

# Student Presentations II

Monday, May 14

Robert Bryant & Kevin Wang: Avuncular Fun

- given set  $S$  of nodes, root  $\in S$
- build tree on  $S$  to minimize  $\sum_{v \in T} \text{depth}(v) \cdot \text{key}(v)$

**NEW  
PROBLEM**

where  $\text{key}(v) = \# \text{ children } \oplus v$

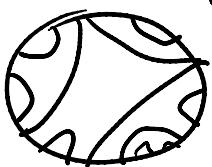
or arbitrary pos. int. if  
 $\# \text{ children}$  is fixed

- Avuncular property:  
 $\forall v, w : \text{key}(v) \leq \text{key}(w) \Rightarrow \text{depth}(v) \leq \text{depth}(w)$
- Problem 1: fixed branching factor
  - $O(\lg^2 n)$  w.c. insertion/deletion to  $S$
  - $O(1)$  space/node
  - $O(\lg n \lg \lg n)$  amortized via indirection
  - $O(\lg n)$  via clever indirection
- Problem 2: variable branching factor
  - seems difficult because of holes
  - lower bound?

**RESEARCH**

## Katherine Lai: Cyclic Union-Split-Find

- union-find:  $\Theta(\alpha(n))$
- union-split-find = predecessor with  $u=n$   
=  $\Theta(d \lg \lg n)$   
↳ on intervals
- cyclic union-split-find:  
chords in polygon or balanced parens.



$(( )(( )(( ))))$

- split = insert
- union = delete
- find = faces on either side of edge
- known:  $\Omega(\log_w n)$
- NEW:  $O(\log_w n)$
- based on strongly weight-balanced B-tree with branching factor  $w^\varepsilon$ .

**RESEARCH**

## Tural Badirkhanli: Worm Detection DSs

- worms are fast (minutes!)
- need automated detection/response
- content-based detection:  
  - FPGA to compute hash on 80-bit sliding window of streaming data
- signature over random combinations  
  ⇒ worm "fingerprinting"
- Rabin-Karp fingerprinting
- bitmap DS to decrease memory & accuracy

## Galen Pickard: Gov't R&D on MST Algs.

- Lincoln Labs: Stellar: system for intrusion via hierarchical clustering
- essentially MST, but dynamic
- $O(n^2)$  via trivial rebuilding of MSTs
- $O(n^{1.2})$  average case [prior work]
- $O(n \lg n)$  via Holm et al. dynamic MST

## IMPLEMENTATION

## Anders Kaseorg: Fully Persistent Arrays

= fully persistent RAM DS

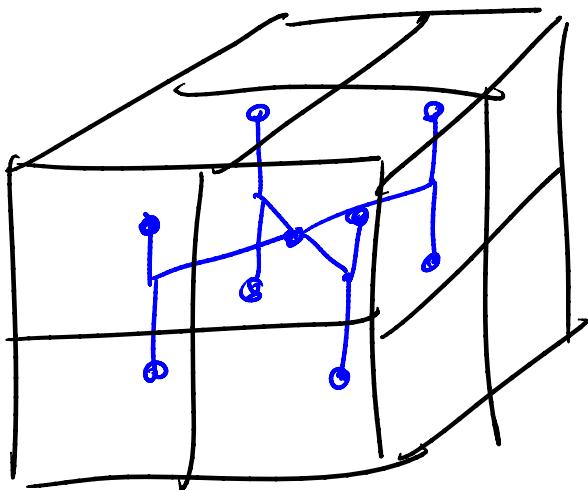
- $O(\lg \lg m)$  lookup,  $O(\lg \lg m)$  expected amortized,  $O(m)$  space [Dietz-WADS 1989]
- Euler tour on version tree
- order query DS on Euler tour,  $O(1)/\text{op.}$
- vFB or  $\gamma$ -fast trees for predecessor on (index, version) pairs  $\Rightarrow O(\lg \lg m)$
- order query DS changes tags which incurs updates in pred. ~ouch
- both order query &  $\gamma$ -fast DSs use indirection with  $\lg n$  chunks @ bottom
- mix tops & bottoms separately

**READING**

## Hui Tang: sculpture

- 8-way branching factor B-tree, depth 4  
 $\approx 40$  hours, 150 ft dowels  
 $\approx \$150$

- star of 4  
from center.  
It's to connect  
to centers of  
8 subcubes



# Mashhood Ishaque: Half-plane Proximity Queries

- preprocess set of points
- query: given point & half-plane,  
find nearest point in half-plane
- want:  $\tilde{O}(n^2)$  preproc. & space ] for diameter  
polylog query MST
- known:
  - simplicial partitioning [Daescu et al. - CGTA 2005]  
 $\tilde{O}(n)$  preproc. & space  
 $O(n^{1/2+\epsilon})$  query
  - points in convex position:
    - $\tilde{O}(n \lg n)$  preproc. & space } [ibid]
    - $O(\lg^2 n)$  query }
    - [poly. preproc.,  $\tilde{O}(n)$  space, } Aronov  
 $O(\lg n)$  query } et al.
  - idea: look at dual arrangement  
 $\Rightarrow \tilde{O}(n^3)$  preproc. & space,  
 $O(\lg n)$  query (general pt. set)
  - minor improvements so far;  
hope to get down to  $\tilde{O}(n^2)$ .

**RESEARCH**