

Problem Set 3 Solutions

Due: Thursday, March 11, 2021

Problem 3.1 [Thinking Outside The Box]. The orthogonal range query data structure described in Lecture 4 supports the following operation:

inside(**p**, **q**): return all *k* points *inside* the bounding box spanned by points **p** and **q**.

Implement the following operation using $O(d)$ calls to *inside*():

outside(**p**, **q**): return all *k* points *outside* the bounding box spanned by points **p** and **q**.

Other than the black-box calls to *inside*(), the running time of your algorithm should be $O(d + k)$.

Solution: The idea is to subdivide the space outside the query range into $O(d)$ disjoint ranges, query each of them, and combine the results.

Without loss of generality, $p_i < q_i$ for all i .

For each dimension, we can partition the space into 3 parts: $L_i = \{x \in \mathbb{R}^d \mid x_i < p_i\}$, $C_i = \{x \in \mathbb{R}^d \mid p_i \leq x_i \leq q_i\}$, and $R_i = \{x \in \mathbb{R}^d \mid x_i > q_i\}$. Then we can write the exterior as

$$\begin{aligned} \mathbb{R}^d \setminus (C_1 \cap \dots \cap C_d) &= L_1 \sqcup R_1 \\ &\sqcup (C_1 \cap L_2) \sqcup (C_1 \cap R_2) \\ &\sqcup (C_1 \cap C_2 \cap L_3) \sqcup (C_1 \cap C_2 \cap R_3) \\ &\vdots \\ &\sqcup (C_1 \cap \dots \cap C_{d-1} \cap L_d) \sqcup (C_1 \cap \dots \cap C_{d-1} \cap R_d). \end{aligned}$$

Each of these $2d$ terms is an orthogonal range. For example, $C_1 \cap L_2 = [p_1, q_1] \times (-\infty, p_2) \times \mathbb{R}^{d-2}$. We call *inside*() on each these ranges in order and concatenate the results.

Creating the first range takes $O(d)$ time. Each successive range differs from the previous range in only $O(1)$ coordinates, so we can update it in $O(1)$ time. Concatenating the results takes $O(k)$ time. So the total running time outside the $2d$ calls to *inside*() is $O(d + k)$.

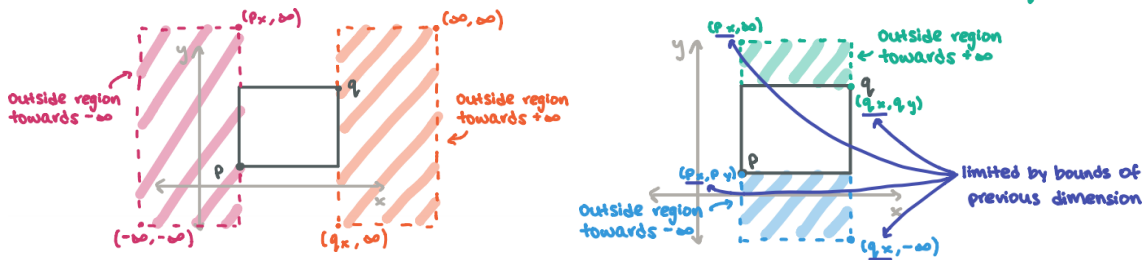


Figure 1: $d = 2$ example. [Figure by Shana Mathew]