and Truth Maintenance", page Justification-and-Truth-Maintenance-SECTION.

**joshua:copy-object-if-necessary object**

Copies the object handed to it if it contains variables, or is otherwise ephemeral.

object

Any object, for example, a list, or a predication

Variables in object are renamed during copying, so that variables in the copy differ from variables in the original.

joshua:copy-object-if-necessary is useful for making copies of predications that may be stack-consed, or whose variables may be temporarily unified. The latter, for example, is true of variables in the query to joshua:ask.

joshua:copy-object-if-necessary creates a separate copy of its argument in the heap.

Examples: Here we reuse some of the examples introduced with joshua:ask. We define some predicates and a rule, then enter some facts into the database.

(define-predicate favorite-meal (eater food))
(define-predicate guzzles (eater food))
(clear)
(tell [and [favorite-meal bears honey]
[favorite-meal mosquitoes people]
[favorite-meal spiders flies]
[favorite-meal monkeys bananas]
[guzzles ted ice-cream]])

Show Joshua Database
True things
[FAVORITE-MEAL BEARS HONEY]
[FAVORITE-MEAL MOSQUITOES PEOPLE]
[FAVORITE-MEAL SPIDERS FLIES]
[FAVORITE-MEAL MONKEYS BANANAS]
[GUZZLES TED ICE-CREAM]
False things
None

(defrule not-finicky (:backward)
  if [guzzles ?eater ?food]
  then [favorite-meal ?eater ?food])

Example 1.
;;;If you don’t copy the query, you lose the information!
(defun collect-answers-wrong ()
  (let ((answers nil))
    (ask [favorite-meal ?eater ?food]
      #'(lambda (backward-support)
          (push (ask-query backward-support) answers)))
    answers))
COLLECT-ANSWERS-WRONG

(collect-answers-wrong)
#<Error printing object CONS 42353464>

Example 2.
;;;Using copy-object-if-necessary saves the information
(defun collect-answers ()
  (let ((answers nil))
    (ask [favorite-meal ?eater ?food]
      #'(lambda (backward-support)
          (push (copy-object-if-necessary
                  (ask-query backward-support)) answers)))
    answers))
COLLECT-ANSWERS
(collect-answers)

([FAVORITE-MEAL TED ICE-CREAM] [FAVORITE-MEAL BEARS HONEY]
[FAVORITE-MEAL MOSQUITOES PEOPLE]
[FAVORITE-MEAL SPIDERS FLIES] [FAVORITE-MEAL MONKEYS BANANAS])

(defun first-solution ()
  (block find-a-solution
    (ask [favorite-meal ?eater ?food]
      #'(lambda (backward-support)
          (return-from find-a-solution
            (copy-object-if-necessary (ask-query backward-support))))
    ;; return nil if nothing succeeded
    nil))
FIRST-SOLUTION

(first-solution)
[FAVORITE-MEAL MONKEYS BANANAS]

Related Functions:
joshua:ask

joshua:define-predicate name args &optional (model-and-other-components *default-predicate-model) &body options

Defines a predicate for use in building predications.

name Any symbol that does not conflict with the name of an existing flavor or presentation type. So, for example, integer, cons, and array are not good predicate names. In fact, they can be disastrous. Doing joshua:define-predicate on these will likely cause problems in both the CL type system and the presentation system.

args A list of symbols, similar to Lisp lambda lists. &optional arguments can be defaulted as in Lisp. Note that, unlike Lisp, &rest arguments can also be defaulted. &rest arguments can be used in "tail" fashion, as in: [foo A B . ?quux], which matches all foo predicates with arguments A and B, followed by anything else. &key, &aux, and other lambda-list keywords are not supported.

model-and-other-components
Lists optional models defined with joshua:define-predicate-model. You can also use any flavor, as long as it doesn't use ordered-instance-variables. The rules of procedure are identical to those of defflavor.
options Any option acceptable to defflavor. :constructor is unlikely to be useful, as joshuadef-predicate already uses it. In addition, see :destruct-into-instance-variables, below.

There are two ways that you can make the predicate arguments lexically available to methods. For frequent use, specify the option :destruct-into-instance-variables in your predicate definition. This keeps the predicate arguments destructured permanently in each predication, taking up more space but providing faster access. For occasional use you can call the macro joshua:with-statement-destructured. Since the macro destructures the arguments each time you call it, it is slower, but such predications take up less space. The latter, for example, is usually appropriate for joshuasay methods. The former might be more appropriate for inner loops.

Examples:

(define-predicate fruit (a-fruit))

(define-predicate bird (bird) (ltms:ltms-predicate-model))

(define-predicate things-to-pack (traveller &rest objects))

(define-predicate gun (range calibre)
  :destruct-into-instance-variables)

(define-predicate has-disease (patient disease &rest symptoms)
  (:destruct-into-instance-variables disease)) ; partial destructuring

Related Functions:

joshua:undef-predicate
joshuamake-predication
joshuapredication

Related Flavor:

joshuapredication

See the section "Joshua Predications", page Joshua-Predications-SECTION.

joshua:defquestion name (control-structure &rest control-
structure-args) pattern &key :code

Macro

Defines a question.

name The name of the question.

control-structure Specifies the direction of chaining the question responds to. Currently, only :backward chaining questions are supported.

control-structure-args Like joshua:defrule, these are arguments to the control structure. Currently supported are :importance and
:documentation. Both work as they do in rules: The former lets you specify the priority in which you want your questions to run (however, they'll always run after rules); the latter lets you add a string to document the meaning of the question. This string can then be retrieved with the Lisp function joshua:documentation.

pattern A single predication. The question triggers when this pattern is matched in an joshua:ask, for :backward question.

Keywords:

:code Any Lisp code. This is for customized versions of joshua:defquestion.

Backward questions behave like backward chaining rules, except that they run after all backward rules. They treat the user as an extension of the database, and solicit more solutions from him. (For the basics of rule operation: See the section "Rules and Inference", page Rules-and-Inference-SECTION.)

Like rules, questions have a name, a trigger pattern, and a body. Like rules, questions are a way of generating information.

When you joshua:ask something with :do-questions joshua::t and the query pattern unifies with pattern in the question, the question body runs. Questions run only after the database has been searched and all appropriate backward rules have been triggered.

If you don't supply the :code keyword, joshua:defquestion supplies a body for you.

At run time, the query unifies with the question trigger. If there are no logic no logic variables in the unified query, a Yes or No question is generated once. The default answer is No. Answering Yes makes the query that triggered the question succeed. Answering No makes the query fail, which can mean either that the query is known to be joshua:*false*, or that it is not known to be joshua:*true*.

If the unified query contains logic variables, the question loops, presenting iterations of an AVV (Accept Variable Values) menu, each soliciting bindings for those variables.

Questions can be used to interact with a user, with some other process running on the machine, or even some other device. For example, a question could go out over the network and ask some other device to answer a question.

Joshua has a default way of asking questions; you can also write your own.

The default version uses either the default joshua:say method to format pattern or a user-defined joshua:say method if available.
Examples:
We define a predicate and then we define a question that triggers on a preication pattern built from this predicate.

\[
\text{(define-predicate \text{foo} (something something-else))}
\]

\[
\text{(defquestion \text{question1} (:backward :documentation "This has no apparent use")}
\]  
\[
\text{[\text{foo 1 ?x}])}
\]

Example 1 is a query with no logic variables in the unified query pattern.

\[
\text{Example 1:}
\]  
\[
\text{(ask \text{foo 1 2} #'\text{print-query} :do-questions t)}
\]  
\[
\text{Is it true that "[\text{FOO 1 2}]"? [default No]: Yes}
\]  
\[
\text{[\text{FOO 1 2]}
\]
\[
\text{NIL}
\]

For example 2 we define a \text{joshua:say} method, and the question uses that method.

\[
\text{Example 2:}
\]  
\[
\text{(define-predicate-method \text{say \text{foo}} (optional (stream *standard-output*)))}
\]  
\[
\text{(with-statement-destructured (something something-else) (}
\]  
\[
\text{(format stream "the arguments -A and -A are correct"}
\]  
\[
\text{something something-else)))}
\]

\[
\text{(ask \text{foo 1 2} #'\text{print-query} :do-questions t)}
\]  
\[
\text{Is it true that "the arguments 1 and 2 are correct"? [default No]: Yes}
\]  
\[
\text{[\text{FOO 1 2]}
\]
\[
\text{NIL}
\]

Example 3 uses a query with logic variables in the query pattern.

\[
\text{Example 3:}
\]  
\[
\text{(ask \text{foo 1 =x} #'\text{print-query} :do-questions t)}
\]  
\[
\text{For what values of =x is it true that "the arguments 1 and =x are correct"?}
\]  
\[
\text{Some solution exists: Yes No}
\]  
\[
\text{Value for =x: 2}
\]  
\[
\text{<ERROR> aborts, <OK> uses these values}
\]
\[
\text{[\text{FOO 1 2]}
\]
\[
\text{What are some more values of =x such that "the arguments 1 and =x are correct"?}
\]  
\[
\text{Some solution exists: Yes No}
\]  
\[
\text{Value for =x: BAR}
\]  
\[
\text{<ERROR> aborts, <OK> uses these values}
\]
\[
\text{[\text{FOO 1 BAR]}
\]
\[
\text{What are some more values of =x such that "the arguments 1 and =x are correct"?}
\]  
\[
\text{Some solution exists: Yes No}
\]  
\[
\text{<ERROR> aborts, <OK> uses these values}
\]
\[
\text{NIL}
\]

To write your own code to do questions, use the \text{code} keyword. This keyword takes arguments and a body, as follows:

\text{args}  \quad \text{(query truth-value continuation &optional query-context)}
The body of a `joshua:defquestion` works like Lisp code in the body of a backward rule. If the value of `body` is `nil`, the query that triggered the question fails. If the value of `body` is non-`nil`, the query succeeds. Calling the `joshua:succeed` function explicitly within the `body` allows the query to succeed many times.

Within `body`, `query` is the query predication given to `joshua:ask`, after the query has been unified with the question’s trigger.

If `truth-value` is `joshua:*true*`, Joshua is trying to determine whether the query is known to be true, as opposed to false or unknown. Similarly for a `truth-value` of `joshua:*false*` Joshua tries to determine whether the query is known to be false, as opposed to true or unknown.

The `query-context` argument can almost always be ignored.

`body` should do the following:

- If there are no logic variables in the query, decide somehow (perhaps by asking the user a question) if the query is true. If so, call `continuation`. You usually rely on the form `(joshua:succeed)` to call `continuation` for you.

- If there are logic variables present, solicit sets of bindings for them from somewhere (for example, the user). For each such set, call `continuation` (usually via `(joshua:succeed)`).

Examples of custom-written questions:

First we define the predicates, a `joshua:say` method, a question, and a backward rule.

```lisp
(define-predicate wrote (author book))
(define-predicate understands (reader book))

(define-predicate-method (say understands)
  (optional (stream *standard-output*)))
(with-statement-destructured (reader book) self
  (format stream ""A understands "A." reader book)))
```
(defquestion writings-of-caesar (:backward) [wrote caesar ?book]
  :code
  ((query truth-value continuation &optional ignore)
   (unless (eq truth-value *true*
     (error "I can only ask positive questions.")))
   (typecase ?book
     (unbound-logic-variable
      ;; asked with ?book unbound
      (loop for prompt = "Tell me something that Caesar wrote: "$
        then "Tell me something else Caesar wrote: "$
        for answer = (accept
          '(((token-or-type ("No more" . no-more))
              (string)))
          :prompt prompt :default "De Bello Gallico")
        until (eq answer 'no-more)
        do (with-unification
          (unify ?book answer)
          (succeed)))
     (otherwise
      ;; asked with ?book bound
      (yes-or-no-p "&Did Caesar write "A? " ?book))))))
(defrule writers-understand-their-work (:backward)
  if [wrote ?author ?work]
  then [understands ?author ?work])

Now we joshua:ask the query.

Tell me something that Caesar wrote: [default "De Bello Gallico"]:
De Bello Gallico
CAESAR understands De Bello Gallico.
Tell me something else Caesar wrote: [default "De Bello Gallico"]:
A Canticle for Leibowitz
CAESAR understands A Canticle for Leibowitz.
Tell me something else Caesar wrote: [default "De Bello Gallico"]: No more NIL

(ask [understands Caesar "Passion on the Nile"] #'say-query :do-questions t)
Did Caesar write Passion on the Nile? (Yes or No) Yes
CAESAR understands Passion on the Nile.
NIL

Related Functions:
  joshua:undefquestion
  joshua:ask
  joshua:ask-questions
  joshua:map-over-backward-question-triggers
  joshua:locate-backward-question-trigger
See the section "Asking the User Questions", page Asking-the-User-Questions-SECTION.

**joshua:defrule**  
`rule-name (control-structure &rest control-structure-args) if if-part then then-part`

Defines a forward or backward chaining rule. The `control-structure` argument specifies the direction of the rule.

Forward chaining rules respond to new facts entered with `joshua:telI`; the response (that is, the rule body or `then-part`), can involve deducing additional facts that are automatically added to the database, or it can involve executing any Lisp program.

Backward chaining rules respond to a goal entered with `joshua:ask` by trying to satisfy it; this can involve satisfying a series of successive subgoals, or any Lisp program. Backward chaining does not automatically add new facts to the database. See the section "Rules and Inference", page Rules-and-Inference-SECTION.

*rule-name*  
Any symbol that uniquely identifies the rule.

*control-structure*  
One of the keywords `:forward` or `:backward` corresponding, respectively, to a forward rule or a backward rule. Future releases may add more possible control structures.

*control-structure-args*  
- **:importance** lets you control the order of rule execution. **:documentation** lets you add a string that documents the meaning of the rule. Future releases may add more keywords.

- **:importance** takes a `value` argument that can be:

  - Numeric; any non-complex number, including `+1e∞` or `-1e∞` (infinity).

  - A symbol (in which case, the system treats it as a special variable whose runtime value should be a number).

  - A form; the compiler enwraps it with `(lambda () ...)` and compiles it. It should return a number when called.

The larger the `value` argument, the higher the priority. Rules with no `value` argument run first, after which rules with a `value` argument are run in order from the highest to the lowest `value`. 
Some expense is associated with ordering using 
:importance. In forward chaining rules it causes a
"best-first" search through a heap of rules according
to the value associated with :importance. Backward
chaining only orders the local "best-first" search of
rules at the current choice point.

A more symbolic type of reasoning, or some level of
modeling are usually preferable to the indiscriminate
use of :importance.

if

The symbol joshua::if.

if-part

Specifies the conditions under which the rule suc-
cceeds. The form of the if-part is identical for forward
and backward rules. Procedurally, the if-parts differ
depending on rule type:

In forward rules the if-part is the trigger part. It can
be one or more predications, joined by joshua::and or
joshua::or. Lisp forms (called procedural nodes) can
be included in the if-part of forward rules, as well.
See the section "The Joshua Rule Compiler" in Joshua

In backward rules the if-part is the action part. It can
be one or more predications as above, as well as any
Lisp construct. These become subgoals.

then

The symbol joshua::then.

then-part

Specifies the conclusions drawn from the rule. The
form of the then-part is identical in forward and back-
ward rules. Procedurally, the then-parts differ depend-
ing on rule type:

In forward rules the then-part is the action part. Can
be one or more predications, joined by joshua::and or
joshua::or, as well as any Lisp construct.

In backward rules this is the trigger part. Must be a
single (not a compound) predication.

Note that the if and then clauses can occur in either order. For example,
some programmers prefer to place the then-part of backward rules first, so
that the trigger (procedure head) always comes first. Either of the arrange-
ments shown below is valid.

If [...] Then [...] and

Then [...] If [...]
A rule's action part (the then-part of forward rules, and the if-part of backward rules) can specify any suitable action(s), such as adding or retracting predications, using Lisp code to perform embedded tests or computations, calling joshua:ask or joshuatell, interacting with the user, or displaying messages. When your Lisp code does iterations, call the function joshuasucceed inside it to let Joshua know that the current set of bindings is correct. Otherwise, Lisp code "succeeds" by returning non-nil. See examples below.

If the action part of a forward rule contains a predication that is not embedded in Lisp code, this newly deduced fact is automatically added to the database when the rule executes (a joshuatell is implicit). Note that such a predication can be backquoted. If the predication is embedded in Lisp, however, you must explicitly use a joshuatell to insert the fact into the database.

The action part of a backward rule has an implicit joshua:ask around it. Backward rule action parts add no predications to the database, unless you explicitly use a joshuatell to accomplish this.

A backward rule's trigger part (the then-part) must consist of a single predication. The trigger can contain logic variables. These variables are bound by the unifier when the trigger part of the rule is matched against the query; they are then passed to the action part (the if-part).

A forward rule’s trigger part (the if-part) may contain an arbitrary number of predications and Lisp forms. The triggers can contain logic variables. A forward rule’s triggers behave as follow:

- If the trigger is a predication, it is satisfied when it has been matched against a predication in the database. The logic variables in the trigger are bound by the unifier when the trigger part of the rule is matched against the database predication.

- The trigger may be a Lisp form (we call such triggers procedural triggers). Such a trigger may be satisfied in two ways: If it returns joshua:?, it is regarded as satisfied. It is also regarded as satisfied each time it calls joshuasucceed.

- If a procedural trigger never calls joshuasucceed, but merely returns joshua: or joshua:nil, then it acts as a filter on the previous triggers (either accepting or rejecting the bindings produced by its predecessors).

- A procedural trigger may also act as a generator, producing several acceptable sets of bindings and calling joshuasucceed for each one.

- Logic variables which occur for the first time in a procedural trigger may be bound by calling joshuaunderify. Logic variables that are referenced in a procedural trigger but which occur in an earlier trigger, are bound to the value established by the earlier trigger during the execution of the Lisp trigger.
• The logical connective *and* can be used to group the triggers into subgroups all of which must be satisfied. The logical connective *or* can be used to group the patterns into subgroups any one of which must be satisfied.

• The trigger part of a forward rule can include the keyword *support* followed by a logic variable after any trigger pattern. During the execution of the rule, this logic variable is bound to the predication that matched the trigger pattern immediately preceding the keyword *support*.

• A procedural trigger may provide an argument to joshua:succeed which should be a database-predication. If it does so, this predication is treated as if it had matched a normal trigger of the rule. If there is a *support* keyword following the procedural trigger, then the logic variable following it will be bound to the database-predication.

Joshua stores and retrieves rules by their triggers. When a new rule is defined, the rule compiler stores the rule’s trigger in a place appropriate to the rule type. The system finds and executes applicable rules by locating their triggers; similarly, it deletes unwanted rules by removing their triggers. See the section “The Joshua Rule Indexing Protocol” in *Joshua Reference Manual*.

Here are some examples. First, here’s how to use the *documentation* keyword. We use a forward rule as an example, but *documentation* works identically for backward rules.

```lisp
(define-predicate reads (person how-much))
(define-predicate is-bookworm (person))

(defrule simple (:forward :documentation "Identifies bookworms")
  if [reads ?person constantly]
  then [is-bookworm ?person])
```

To retrieve the documentation string of this rule, use the Lisp function joshua::documentation.

```lisp
(documentation 'simple)
"Identifies bookworms"
```

Here are some examples of forward chaining. This first a simple declarative rule:

```lisp
(defrule good-cake (:forward)
  if [and [rises ?cake justright] [color ?cake evenly-gold] [texture ?cake moist] [taste ?cake justright]]
  then [good ?cake])
```

Next is an example of using the *support* keyword to allow the body of the rule to reference the triggering facts:
(defrule good-cake (:forward)
  if [and [rises ?cake justright] :support ?f1
      [color ?cake evenly-gold] :support ?f2
      [texture ?cake moist] :support ?f3
  then [and (Format t "-%The reason I thing that ~s is good is that:" ?cake)
           (say ?f1) (say ?f2) (say ?f3) (say ?f4)
           [good ?cake]])

Here we show how a Procedural Trigger can be used as a generator. Once all triggers before the procedural trigger are matched, it executes and generates two acceptable bindings for ?color.

(defrule good-cake (:forward)
  if [and [rises ?cake justright] [texture ?cake moist]
       (loop for color in '(evenly-gold nicely-brown)
             do (unify ?color color)
             (succeed))
       [taste ?cake justright]]
  then [and (format t "-%~s is a good cake with color ~s" ?cake ?color)
           [good ?cake]])

Here is an example of a procedural trigger being used as a filter:

(defrule check-temperature (:forward)
  if [and [temperature-used ?object ?temp]
       (< 325 ?temp 400)] ; example of Lisp used as a filter
  then [correct-temperature-used ?object ?temp])

(defun check-oven-setting ()
  (clear)
  (tell [temperature-used jelly-roll 375])
  (ask [correct-temperature-used jelly-roll ?temp] #'print-query))

(check-oven-setting)
[CORRECT-TEMPERATURE-USED JELLY-ROLL 375]
NIL

Finally, here is an example using nested and's and or's:

(defrule deduce-ancestry (:forward)
  if [or [is-parent-of ?old ?young]
       [and [is-ancestor-of ?old ?middle]
             [is-parent-of ?middle ?young]]]
  then [is-ancestor-of ?old ?young])

Here are some examples using backward chaining:
(defrule sailor-alert (:backward)
  if [or [condition-of wind gusting]
      [weather-forecast squalls]]
  then [issue-warning small-craft alert])

;;; Lisp code in action part of backward rule
(define-predicate good-to-read (book))
(defparameter *books* '(decameron canterbury- TALES gargantua- and- pantagruel
tom-jones catch-22))

(defrule reading-list (:backward)
  if (typecase ?candidate-book
    (unbound-logic-variable
     (loop for book in *books*
          doing (with-unification
                   (succeed)))))
  (otherwise
   (member ?candidate-book *books*)))))
  then [good-to-read ?candidate-book])

(ask [good-to-read ?x] #'print-query)
[GOOD-TO-READ DECAMERON]
[GOOD-TO-READ CANTERBURY- TALES]
[GOOD-TO-READ GARGANTUA- AND- PANTAGRUEL]
[GOOD-TO-READ TOM-JONES]
[GOOD-TO-READ CATCH-22]
NIL

You can inhibit backward chaining rule invocation by passing joshua::nil as the
:do-backward-rules keyword argument to joshua:ask (the default value is
joshua::t). In this case the system does only a database lookup.

You can cause backward question invocation by passing joshua::t as the
:do-questions keyword argument to joshua:ask (the default is joshua::nil).
Advanced Concepts Note:
Six built-in models are available for predicates in joshua:ask goals. These flavors do subsets of what joshua:ask normally does, by leaving out one or more of the steps joshua:ask-data, joshua:ask-rules, or joshua:ask-questions. Thus the models save a certain amount of overhead when their predicates are used as goals to joshua:ask. The steps that are done are indicated by the names:

- joshua:ask-data-only-mixin
- joshua:ask-rules-only-mixin
- joshua:ask-questions-only-mixin
- joshua:ask-data-and-rules-only-mixin
- joshua:ask-data-and-questions-only-mixin
- joshua:ask-rules-and-questions-only-mixin

Related Functions:

joshua:undefrule  
joshuatell  
joshua:ask  
joshua:ask-rules


\textbf{joshua:different-objects object1 object2}

\textit{Function}

Returns \texttt{nil} if the arguments are \texttt{eql} or if either argument is an uninstantiated logic variable (in the latter case the two objects can potentially be \textit{made} to be the same by the unifier). Otherwise, \texttt{joshua:different-objects} returns \texttt{t}.

\texttt{object1} A Lisp object.  
\texttt{object2} A Lisp object.

This function is useful in making rules that weed out inappropriate self-referential behavior. For example, in a program simulating the behavior of a monkey that can pick up objects, it is important to ensure that the monkey does not try to pick up itself.

This function is often useful in the if-part of rules, or in Lisp code.

(defrule pick-up (:backward)  
if (different-objects ?obj 'monkey)  
thenn [can-pick-up monkey ?obj])

To invoke this rule, you would type something like:

(ask [can-pick-up monkey wrench] #'print-query)

See the section "Using Joshua Within Lisp Code", page Using-Joshua-Within-Lisp-Code-SECTION.
Disable Joshua Tracing Command

Turns off Joshua tracing.

Type of Tracing  The type of tracing to disable. It can be one of forward rules, backward rules, predications, TMS operations, or all. The type of tracing defaults to all.

Disable Joshua Tracing turns off the Joshua tracing facility.

Related Commands:

"Enable Joshua Tracing Command"
"Reset Joshua Tracing Command"

Enable Joshua Tracing Command

Turns on Joshua Tracing.

Type of Tracing  The type of tracing to enable. You can enable the tracing of forward rules, backward rules, predications, TMS operations, or All. Unless otherwise specified (by using the :Menu option for example), tracing is turned on with the same options and tracing events that were in effect the last time you used tracing.

:Menu  Brings up a menu of detailed tracing options for the type of tracing being enabled. This menu provides a greater degree of control over exactly what gets traced and when the tracing facility interacts with the user.

:Trace Events  When enabling a particular type of tracing this option allows you to specify precisely which events will be displayed during tracing. These can also be set by using the :Menu option.

:Step Events  Allows you to specify at which events the tracing facility will stop and prompt for interaction. These can also be set by using the :Menu option.

The Enable Joshua Tracing command turns on the Joshua tracing tools and allows you to customize tracing to your particular application or preference. The Joshua tracing facility is very flexible. You can, for example, trace just forward rules that are triggered by predications matching a particular pattern:

Enable Joshua Tracing Forward Rules :Menu Yes

Forward Rules

Forward Rules Options
Trace forward rules: All  Selectively
Trace forward rules: None  forward rules
Trace forward rule triggers: None  [TME:LOVES DEMETRIUS HERMIA]
Traced Events: Fire  Exit  Queue  Dequeue
Stepped Events: Fire  Exit  Queue  Dequeue
ódigo aborts, IDO uses these values
Or, you can even just trace predications built on a particular model:

Enable Joshua Tracing Predications : Menu Yes

**Predications**

**Predications Options**
- Trace Predications: All **Selectively**
- Trace Predicates of flavor(s): None **LTMS:LTMS-PREDICATE-MODEL**
- Trace Facts Matching: **None** predicates
- Traced Events: Ask Tell Untell Truth value change Justify Unjustify
- Stepped Events: Ask Tell Untell Truth value change Justify Unjustify

The best way to familiarize yourself with this facility is to type Enable Joshua Tracing All :Menu Yes. This brings up a menu of all the types of Joshua tracing and the options available for each one. By moving the mouse over each option you can see the documentation for that option in the mouse documentation line.

Related Commands:

"Disable Joshua Tracing Command"
"Reset Joshua Tracing Command"

See the section "Tracing Predications", page Tracing-Predications-SECTION. See the section "Tracing Rules", page Tracing-Rules-SECTION.

**Explain Predication Command**

Traces the chain of TMS justifications for *database-predication* through rules to primitive support (premises and assumptions).

*database-predication* A predication object that is in the database. Must be the actual database object, and not a copy of it.

*depth* Specifies how many layers deep into the explanation to go before cutting off.

This is a command interface to Joshua's *joshua:explain* function.

**joshua:explain** *database-predication* &optional *depth* (stream *standard-output*)

Traces the chain of TMS justifications for *database-predication* through rules to primitive support (premises and assumptions).

*database-predication* A predication object that is in the database. Must be the actual database object, and not a copy of it.

*depth* Specifies how many layers deep into the explanation to go before cutting off.

*stream* Specifies a stream to which to display the output.

In general, *joshua:explain* is useful only if *database-predication* is built on some model that supports the TMS protocol.
Examples:

(define-predicate higher-in-food-chain (eater lower-in-food-chain)
  (ltms:ltms-predicate-model))

(define-predicate favorite-meal (eater food) (ltms:ltms-predicate-model))

; A good example of how to implement transitive relations

(defrule basic-food-chain (:forward)
  if [favorite-meal ?eater ?eatee]
  then [higher-in-food-chain ?eater ?eatee])

(defrule transitive-food-chain (:forward)
  if [and [favorite-meal ?eater ?eatee]
       [higher-in-food-chain ?eatee ?food]]
  then [higher-in-food-chain ?eater ?food])

(defun meals ()
  (clear)
  (tell [and [favorite-meal red-herring worm]
           [favorite-meal worm algae]])
  (tell [favorite-meal Miss-Marple red-herring] :justification :assumption)
  (cp:execute-command "Show Joshua Database")

(meals)
True things
[HIGHER-IN-FOOD-CHAIN MISS-MARPLE RED-HERRING]
[HIGHER-IN-FOOD-CHAIN MISS-MARPLE WORM]
[HIGHER-IN-FOOD-CHAIN MISS-MARPLE ALGAE]
[HIGHER-IN-FOOD-CHAIN WORM ALGAE]
[HIGHER-IN-FOOD-CHAIN RED-HERRING ALGAE]
[HIGHER-IN-FOOD-CHAIN RED-HERRING WORM]
[FAVORITE-MEAL MISS-MARPLE RED-HERRING]
[FAVORITE-MEAL WORM ALGAE]
[FAVORITE-MEAL RED-HERRING WORM]
False things
None
(ask [higher-in-food-chain Miss-Marple ?food]
  #'(lambda (backward-support)
    (explain (ask-database-predication backward-support))))
[HIGHER-IN-FOOD-CHAIN MISS-MARPLE RED-HERRING] is true
  It was derived from rule BASIC-FOOD-CHAIN
[FAVORITE-MEAL MISS-MARPLE RED-HERRING] is true
  It is an :ASSUMPTION
[HIGHER-IN-FOOD-CHAIN MISS-MARPLE WORM] is true
  It was derived from rule TRANSITIVE-FOOD-CHAIN
[FAVORITE-MEAL MISS-MARPLE RED-HERRING] is true
  It is an :ASSUMPTION
[HIGHER-IN-FOOD-CHAIN RED-HERRING WORM] is true
  It was derived from rule BASIC-FOOD-CHAIN
[FAVORITE-MEAL RED-HERRING WORM] is true
  It is a :PREMISE
[HIGHER-IN-FOOD-CHAIN MISS-MARPLE ALGAE] is true
  It was derived from rule TRANSITIVE-FOOD-CHAIN
[FAVORITE-MEAL MISS-MARPLE RED-HERRING] is true
  It is an :ASSUMPTION
[HIGHER-IN-FOOD-CHAIN RED-HERRING ALGAE] is true
  It was derived from rule TRANSITIVE-FOOD-CHAIN
[FAVORITE-MEAL RED-HERRING WORM] is true
  It is a :PREMISE
[HIGHER-IN-FOOD-CHAIN WORM ALGAE] is true
  It was derived from rule BASIC-FOOD-CHAIN
[FAVORITE-MEAL WORM ALGAE] is true
  It is a :PREMISE

Related Functions:
  joshua:graph-tms-support

See the section "Explaining Program Beliefs", page Explaining-Program-
Beliefs-SECTION.

joshua:*false*

A named constant used by Joshua to denote a truth value of false. You can
compare truth values using eql.

Related Topics:
  joshua:*true*
  joshua:*unknown*
  joshua:*contradictory*
  joshua:truth-value
  joshua:predication-truth-value

See the section "Truth Values", page Truth-Values-SECTION.
**joshua:graph-query-results** `backward-support &key (orientation :vertical) (stream *standard-output*)`

A convenience function for use in an joshua:ask continuation. joshua:graph-query-results draws a graph of the support information in backward-support, that is, the successful query, and the reasons it succeeded.

backward-support is fully described in the dictionary entry for joshua:ask. joshua:graph-query-results both extracts and interprets the information for you.

backward-support A support argument passed by joshua:ask to a continuation.

:orientation Specifies the graph orientation. Default is vertical.

:stream The stream on which the graph is output. Default is *standard-output*.

The convenience function joshua:print-query-results prints the same information as joshua:graph-query-results.

The accessor function joshua:ask-derivation extracts all the support for a satisfied query but without interpreting it. For the sake of comparison we'll use the same examples to illustrate all three of these functions.

Examples: First, a query satisfied from the database. The graph shows the database predication that matched the query.

```
(define-predicate edible (object))
(define-predicate is-food (object))
(define-predicate contains (object substance))
(define-predicate sweet (object))
(define-predicate type-of (object type))

(tell [edible chocolate-coated-ants])
(tell [contains chocolate-coated-ants honey])
```

![Ask Edible Query Results]

```
(ask [edible #what] #\'graph-query-results)
```

The next example shows the support for a query that is satisfied from rules. We have a rule, dessert?, that determines if a given food is a dessert. Each of this rule's subgoals is derived from other rules. Here are the rule definitions.

```
(defrule food? (:backward)
  if [edible ?object]
  then [is-food ?object])
```
(defrule sweet? (:backward)
  if [or [contains ?object chocolate]
      [contains ?object sugar]
      [contains ?object honey]]
  then [sweet ?object])

(defrule dessert? (:backward)
  if [and [is-food ?object]
         [sweet ?object]]
  then [type-of ?object dessert])

Now we joshua:ask what foods qualify as desserts and why. In the graph, ovals denote queries that were not satisfied directly by the database. Rectangles denote queries that were satisfied by the database.

The top of the graph shows the satisfied goal, and names the rule that satisfied it. The rest of the graph shows successive subgoals and how each was satisfied.

Since backward chaining stops when it finds database predications, the bottom leaves of the graph tree are queries that were satisfied by the database. Hence they are rectangles, whereas intermediate nodes are ovals.

The arrows move in the if-then (logical conclusion) direction.

Here's an extension to the previous example, to show how the graph displays truth values of joshua:*false*. We add a rule to eliminate first course choices: the rule says that things that are liquid and are not desserts are not a main course.

(define-predicate is-consistency-of (food consistency))

(defrule soup? (:backward)
  if [and [not [type-of ?food dessert]]
       [is-consistency-of ?food liquid]]
  then [not [type-of ?food main-course]]
(tell [not [type-of chicken-broth dessert]])
(tell [is-consistency-of chicken-broth liquid])

(ask [not [type-of chicken-broth main-course]] 'graph-query-results)
  (not [TYPE-OF CHICKEN-BROTH MAIN-COURSE])
    Rule SOUP?
  (not [TYPE-OF CHICKEN-BROTH DESSERT])
    Database (False)
      [TYPE-OF CHICKEN-BROTH DESSERT]
    Database (False)
      [IS-CONSISTENCY-OF CHICKEN-BROTH LIQUID]

The graph displays the satisfied query prefixed by [not ...]. The database predicate matching the query appears without the prefix, just as it would in the database display. The label above it indicates that its truth value is joshua:*false*. (Predicates with a truth value of joshua:*true* are not labelled as such in the graph Database heading.)

Related Functions:
  joshua:ask
  joshua:print-query-results

See the section "Querying the Database", page Querying-the-Database-SECTION. See the section "Explaining Backward Chaining Support", page Explaining-Backward-Chaining-Support-SECTION.

joshuagraph-tms-support &rest predicates Function
Displays a graph of the TMS support for predicates, that is, of the dependency information which a Truth Maintenance System stores in the database along with predicates. The graph traces the support chain through the dependency records created by forward rules (or other callers of joshua:justify such as the the justification keyword argument to joshuactell) to the underlying primitive support (assumptions and premises). (Backward chaining support is not graphed, since the rule result is not stored in the database. For that, you probably want joshuagraph-query-results.)

Example:

(define-predicate dreams-in (language dreamer) (ltms:ltms-predicate-model))
(define-predicate counts-in (language person) (ltms:ltms-predicate-model))
(define-predicate native-speaker-of (language speaker)
  (ltms:ltms-predicate-model))

(defrule native-speaker? (:forward)
  if [and [dreams-in ?language ?person]
        [counts-in ?language ?person]]
  then [native-speaker-of ?language ?person])
(tell [dreams-in Spanish Violet] :justification :assumption)
(tell [counts-in Spanish Violet])

Show Joshua Database (matching pattern [default All])
   [native-speaker-of ?x ?y] (opposite truth-value too? [default Yes]) Yes
   True things
   [NATIVE-SPEAKER-OF SPANISH VIOLET]
   False things
   None

-exec (graph-tms-support [NATIVE-SPEAKER-OF SPANISH VIOLET])-

NIL

You must give joshualanguage-tms-support the actual predication object that
resides in the database, rather than a copy of it. In our example we re-
trieve the predication object by clicking the mouse over it in the database
display.

Since the support graph traces the support for facts that are in the
database, all nodes are rectangles. (Compare the display of joshualanguage-
query-results.) The top of the graph tree shows the predication whose sup-
port we want to know about. We see that this predication was derived from
a forward rule, which in turn was derived from some predications. The bot-
tom leaves of the graph tree show primitive support (premise or assump-
tion) denoting the end of the forward chaining process. The arrows point in
the if-then (logical conclusion) direction.

Here’s an example showing the support graph for a predication whose truth
value is joshua:*false*.

(define-predicate has-ticket (claimant)(lms:lms-predicate-model))
(define-predicate admissible (claimant)(lms:lms-predicate-model))

(defrule no-free-lunch (:forward)
    if [not [has-ticket ?x]]
    then [not [admissible ?x]]])

(tell [not [has-ticket Jane]]
Show Joshua Database (matching pattern [default All]) All

True things

None

False things

[ADMISSIBLE JANE]
[HAS-TICKET JANE]

(graph-tms-support ¬[ADMISSIBLE JANE])

[ADMISSIBLE JANE]
(False) Rule NO-FREE-LUNCH

[HAS-TICKET JANE]
(False) :PREMISE

NIL

Predications with a truth value of joshua:*false* appear with an indication that they are false.

See the section "Explaining Program Beliefs", page Explaining-Program-Beliefs-SECTION.

joshuaknown proposition

Joshua Predicate

This modal operator checks if proposition is known to be either joshua:*true* or joshua:*false*.

proposition

A Joshua predication pattern to match.

The query: (ask [known [foo ?x]] #’ ...) succeeds when: either [foo ?x] or [not [foo ?x]] succeed

If successful, joshuaknown calls the continuation on the instantiated query.

Examples:

We use the predicate shape-of and the statements about shapes that we used to illustrate the predicate joshuaprovable. Here they are.

(define-predicate shape-of (object shape))

(tell [and [shape-of door oval]
[not [shape-of leaf pointed]]]
[AND [SHAPE-OF DOOR OVAL] [NOT [SHAPE-OF LEAF POINTED]]]

Show Joshua Database

True things

[SHAPE-OF DOOR OVAL]

False things

[SHAPE-OF LEAF POINTED]]

The database contains one statement about shapes that is joshua:*true* and one that is joshua:*false*. joshuaknown succeeds in each case, re-
turning the instantiated query. Note that there is no indication of truth value in the instantiated query. That is because when we ask if something is **joshua:known**, we are interested only in the existence of an answer, not in its particular truth value. (**backward-support** for the **joshua:ask** does indicate what the truth value of the instantiated query was.)

(ask [known [shape-of ?object ?shape]] #'print-query)

[KNOWN [SHAPE-OF DOOR OVAL]]

[KNOWN [SHAPE-OF LEAF POINTED]] ; argument was actually false

A more interesting question is to ask whether a predication is *not* known to Joshua.

The query: (ask [not [known [foo ?x]]]] #'...

Succeeds when: [foo ?x] and [not [foo ?x]] both fail

Examples:

; The proposition is not in the database or in rules

(ask [not [known [shape-of nose pointed]]] #'print-query)

[not [KNOWN [SHAPE-OF NOSE POINTED]]]

**joshua:known** can also be used in backward rules. The goal of the very inconsiderate rule in the next example is to select a dancing partner. The rule filters out those whose ability at ?activity is unknown, keeping those who are good or bad.

(define-predicate need-a-partner (activity))
(define-predicate is-good-at (activity person))
(define-predicate use-as-partner (person activity))

(defrule two-left-feet-will-do (:backward)
 if [and [need-a-partner ?activity]
       [known [is-good-at ?activity ?person]]]
 then [use-as-partner ?person ?activity])

(defun test-known ()
 (clear)
 (tell [and [need-a-partner dancing]
           [is-good-at dancing Tom]
           [not [is-good-at dancing Fred]]])

'Done.)

(test-known)
DONE.

(ask [use-as-partner ?person ?activity] #'print-query)

[USE-AS-PARTNER TOM DANCING]

[USE-AS-PARTNER FRED DANCING]
The goal of the rule in the next example is to hire an applicant if his/her qualifications are excellent, even if nothing is known about the applicant's experience level.

(define-predicate has-qualifications (person qualifications))
(define-predicate previous-experience (person experience))
(define-predicate hire-candidate (name))

(tell [and [has-qualifications Fred poor]
        [has-qualifications Joan excellent]]]
[AND [HAS-QUALIFICATIONS FRED POOR] [HAS-QUALIFICATIONS JOAN EXCELLENT]]

(defrule inexperience-no-obstacle (:backward)
  if [and [has-qualifications ?applicant excellent]
       [not [known [previous-experience ?applicant ?how-much]]]]
  then [hire-candidate ?applicant])

(ask [hire-candidate Fred] #'print-query)
(ask [hire-candidate ?applicant] #'print-query)
[HIRE-CANDIDATE JOAN]

Related Predicate:
  joshua:provable

### joshua:make-predication statement &optional area

**Function**

Construct a predication out of the specified `statement` (in the optional `area` supplied). The newly constructed predication is *not* entered in the database, unless you combine `joshua:make-predication` with `joshua:tell`.

You should seldom need to know about this, as the `[ ]` syntax is used in Joshua contexts as a reader macro for `joshua:make-predication`.

- `statement`: A list whose first element is the name of a (defined) predicate. The rest of the list elements are the arguments to the predicate.
- `area`: Storage area to cons in

#### Examples:

(define-predicate shape-of (object shape))

(make-predication '(shape-of window round))
[SHAPE-OF WINDOW ROUND] ; this is not in the database

(tell (make-predication '(shape-of window round)))
[SHAPE-OF WINDOW ROUND] ; new predication added to the database
joshuamake-predication is useful for constructing Joshua predications from data generated within Lisp code. (Still, backquoting [ ] expressions should suffice most of the time.)

Related Functions:

joshua:define-predicate
See the section "Predications and Predicates", page Predications-and-Predicates-SECTION.

joshua:map-over-database-predications

Predicate pattern function
A convenience macro for joshua:ask. Use it whenever you want to find an answer to a query in the database without using rules or questions.

joshua:map-over-database-predications finds all database predications that unify with predication-pattern and applies function to each.

predication-pattern A pattern to match against database predications.
function Specifies the operation to do on each database predicate that unifies with predication-pattern. Should be a function of one argument.

(map-over-database-predications <predication> <continuation>) is equivalent to:

(ask [foo ?x]
  #'(lambda (support)
     (funcall <cont>
       (ask-database-predication support)))
     :do-backward-rules nil)

Example:
We'll build an author-title index for a library, using joshuatell statements.
We'll include an LTMS in our predicate definitions, so that we can later get joshua:explain to tell us about some database predications.

(define-predicate author-of (work author) (ltms:ltms-predicate-model))

(defun build-author-title-index1 ()
  (clear)
  (tell [and [author-of "The Interpretation of Dreams" Freud]
           [author-of "Hedda Gabler" Ibsen]
           [author-of "Totem and Taboo" Freud]
           [author-of "A Doll's House" Ibsen]]
           (cp:execute-command "Show Joshua Database"))
  BUILD-AUTHOR-TITLE-INDEX1
The first example looks in the library database and removes from it all of Freud’s books (perhaps for rebinding due to overuse). We use `joshua:map-over-database-predications` to get our hands on the actual predication objects so that we can remove them.

To allow easy replacement of this information we’ll `joshua:unjustify` the facts rather than actually removing them with `joshua:untell`. The truth value of each of these facts becomes `joshua:*unknown*`, even though they physically remain in the system.

```lisp
(defun away-with-sigmund ()
  (map-over-database-predications [author-of ?work Freud] #'unjustify)
  (cp:execute-command "Show Joshua Database"))

AWAY-WITH-SIGMUND

(defun away-with-SIGMUND ()
  (map-over-database-predications [author-of ?work Shakespeare] #'unjustify)
  (cp:execute-command "Show Joshua Database"))

AWAY-WITH-SIGMUND

Let’s add a forward rule that says the library owns any work that was authored by Shakespeare, and then build another database.

```lisp
(define-predicate owns-library (work) (ltms:ltms-predicate-model))
```
(defrule Shakespeare-holdings (:forward) 
  if [author-of ?work Shakespeare] 
  then [owns-library ?work])

(defun build-author-title-index2 () 
  (clear) 
  (tell [and [author-of "King Lear" Shakespeare] 
            [author-of "Hedda Gabler" Ibsen] 
            [owns-library "Trumpeting Joshua"] 
            [author-of "A Doll's House" Ibsen]]) 
  (cp:execute-command "Show Joshua Database") 
BUILD-AUTHOR-TITLE-INDEX2 

(build-author-title-index2) 

True things 
[OWNS-LIBRARY "Trumpeting Joshua"] [AUTHOR-OF "Hedda Gabler" IBSEN] 
[OWNS-LIBRARY "King Lear"] [AUTHOR-OF "King Lear" SHAKESPEARE] 
[AUTHOR-OF "A Doll's House" IBSEN] 
False things 
None 

We can now ask Joshua to joshua:explain the database predications about 
works the library owns. 

(map-over-database-predications [owns-library ?work] #'explain) 
[OWNS-LIBRARY "Trumpeting Joshua"] is true 
It is a :PREMISE 
[OWNS-LIBRARY "King Lear"] is true 
It was derived from rule SHAKESPEARE-HOLDINGS 
[AUTHOR-OF "King Lear" SHAKESPEARE] is true 
It is a :PREMISE 

Here's an example showing the display when the database predication has a 
truth value of joshua:*false*. The predication displays without indicating 
its truth value; that information is supplied by the accompanying explana-
tion. 

(tell [not [owns-library "Everyday Sanskrit"]]) 

¬[OWNS-LIBRARY "Everyday Sanskrit"] 
T 

(map-over-database-predications [not [owns-library ?work]] #'explain) 
[OWNS-LIBRARY "Everyday Sanskrit"] is false 
It is a :PREMISE 

The accessor function joshua:ask-database-predication can also be used to 
extract database predications from the backward support supplied to the 
joshua:ask continuation. Most of the time joshua:map-over-database-
predications probably serves just as well, and it is easier to use. For com-
parison we are using the same examples to illustrate both functions.
Related Functions:

**joshua:ask**

See the section "Querying the Database", page Querying-the-Data-Base-SECTION.

**joshua:predication**

*Flavor*

The non-instantiable base flavor for all predications in Joshua. It is mixed into new predications via **joshua:define-predicate**.

You can test for this flavor by using **typep** or **joshua:predicationp** (into which **typep** is optimized).

Related Presentation Types:

**joshua:predication**

**joshua:database-predication**

**joshua:predicationp object**

*Function*

Checks whether **object** is a Joshua predication, that is, whether the object is built on the base flavor **joshua:predication**. **joshua:predication** is the root of the Joshua model tree.

**joshua:predicationp** returns **t** if the object is a Joshua predication, otherwise **nil**.

**object**

An object in the Lisp world.

Examples:

```lisp
(define-predicate valid-word (word language))

(tell [valid-word incarnadine English]
[VALID-WORD INCARNADINE ENGLISH]
T)

(predicationp [VALID-WORD INCARNADINE ENGLISH])
; click on object returned by tell
(predicationp FLAVOR:VANILLA)

(ask [valid-word incarnadine ?language]
#'(lambda (backward-support)
  (when (predicationp (ask-database-predication backward-support))
    (print (ask-database-predication backward-support)))
[VALID-WORD INCARNADINE ENGLISH])
You can use **typep** to do the same test as **joshua:predicationp**. In fact, the compiler optimizes the form:

(typep x 'predication)

into the form:

(predicationp x)
For example:

```lisp
(ask [valid-word incarnadine ?language]
  #'(lambda (backward-support)
      (when (typep (ask-database-predication backward-support) 'predication)
        (print (ask-database-predication backward-support)))))
[VALID-WORD INCARNADINE ENGLISH]
```

Related Functions:

- `joshua:predication`
- `typep`

**joshua:print-query** `backward-support &optional (stream *standard-output*)`

A convenience function for use in an `joshua:ask` continuation. `joshua:print-query` displays the `joshua:ask` query with its variables instantiated.

*backward-support* The backward support supplied to the `joshua:ask` continuation.

*stream* A stream to which to output the information. Defaults to `*standard-output*`.

Examples:

```lisp
(define-predicate type-of (object type))
(tell [type-of Iliad epic])

[TYPE-OF ILIAD EPIC]
```

If you want to use the instantiated query in ways other than printing it, extract it yourself using the accessor function `joshua:ask-query`.

Related Functions:

- `joshua:ask`
- `joshua:graph-query-results`
- `joshua:print-query-results`
- `joshua:say-query`

See the section "Querying the Database", page Querying-the-Database-SECTION.

**joshua:print-query-results** `backward-support &key (:stream *standard-output*) (:printer #'print1)`

A convenience function for use in an `joshua:ask` continuation. `joshua:print-query-results` displays and interprets the support information
in the **joshua:ask** continuation argument, **backward-support**; that is, it tells you what queries succeeded, and why.

**backward-support**  A list containing the satisfied query and information about its support.

**stream**  A stream to which to output the information. Default is `*standard-output*`.

**printer**  A function of two arguments, like `print1`, that is used to print elements of the support. `print1` is the default, but another reasonable value to give is `joshua:say`.

Use `joshua:graph-query-results` to see a graph of the information provided by `joshua:print-query-results`.

The accessor function `joshua:ask-derivation` extracts the support portion of `backward-support` but does not interpret the information.

For comparison, we use the same examples to illustrate all three functions.

Examples:

The first example shows a query satisfied by database lookup. Both the instantiated query and its support (here the matching database predication) are printed.

```
(define-predicate type-of (object type))

(tell [type-of Iliad epic])

[TYPE-OF ILIAD EPIC] succeeded: [TYPE-OF ILIAD EPIC] was TRUE in the database
```

The next example shows the support for a query that is satisfied from rules. We have a rule, `dessert?`, that determines if a given food is a dessert. Each of this rule’s subgoals is derived from other rules. Here are the definitions.

```
(define-predicate edible (object))
(define-predicate is-food (object))
(define-predicate contains (object substance))
(define-predicate sweet (object))

(defrule food? (:backward)
  if [edible ?object]
  then [is-food ?object])
```
(defrule sweet? (:backward)
  if [or [contains ?object chocolate]
      [contains ?object sugar]
      [contains ?object honey]]
  then [sweet ?object])

(defrule dessert? (:backward)
  if [and [is-food ?object]
        [sweet ?object]]
  then [type-of ?object dessert])

; tell some sticky facts
(tell [edible chocolate-coated-ants])
(tell [contains chocolate-coated-ants honey])

Now we joshua:ask what foods qualify as desserts and why. A single food,
chocolate-covered-ants, succeeded. The display shows the instantiated
query, explaining why it succeeded: support is traced backward from rule
dessert? that satisfied the query, through the support used to satisfy parts
of the rule body.

(ask [type-of ?object dessert] #'print-query-results)
[TYPE-OF CHOCOLATE-COOED-ANTS DESSERT] succeeded
  It was derived from rule DESSERT?
[IS-FOOD CHOCOLATE-COOED-ANTS] succeeded
  It was derived from rule FOOD?
[EDIBLE CHOCOLATE-COOED-ANTS] succeeded
  [EDIBLE CHOCOLATE-COOED-ANTS] was true in the database
[SWEET CHOCOLATE-COOED-ANTS] succeeded
  It was derived from rule SWEET?
[CONTAINS CHOCOLATE-COOED-ANTS HONEY] succeeded
  [CONTAINS CHOCOLATE-COOED-ANTS HONEY] was true in the database

Related Functions:
  joshua:ask
  joshua:graph-query-results
  joshua:print-query
  joshua:say-query

See the section "Querying the Database", page Querying-the-Database-
SECTION. See the section "Explaining Backward Chaining Support", page
Explaining-Backward-Chaining-Support-SECTION.

joshua:provable proposition  Joshua Predicate
Checks if proposition is known to be joshua:*true*, (or if it is known to be
joshua:*false*, if [not ...] is wrapped around it.)
This is a modal operator, [provable ...] and [not [provable ...]] correspond to the "box" and "diamond" operators of some modal logics.

**proposition**

A Joshua predication pattern to match.

The query: (ask [provable [foo ?x]] #’ ...)  
Succeeds when: [foo ?x] would succeed

The query: (ask [provable [not [foo ?x]] #’ ...])  
Succeeds when: [not [foo ?x]] would succeed

If successful, **joshua:provable** calls the continuation on the instantiated query.

**Examples:**

Let’s define a predicate, shape-of, **joshua:tell** some statements about the shape of objects, and then display the database.

\[
\text{(define-predicate shape-of (object shape))}
\]

\[
\text{(tell [and [shape-of door oval]}
\text{[not [shape-of leaf pointed]]])}
\]

\[
\text{[AND [SHAPE-OF DOOR OVAL] [NOT [SHAPE-OF LEAF POINTED]]]}
\]

Show Joshua Database  
True things  
[SHAPE-OF DOOR OVAL]  
False things  
[SHAPE-OF LEAF POINTED]]

Now we can check which statements about shapes are **joshua:*true*», and which are **joshua:*false*».

```lisp
;; Check if the proposition is **joshua:*true*»  
(ask [provable [shape-of door oval]] #’print-query)  
[PROVABLE [SHAPE-OF DOOR OVAL]]
```

```lisp
;; Comparing provable to known  
(ask [provable [shape-of leaf pointed]] #’print-query)  
; this fails  
(ask [known [shape-of leaf pointed]] #’print-query)  
[KNOWN [SHAPE-OF LEAF POINTED]]
```

```lisp
;; Check if the proposition is **joshua:*false*»  
(ask [provable [not [shape-of leaf pointed]]]] #’print-query)  
[PROVABLE [NOT [SHAPE-OF LEAF POINTED]]]
```

```lisp
(ask [provable [not [shape-of ?object ?shape]]]] #’print-query)  
[PROVABLE [NOT [SHAPE-OF LEAF POINTED]]]
```
Reset Joshua Tracing Command

Resets the tracing options to the original defaults.

*Type of Tracing* Which type of tracing to reset. The possible types are forward rules, backward rules, predications, TMS operations and All.

*Include events* Whether to reset the traced and stepped events for the *Type of Tracing* as well.

The Reset Joshua Tracing command sets the Joshua tracing options back to their initial defaults. This command is useful if you have been selectively tracing rules or predications and would like to go back to tracing all rules or all predications. The *Include events* option comes in handy when you have been tracing or stepping particular events and would like to go back to just tracing the default events. This command does not disable or enable tracing, it just affects which things are traced.

Related Commands:

*Enable Joshua Tracing Command*

*Disable Joshua Tracing Command*

*\say* *predication*  

A *format* directive that makes it easy to combine the use of *joshuasay* with other kinds of formatted output. It takes one format-argument, the predication to be *joshuasay’d* to the output stream.
Examples:

(format t " "&The registry of deeds says that "\say\\." [frobozz Prospero 1616 remote-island])

This would print the following sentence:
The registry of deeds says that PROSPERO was an owner of a frobozz in 1616 at REMOTE-ISLAND.

You can also use "\say\\" in other places format strings are used, for instance prompt-and-accept:

(prompt-and-accept 'integer "For what values of "S is it true that "\say\\?" ?x [Riemann-zeta 3 ?x])

Related Functions:

joshua:say

See the section "Formatting Predications: the SAY Method", page Formatting-Predications-the-SAY-Method-SECTION.

**joshua:say** predication &optional (stream *standard-output*)

*Function*

Prints out predication on stream, possibly in a way other than print would. This is good for printing the meaning of a predication in natural language, as opposed to the predicate calculus notation in which programs are written. However, you needn't restrict your thinking about joshua:say to just natural language. For example, joshua:say could present a predication as a piece of graphics; see examples below. Judicious use of joshua:say methods can make it easier to generate user interfaces.

It usually doesn't matter what value the implementations of joshua:say return, since joshua:say is usually done for side-effect. The exception is that if stream is explicitly supplied as nil, the implementations should do what format would do, that is, return a string if possible. (Graphical joshua:say methods can't do this.)

Examples:

(define-predicate frobozz (who when where) ()
 :destruct-into-instance-variables)

(define-predicate-method (say frobozz) (optional (stream *standard-output*))
 (format stream ""S was an owner of a frobozz in "S at "S." who when where))

(say [frobozz Prospero 1616 remote-island])

prints the sentence:

PROSPERO was an owner of a frobozz in 1616 at REMOTE-ISLAND.

An example using graphics would be:
(define-predicate-method (say frobozz) (optional (stream *standard-output*))
  (dw:with-output-as-presentation
   (:stream stream :object self :type (type-of self))
   (format-graph-from-root (list who (list where) (list when))
     #'(lambda (x s) (prin1 (car x) s))
     #'cdr
     :stream stream)))

The **joshua:say** method now draws a graph representing Prospero's relationship to his property and the time at which he owned it.

(C: (say [frobozz Prospero 1616 remote-island])

```
PROSPERO

REMOTE-ISLAND

1616
```

#:<DW::DISPLAYED-PRESENTATION [FROBOZZ PROSPERO 1616 REMOTE-ISLAND]>

Related Functions:

"""\Say\"

See the section "Formatting Predications: the SAY Method", page Formatting-P predications:-the-SAY-Method-SECTION.

**joshua:say-query** backward-support &optional (stream *standard-output*)

A convenience function for use in a **joshua:ask** continuation. **joshua:say-query** displays the instantiated query using a user-defined **joshua:say** method if available, or the default **joshua:say** method. The latter simply prints the instantiated query.

**backward-support** The support supplied to the **joshua:ask** continuation.

**stream** A stream to which to output the information. The default is *standard-output*.

Examples:

```
;; say-query with default say method
(define-predicate loves (person object))

(tell [loves Bob chocolate])

(ask [loves Bob ?x] #'say-query)
[LOVES BOB CHOCOLATE]

;; say-query with user-defined say method
(define-predicate type-of (object type))
```
(define-predicate-method (say type-of) (&optional (stream *standard-output*))
  (with-statement-destructured (object type) ()
    (format stream
      "%- The "A is an example of the "A literary form." object type))
  )
(tell [type-of Iliad epic])
[TYPE-OF ILIAD EPIC]

The Iliad is an example of the Epic literary form.

To use the instantiated query in some other way rather than joshua:saying it, extract it from the continuation argument using the accessor function joshua:ask-query, and interpret the information.

Related Functions:
  joshua:ask
  joshua:graph-query-results
  joshua:print-query
  joshua:print-query-results

See the section "Querying the Database", page Querying-the-Database-SECTION.

Show Joshua Database Command

Displays the contents of the Database, or a subset of the contents matching a certain pattern.

matching pattern Specifies the predication patterns to display. The default is the entire database.

The display groups predications under the headings True and False, for predications with a truth value of joshua:*true* and joshua:*false*, respectively.

When specifying a pattern you can further limit the display to patterns with either truth value.

Examples:

Show Joshua Database (matching pattern [default All]) All
  True things ; indication of truth value is in the heading
    [DREAMS-IN SPANISH LUCINDA] [NATIVE-SPEAKER-OF SPANISH LUCINDA]
    [DREAMS-IN SUMERIAN DR-PARCHMENT] [NATIVE-SPEAKER-OF GERMAN DR-PARCHMENT]
    [COUNTS-IN SPANISH LUCINDA]
  False things
    [COUNTS-IN GERMAN HENRY] ; indication of truth value is in the heading
Show Joshua Database (matching pattern [default All]) [dreams-in ?x ?y] (opposite truth-value too? [default Yes]) Yes
True things
[DREAMS-IN SPANISH LUCINDA]
[DREAMS-IN SUMERIAN OR-PARCHMENT]
False things
None

See the section "Entering and Displaying Predications in the Database", page Entering-and-Displaying-Predications-in-the-Database-SECTION.

Show Joshua Predicates Command

Shows the currently defined Joshua predicates.

:Include Models Whether to include predicates that are used as base flavors for building other predicates in the output.

:Matching Show only predicates whose names contain a substring or substrings.

:Output Destination Where to display the information.

:Packages Only show predicates in the specified package or packages. Supply a value of All to see all the currently defined Joshua predicates. Unless you otherwise specify the package, you see only the predicates defined in the current package.

:Search Inherited Symbols Whether to include predicates that are inherited by the packages specified in :Packages.

:System Show only the predicates that are defined in a particular system.

The Show Joshua Predicates command provides a convenient tool for browsing through all the predicates defined in the current world. The output is a table of predicate names and arguments. There are a number of mouse behaviors defined for the predicate names that this command displays. These can be seen by mousing right on the name.

Show Joshua Predicates :Packages TME
TME:ABNORMAL (WHO FOR-WHAT) TME:LOVES (LOVER LOVEE)
TME:BIRD (BOID) TME:ONE-PER-ROW-OR-COL (R-OR-C INDEX)
TME:FLY (BOID) TME:PENGUIN (BOID)
IS-EXAMPLE-OF (NAME TYPE) PROVABLE (PROPOSITION)
TME:JEALOUS (WHO) TME:QUEEN (ROW COL)
TME:KILLS (KILLER VICTIM) TME:TRAGEDY (EVENT)
KNOWN (PROPOSITION)

Related Commands:
"Show Joshua Rules Command"
"Show Joshua Tracing Command"

Show Joshua Rules Command

Displays the currently defined rules.

:Triggered By    Show rules with one or more triggers that unify with the specified predication.

:Matching        Show rules with names containing one or more substrings.

:Output Destination Where to display the output from this command.

:Packages        Show the rules defined in which package or packages. This defaults to the current package.

:Search Inherited Symbols Include rules that are inherited by Packages.

:System          Show only the rules defined in a particular system.

:Type            Show only backward or forward rules. By default the command shows both backward and forward rules.

The Show Joshua Rules command provides a tool for browsing through all the Joshua rules. It displays a table of all the rules satisfying the given arguments. Mousing middle on a rule name displays the most recent definition of that rule.

Example:

Show Joshua Rules :Triggered By [tme:loves ? ?] :Packages All

Forward Rules:
JEALOUSY    LOVE-IN-IDLENESS ONLY-ONE-LOVE QUALITY-NOT-QUANTITY
UNREQUITED-LOVE

The above example lists all of the rules that could be triggered by a predication of the form [tme:loves ? ?].

Related Commands:

"Show Joshua Predicates Command"
"Show Joshua Tracing Command"

Show Joshua Tracing Command

Shows information about Joshua tracing.

Type of Tracing Which type of tracing to describe. It can be one of forward rules, backward rules, predications, TMS operations, or all.

:Output Destination Where to display the output from this command.

The Show Joshua Tracing command describes the current state of Joshua tracing,
saying whether each Type of Tracing is on or off. For each active Type of Tracing, Show Joshua Tracing prints out information about the current options and traced events.

Example:

```
Show Joshua Tracing (type of tracing) All
```

**Forward Chaining** tracing is on
Tracing all forward rules triggered by a predication matching:

[TME:LOVES DEMETRIUS HERMIA]
Traced events: Fire and Queue

---

**TMS** tracing is off

---

**Predication** tracing is on
Tracing predications of flavor: LTMS:LTMS-PREDICATE-MODEL
Traced events: Ask and Tell

---

**Backward Chaining** tracing is off

---

Related Commands:

"Show Joshua Rules Command"
"Show Joshua Predicates Command"

**Show Rule Definition Command**

Shows the latest definition of a Joshua rule.

- **Rule**: Show the definition of which rule or rules.
- **:Load**: This argument controls the behavior of the command when the desired rule definition is not currently in an editor buffer. If you enter Yes, the command loads the definition into an editor buffer. If you enter No, it does not. The value of Load defaults to Query, meaning the command should ask you before loading any file into the editor.

- **:Output Destination**: Where to display the output from this command.

The Show Rule Definition command allows you to see the definition of a Joshua rule in a Lisp Listener without having to enter the editor. When the rule definition can be found in the editor the command displays the latest version. Otherwise, depending on the value of Load, the command offers to read in the latest definition from the file containing the rule definition.

Example:
Show Rule Definition JEALOUSY

Rule Jealousy:
(defrule jealousy (:forward :importance 3)
  (IF [and [jealous ?x]
      [loves ?x ?y]
      [loves ?z ?y]
      (different-objects ?x ?z)]
  THEN [kills ?x ?z])

joshua:succeed &optional support  Function

Joshua is a success-continuation-passing language. In most places, calling
the continuation means "go ahead with the rest of the computation". Based
on context, the form joshua:succeed finds the continuation and calls it ac-
cordingly.

You can use joshua:succeed within Lisp code embedded in:

- The if-part of rules (in Lisp code in forward rules, and in multiply-
succeeding Lisp forms of backward rules)
- The body of a joshua:defquestion

It makes no sense to call joshua:succeed elsewhere.

The optional support argument allows the Lisp code to specify the deriva-
tion information for the query.

Example:

(define-predicate good-to-read (book))

(defparameter *books* `(decameron canterbury- tales gargantua-and-pantagruel
tom-jones catch-22))

(defrule reading-list (:backward)
  if (typecase ?candidate-book
    (unbound-logic-variable
      (loop for book in *books*
        doing (with-unification
          (succeed 'Humor-101-reading-list)))))
  (otherwise
    (when (member ?candidate-book *books*)
      (succeed (succeed 'Humor-101-reading-list))))
  then [good-to-read ?candidate-book])
(ask [good-to-read ?x] #'print-query-results)
[GOOD-TO-READ DECAMERON] succeeded
  It was derived from rule READING-LIST
  HUMOR-101-READING-LIST
[GOOD-TO-READ CANTERBURY-TALES] succeeded
  It was derived from rule READING-LIST
  HUMOR-101-READING-LIST
[GOOD-TO-READ GARGANTUA-AND-PANTAGRUEL] succeeded
  It was derived from rule READING-LIST
  HUMOR-101-READING-LIST
[GOOD-TO-READ TOM-JONES] succeeded
  It was derived from rule READING-LIST
  HUMOR-101-READING-LIST
[GOOD-TO-READ CATCH-22] succeeded
  It was derived from rule READING-LIST
  HUMOR-101-READING-LIST

Related Functions:

  joshua:unify
  joshua:with-unification

**joshua:support database-predication &optional filter**

Examines the TMS justification structures currently supporting belief in database-predication, tracing them back to primitively justified predications (i.e. to those whose support does not depend on any other predications). Returns a list of the primitive support (assumptions and premises). Filter, if provided, is a predicate to be applied to the support. Only those elements of the primitive support which satisfy the predicate are collected.

*database-predication*  A predication object that is in the database. Must be the actual database object, and not a copy of it.

*filter*  If filter is not supplied the value default to nil which means that all the primitive support should be collected and returned. Otherwise, filter should be a function of one argument that returns non-nil on the support you want. (For example, you might want to look at just the assumption support of database-predication.) When the database-predication argument is based on a TMS, this function is passed a justification as its argument. It may examine the justification using joshua:destructure-justification.

Examples:

Prospero, curious about his daughter’s relationship with Caliban, might do:
(ask [is-friend-of Miranda ?]
  #'(lambda (backward-support)
      (format t "&The support for ~S is ~S."
             (ask-database-predication backward-support)
             (support (ask-database-predication backward-support))))
  :do-backward-rules nil)

If he wanted to see just the assumptions underlying it, he would do:

(ask [is-friend-of Miranda ?]
  #'(lambda (backward-support)
      (format t "&The support for ~S is ~S."
             (ask-database-predication backward-support)
             (support (ask-database-predication backward-support)
                      #'(lambda (justification)
                          (multiple-value-bind (ignore mnemonic)
                                          (eq mnemonic :assumption)))))
      :do-backward-rules nil)


**joshuatell** *predication* &key :justification

*Function*

Puts a predication into the virtual database.

**Note:** *joshuatell* is a macro, and as such it cannot be used as an argument to the function **funcall**.

*predication* should be thought of as a pattern argument, not as the actual data in the database. If something already exists in the database that is a *joshua:variant* of *predication*, the returned (canonical) value will not be **eq** to *predication*. Thus *joshuatell* serves as an interner for *predication*, that is, it gives you the canonical copy in the database, creating it if necessary.

If *predication* is not already in the database, the returned values are *predication* and the symbol **t**.

If something already exists in the database that is a *joshua:variant* of *predication*, *predication* is not put into the database, since that would be duplication. Instead, the canonical version found in the database is returned, along with the symbol **nil**.

**Justification** can be one of the following:

- **nil**, in which case a default justification is used. If the *joshuatell* occurs outside a rule, the default justification is **premise**. If the *joshuatell* is inside a rule, the default justification includes the rule name and the current support set.

- A **symbol**. A justification which is a symbol means that the truth-value of *predication* does not depend on that of any other predication; we say that *predication* has a *primitive justification*, in such a case. One primitive justification is specially treated by the LTMS provide with Joshua,
namely \texttt{:premise}. \texttt{:premise} justifications will never be removed by the LTMS without querying the user. Other primitive justifications are treated as assumptions that can be removed by the LTMS if necessary to resolve a contradiction.

- **A List of Four fields.** These are identical to the arguments to the Joshua protocol function \texttt{joshua:justify}, namely a mnemonic, \textit{true-support}, \textit{false-support} and \textit{unknown-support}. These fields are used (or discarded) by whatever TMS is present.

The database into which \textit{predication} is put depends on the data model of its predicate. The default is the discrimination net.

Examples:

\begin{verbatim}
(tell [is-magician Prospero])
(tell [not [is-magician Caliban])
(tell [is-daughter-of Miranda Prospero])
(tell [is-servant-of Caliban Prospero] :justification :premise)
(tell [is-friend-of Miranda Caliban] :justification :assumption)
; later retracted!
(tell '[is-exiled-from Prospero , (find-exile-country 'Prospero])
\end{verbatim}

\begin{center}
\begin{tabular}{|l|}
\hline
Note:
\text{Chances are that you seldom want to define a method that takes over the entire functionality of \texttt{joshua:tell}. It's more likely that you would want to define a method for one of the generic functions it calls, such as \texttt{joshua:insert}, \texttt{joshua:justify}, or \texttt{joshua:map-over-forward-rule-triggers}.}
\hline
\end{tabular}
\end{center}

Related Functions:
\begin{verbatim}
joshua:untell
joshua:clear
joshua:ask
joshua:justify
\end{verbatim}

See the section "Entering and Displaying Predications in the Database", page Entering-and-Displaying-Predications-in-the-Database-SECTION.


See the section "Customizing the Data Index" in Joshua Reference Manual.


\texttt{joshua:*true*} \hspace{1cm} \textit{Variable}

A named constant used by Joshua to denote a truth value of true. You can compare truth values using \texttt{eql}.

Related Topics:
joshua:*false*
joshua:*unknown*
joshua:*contradictory*
joshua:truth-value
joshua:predication-truth-value

See the section "Truth Values", page Truth-Values-SECTION.

**joshua:undefine-predicate** name  
*Macro*

"Undoes" a predicate definition. Predications built with this definition remain in the world, but an attempt to do almost anything to them results in an error.

**Example:**

```
(define-predicate fruit (a-fruit))
(undefine-predicate 'fruit)
```

You can perform the same operation from the Zmacs editor. Place your cursor on the predicate definition to be removed and use the command `m-X` Kill Definition. The system asks for confirmation in the minibuffer; then it offers you the options of removing the definition from the editor buffer itself, and of inserting the `joshua:undefine-predicate` command into the editor buffer.

**Example:**

### 1. Interaction During m-X Kill Definition

```lisp
;;; -x Mode: Joshua; Package: JOSHUA-USER; Syntax: Joshua; Usp: 0 -x
;;; Created 09/07 14:15:27 by Covo running on LRED-PERSEPHONE at SRED

(define-predicate needs-a-vacation (person))
```

Extended command:

Remove predicate NEEDS-A-VACATION from the current world? (Y or N) Yes.
Remove predicate NEEDS-A-VACATION from the editor buffer? (Y or N) No.
Insert form to kill predicate NEEDS-A-VACATION into the editor buffer? (Y or N) Y

### 2. Zmacs Buffer After Completion of m-X Kill Definition

```lisp
;;; -x Mode: Joshua; Package: JOSHUA-USER; Syntax: Joshua; Usp: 0 -x
;;; Created 09/07 14:15:27 by Covo running on LRED-PERSEPHONE at SRED.

(UNDEFINE-PREDICATE NEEDS-A-VACATION)
(define-predicate needs-a-vacation (person))
```

```
Zmacs (Joshua) questions-examples.lisp) joshua/doc/examples 0; (15) * [More below]
Predicate NEEDS-A-VACATION removed from the current world.
```
Related Functions:

**joshua:define-predicate**
"Zmacs Command: Kill Definition"

**joshua:undefquestion name**  Function

Removes a question definition from the system.

*name*  The name of the question

```lisp
(define-predicate foo (something something-else))

(defun question1 ( :backward) [foo 1 ?x])

(ask [foo 1 2] #'print-query :do-questions t)
Is it true that "[FOO 1 2]"? [default No]: Yes

[FOO 1 2]

(undefquestion 'question1)

QUESTION1

(ask [foo 1 2] #'print-query :do-questions t)
```

To kill a question definition from a Zmacs buffer, use the command *-K Kill Definition. For a sample interaction with the command: See the macro **joshua:undef-predicate**, page joshua:undef-predicate:FUNCTION.

Related Functions:

**joshua:defquestion**
"Zmacs Command: Kill Definition"

See the section "Asking the User Questions", page Asking-the-User-Questions:SECTION.

**joshua:undefrule**  Function

Removes a rule definition so that the rule cannot execute.

You can also remove a rule from a Zmacs buffer with *-K Kill Definition. For a sample interaction with the command: See the macro **joshua:undef-predicate**, page joshua:undef-predicate:FUNCTION.

*rule-name*  The name of the rule to be removed.

Examples:

```lisp
(defrule parched (:forward)
  if [condition-of plant-soil dry]
  then [needs plant-soil water])

(undefrule 'parched)
```
Modeling Note:

`joshua:undefrule` calls one of the generic functions `joshua:delete-forward-rule-trigger` or `joshua:delete-backward-rule-trigger` which removes the rule’s trigger from its storage place, so that it is no longer found by the trigger locating and trigger mapping functions.

See the section "The Contract of the Trigger Deleting Functions" in *Joshua Reference Manual*.

Related Functions:

`joshua:undefrule`

`joshua:clear`

"Clear Joshua Database Command"

"Zmacs Command: Kill Definition"

See the section "Rules and Inference", page Rules-and-Inference-SECTION.

`joshua:unify` `object1` `object2`

*Function*

If `object1` and `object2` unify, does so, while side-effecting any logic variables for the duration of the unification.

`object1` A pattern in Joshua, that is, a predication containing other predications, lists, symbols, numbers, or logic variables.

`object2` Another pattern.

Pattern matching underlies the inferencing process. In forward chaining, Joshua matches rule trigger patterns with database predications. In backward chaining, it matches goals with database predications and with rule and question trigger patterns.

Two patterns containing no logic variables *match* if they are structurally equivalent (if they "look the same").

Two patterns containing logic variables *unify* when one can substitute values for the variables so that both patterns become structurally equivalent. The process of doing so is called *unification*.

`joshua:unify` is useful for assigning values to logic variables within Lisp code in rule bodies. If the expressions are unifiable, the appropriate substitutions are made and rule execution continues.

If the expressions are not unifiable, rule execution fails. "Fails" means that it throws to the nearest (dynamically) containing `joshua:with-unification` clause.

Always wrap the macro `joshua:with-unification` around `joshua:unify` (or calls to functions that call `joshua:unify`) to establish the scope within which the substitutions remain in effect.
The Joshua unifier does what is called an occur check, that is, prevents the formation of certain circular structures by refusing to unify a logic variable with a structure in which it occurs. For example, if you tried to unify \( ?x \) with \([f ?x] \), you would get something whose printed representation would look (partially) like this:

\[
[f [f [f [f [f ...]
\]

This is exactly the same thing that happens when you make certain conses point at themselves — you get circular lists.

To see how this might happen, consider example 3 below.

Examples:

Example 1:
(defrule test-1 (:forward)
  if [and [balance-due ?applicant ?balance]
    [yearly-salary ?applicant ?salary]
    (unify ?cash-flow (- ?salary ?balance))
    (≤ ?cash-flow ?balance)]
  then [and [deny-credit ?applicant]
    (format t "~% Sorry, ~S, your cash-flow of ~S is insufficient."
      ?applicant ?cash-flow)])

(defun test-it ()
  (clear)
  (tell [yearly-salary Fred 20000])
  (tell [balance-due Fred 15000])
  (tell [yearly-salary George 20000])
  (tell [balance-due George 15000])
  'done-testing)
TEST-IT

(test-it)
Sorry, FRED, your cash flow of 5000 is insufficient.
DONE-TESTING
Show Joshua Database

True things

[BALANCE-DUE FRED 15000]
[YEARY-SALARY FRED 20000]
[YEARY-SALARY GEORGE 200000]
[BALANCE-DUE GEORGE 15000]
[DENY-CREDIT FRED] ;Inference added to database

False things

None

Example 2:
(with-unbound-logic-variables (x)
 (let ((p1 '[foo ,x])
 (p2 [foo 1]))
 (with-unification
 (unify p1 p2)
 ; If p1 and p2 don't unify, the next
 ; expression is not executed
 (format t """"&The value of x is """"s."""" x))))

The value of x is 1.
NIL

Example 3 shows a case where the occur-check feature makes the unification fail.

Example 3:
(define-predicate f (arg))
(define-predicate g (arg1 arg2))

(defun test-occur ()
 (with-unbound-logic-variables (x y)
 (with-unification
 (unify '+[g ,x ,x] '+[g ,y [f ,y]])
 ; ; if you get here, print Y and return
 (format t """"&You blew it. Y is now circular: """"s"""" y)
 (return-from test-occur :loser))
 ; ; if you got here, the unification failed
 :occur-check-forbids))

(test-occur)
:OCUR-CHECK-FORBIDS

This function attempts to unify [g ?x ?x] with [g ?y [f ?y]]. If it unifies, the function prints an abusive message and returns the symbol :loser. If the unification fails, it returns the symbol :occur-check-forbids.

Let's follow the unification and see what happens:
• The predicates in both places are $g$, so the unifier goes on to inspect the arguments.

• The first argument on the left is $?x$ and the first on the right is $?y$. The unifier unifies $?y$ and $?x$, which we can write as the equation $?x = $?y.

• The next argument on the left is $?x$ and the next on the right is $[f ?y]$. Thus the unifier attempts to enforce the equation $?x = [f ?y].

We thus have the two equations $?x = $?y and $?x = [f ?y]. Combining them, we have the single equation $?y = [f ?y], whose only solution is to unify $?y to a structure containing itself, that is, a predication that structurally resembles a circular list: $[f [f [f \ldots$. The unifier forbids this and fails. When the unifier fails, it throws to the nearest containing joshua:with-unification. Thus the function above returns :occur-check-forbids.

(test-occur) -> :occur-check-forbids

Why should Joshua attempt to avoid creating such circular structures, though? (The check does have a cost in performance, which is why most versions of Prolog won’t do it.) The answer is that if it were permitted, certain incorrect inferences could be made. Here’s an example. Suppose we have a predicate is-parent-of, which takes two people as arguments:

(define-predicate is-parent-of (kid parent))

This means that parent is a parent of kid. We can then make the (unsurprising) statement that every person has a parent:

\[ \forall x \exists y : \text{has-parent}(x, y) \]

or, in quantifier-free language,

[has-parent ?x (p ?x)]

where $p$ is the Skolem function for the existential variable $y$. (You can think of it as a notation for finding the parent of its argument.)

Now try to unify the above statement with [has-parent ?z ?z]. In the absence of the occur check, we get the equations:

?z = ?x

and

?z = (p ?x)

(This would end up with $?x = (p ?x) = (p (p (p \ldots$. Now substitute for the arguments in $[p ?z ?z]$ using those equations, to get:

[has-parent (p ?x) ?x]

which is just the original statement with the arguments reversed. This is unsound. It is not justifiable to infer that has-parent is a symmetric predicate. (Indeed, it is not, since no one is his own parent!) Thus, to be sound, Joshua must forbid occur-check-type matches.
Related Functions:

  joshua:with-unification
  joshua:succeed

See the section "Pattern Matching in Joshua: Unification", page Pattern-Matching-in-Joshua:-Unification-SECTION.

joshua:unjustify database-predication &optional justification  Generic Function
Removes a justification from a predication in the database. For example, if you joshuatell predication and then later change your mind about it, you can use joshua:unjustify to remove justification from the possible supports. This does not automatically remove all support for database-predication, as there might be other justifications for it as well.

database-predication A predication object that is in the database. Must be the actual database object, and not a copy of it.

justification Specifies the justification to be removed. If justification is not supplied, implementations of joshua:unjustify should default it to the justification currently being used to support database-predication.

In general, joshua:unjustify is useful only if database-predication is built on some model that supports the TMS protocol.

In the default (non-TMS) Joshua model, joshua:unjustify just sets the truth-value of its argument to joshua:*unknown*.

Examples:

When Prospero is reconciled to his countrymen, he will cast the following spell:

(map-over-database-predications [is-exiled-from Prospero ?] #’unjustify)

(map-over-database-predications [is-exiled-from Miranda ?] #’unjustify)

(map-over-database-predications [is-friend-of Miranda Caliban] #’unjustify)

joshua:unjustify and joshua:untell work in similar fashion, but with very different results. See the generic function joshua:untell, page joshua:untell-FUNCTION. joshua:unjustify keeps the unjustified fact in the database. If the fact is later given again to joshuatell, it is not considered as a new predication, but rather as a variant of an existing one, and no forward rules are run.

joshua:untell, on the other hand, actually removes the fact from the database, freeing up storage, and causing the database to lose previous knowledge of it; if the fact is later given to joshuatell again, it is considered as a new fact, and forward rules are rerun.

Related Functions:
joshua:untell
joshua:uninsert
See the section "Revising Program Beliefs", page Revising-Program-Beliefs-SECTION.

See the section "Retracting Predications with joshua:unjustify", page Retracting-Predications-with-JOSHUA:UNJUSTIFY-SECTION.

joshua:*unknown*  Variable
A named constant used by Joshua to denote a truth value of joshua:*unknown*. You can compare truth values using eql.

A predication is joshua:*unknown* when there is no valid reason that supports it. The predication may or may not remain in the database, but is conceptually "not seen" until its truth value changes to joshua:*true* or joshua:*false*.

Related Topics:
- joshua:*true*
- joshua:*false*
- joshua:*contradictory*
- joshua:truth-value
- joshua:predication-truth-value

See the section "Truth Values", page Truth-Values-SECTION.

joshua:untell database-predication  Generic Function
Removes a single predication from the database, clearing up storage space.
(This function is a dual of joshuatell, which adds a predication to the database.)

database-predication A predication. Must be the actual predication object that is in the database, not a copy of it.

joshua:untell first calls joshua:unjustify to make the fact no longer believed (joshua:*unknown*), clears some internal caches, then calls joshua:uninsert to remove the fact from the database. The surgical properties of joshua:untell in actually removing the predication as opposed to only removing its justification have two effects:

1. Some storage may become garbage-collectible. This can lower the virtual-memory requirements of your program. Of course, you pay for it by doing the extra work of joshua:uninsert.

2. The predication is no longer in the database. This means that if you re-joshuatell it, joshuatell returns a second value of T, denoting it has never seen this predication before; in consequence, joshuatell will also run forward rules again.

(If, on the other hand, you merely joshua:unjustify the predication, then
joshuatell it once again, joshuatell returns a second value of nil, denoting
the predication already existed in the database; joshuatell does not run
forward rules when an existing predication is retold.) However, if a TMS is
present, the consequences of running those rules will be brought back in.

Examples:

(define-predicate has-eye-color (creature color))

(tell [and [has-eye-color cat green]
       [has-eye-color rat black]])

Show Joshua Database
True things
[HAS-EYE-COLOR CAT GREEN]
[HAS-EYE-COLOR RAT BLACK]
False things
None

;; untell a predication by clicking left on it in the database display
(untell [HAS-EYE-COLOR CAT GREEN])
NIL

Show Joshua Database (matching pattern [default All]) All
True things
[HAS-EYE-COLOR RAT BLACK]
False things
None

;; untell using the predication object returned as the query support
(ask [has-eye-color rat black]
     #'(lambda (backward-support)
         (untell (ask-database-predication backward-support)))
     :do-backward-rules nil)

Show Joshua Database (matching pattern [default All]) All
True things
None
False things
None

Note that in the last example above you probably should have used
(map-over-database-predications [has-eye-color rat black] #'untell)

Compare the following examples to see the difference between
joshua:untell and joshua:unjustify.

(define-predicate is-uncle-of (uncle niece-or-nephew) (ltms:ltms-predicate-model))
(define-predicate is-nephew-of (nephew uncle) (ltms:ltms-predicate-model))
(defrule notice-uncles (:forward)
  if [is-uncle-of ?uncle ?nephew]
  then [and (format t "%I note that -A is the uncle of -A." ?uncle ?nephew)
            [is-nephew-of ?nephew ?uncle)])

First we'll joshuatell an avuncular fact, joshua:untell it, and then
re-joshuatell it. After the first joshuatell the fact fires the forward rule.
After the second joshuatell the forward rule fires again, since joshuatell
sees the predication as T.

  (setq canonicalized-uncle-fact (tell [is-uncle-of Judah Manasseh]))
I note that JUDAH is the uncle of MANASSEH.
[IS-UNCLE-OF JUDAH MANASSEH]
T

Show Joshua Database
True things
  [IS-UNCLE-OF JUDAH MANASSEH]
  [IS-NEPHEW-OF MANASSEH JUDAH]
False things
   None

(untell canonicalized-uncle-fact)

Show Joshua Database
True things
   None
False things
   None

(tell [is-uncle-of Judah Manasseh]); this fires the rule again!
I note that JUDAH is the uncle of MANASSEH.
[IS-UNCLE-OF JUDAH MANASSEH]
T

Now we'll use a variation of this example.

We start with the fact we just entered in the database above and which
fired the forward rule. Now we joshua:unjustify the fact and then
joshuatell it again.

After the joshua:unjustify, the fact changes its truth value from
joshua:*true* to joshua:*unknown*, but remains in the database. When we
joshuatell the fact once again, its truth value changes from
joshua:*unknown* to joshua:*true*, but joshuatell already knows about
the fact, and no forward rules fire. Note, however, that the TMS brings the
is-nephew-of deduction back in. We can tell it does so without re-executing
the rule, since the side-effect (the format message) in the rule-body did not
recur.
Show Joshua Database
  True things
  [IS-UNCLE-OF JUDAH MANASSEH]
  [IS-NEPHEW-OF MANASSEH JUDAH]
  False things
  None

  (unjustify [IS-UNCLE-OF JUDAH MANASSEH])
  NIL

Show Joshua Database
  True things
  None
  False things
  None

  (tell [is-uncle-of Judah Manasseh])
  ; tell knows this fact is old, and it doesn’t rerun the forward rule
  [IS-UNCLE-OF JUDAH MANASSEH]
  NIL

Show Joshua Database
  True things
  [IS-UNCLE-OF JUDAH MANASSEH]
  [IS-NEPHEW-OF MANASSEH JUDAH]
  False things
  None

In sum, joshua:unjustify and joshua:untell do similar things, but with significant differences. If you want to change your mind about believing a fact but reserve your right to return to that fact later, you probably want to use joshua:unjustify. If, on the other hand:

- You just did a scratch calculation and want to flush it now that you have the answer, or
- You want the storage back, or
- You don’t intend to come back and raise the issue of re-running rules.

you probably want to use joshua:untell.

Related Functions:
  joshua:tel1
  joshua:unjustify
   "Clear Joshua Database Command"

See the section "Removing Predications From the Database", page Removing-Predications-From-the-Database-SECTION.
See the section "The Joshua Database Protocol" in *Joshua Reference Manual*.

See the section "Customizing the Data Index" in *Joshua Reference Manual*.

**joshua:variant** object1 object2  

Function

Two predications that differ *only in the names* of the logic variables they contain are equivalent, and are said to be *variants* of each other.

The function **joshua:variant** checks whether two objects are *variants* of each other. If they are, it returns **t**, otherwise **nil**.

When **joshua:tell** has to add a predication to the database it uses **joshua:variant** to determine if the predication is already there.

object1 A predication

object2 Another predication

"Variant" means there is a *renaming* of variables that makes one variable look like the other. For example:

```
(define-predicate foo (object))

(variant [foo 1 ?x] [foo 1 ?y])
T ; you can rename ?x → ?y.
```

```
(variant [father ?x ?y] [father ?a ?b])
T
```

**joshua:variant** should not be confused with **joshua:unify**. The latter tries to see if two objects can be *made to be* the same. **joshua:variant** checks if they *are* the same. It doesn’t ever *set* logic variables, but merely looks for a renaming. **joshua:variant** is based on the notion that it should not matter what the names of logic variables are, so long as the structures are the same. This is a much stronger condition than **joshua:unify**. Thus, every pair that satisfies **joshua:variant** also satisfies **joshua:unify**, but not vice versa.

```
(variant [foo 1 ?x] [foo 1 bar])
NIL ; these unify, but are not variants
 ; variables cannot be renamed
```

```
(variant [father ?x ?x] [father ?a ?b])
NIL
```

**joshua:variant** also works on other structures such as lists.

Examples:

```
(variant 'a 'a)
T
```
(variant '([foo baz] [foo bar]) '([foo baz] [foo bar]))

T

See the section "Variables and Scoping in Joshua", page Variables-and-
Scoping-in-Joshua-SECTION.

**joshua:with-statement-destructured arglist predication &body**

*Macro*

body

Provides access to the `arglist` of `predication`. Wrap this macro around a `body`
of code within methods in which you want to refer to the arguments of a
predication that are not already in instance variables. (This macro works
outside of methods, too.)

**arglist**

The argument list of the specified `predication`. This can
be anything suitable for deconstructing-bind.

**predication**

A Joshua predication.

For example, inside a `joshua:say` method for the predication `foo`:

(define-predicate enough-already (number-of servings food))

(define-predicate-method (say enough-already)
  (optional (stream *standard-output*))
  (with-statement-destructured (how-many servings food) self
    (format stream "{-% You’ve just had ^A ^A of ^A. Hadn’t you better quit?"
     how-many servings food)))

(say [enough-already 5 platters pickled-pigs-feet])

You’ve just had 5 PLATTERS of PICKLED-PIGS-FEET. Hadn’t you better quit?
NIL

Related Functions:

  joshua:define-predicate

**joshua:with-unbound-logic-variables variable-list &body body**

*Macro*

This macro provides a way to generate a set of logic variables for use in
code. Each (Lisp) variable within the `variable-list` is bound within the scope
of the macro to a distinct, non-unified logic variable within the `body` of the
macro. In essence a Lisp variable in `variable-list` has as its Lisp value a log-
ic variable, for the duration of `body`.

**variable-list**

Is a list of variables

**body**

Is any lisp form

Example:
The predicate presidential-candidate is defined in the following example. The macro is used to temporarily set anybody to be a logic variable. Then two predications are compared to see if they unify with one another. Unification occurs in this case so the format statement prints the value of anybody.

    (define-predicate presidential-candidate (somone))

    (with-unbound-logic-variables (anybody)
      (with-unification
        (unify '[presidential-candidate ,anybody] [presidential-candidate Abe])
        (format t "&The value of anybody is ~s." anybody)))

The value of anybody is ABE.
NIL

**joshua:with-unification** &body body

*Macro*

Establishes the scope within which substitutions specified by the **joshua:unify** function take effect. This temporary unifying mechanism is useful within Lisp code in the body of Joshua rules, since it lets the programmer try out a variety of different matching options.

Whenever unification fails, **joshua:unify** goes to the end of the dynamically innermost **joshua:with-unification** and undoes all the bindings established so far.

Thus, **joshua:with-unification** establishes both of the following:

- The scope of unifications done in its body
- A place to be thrown to if a unification in its body fails

**Examples:**

    (define-predicate candidate-word (a-word))
    (define-predicate is-computer-jargon (some-word))
    (defvar *computer-jargon* '(foo bar baz quux))
(defrule jargon-finder (:backward)
  IF (typecase ?candidate-word
       (unbound-logic-variable
         (loop for word in *computer-jargon*
               doing (with-unification
                       (unify ?candidate-word word)
                       (succeed)))))
  (otherwise
   (member ?candidate-word *computer-jargon*))
  THEN [is-computer-jargon ?candidate-word])

(ask [is-computer-jargon ?x] #'print-query)
[IS-COMPUTER-JARGON FOO]
[IS-COMPUTER-JARGON BAR]
[IS-COMPUTER-JARGON BAZ]
[IS-COMPUTER-JARGON QUUX]

Related Function:
  joshua:unify

See the section "Pattern Matching in Joshua: Unification", page Pattern-Matching-in-Joshua-Unification-SECTION.
Appendix A
A Figure of the Joshua Protocol of Inference
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