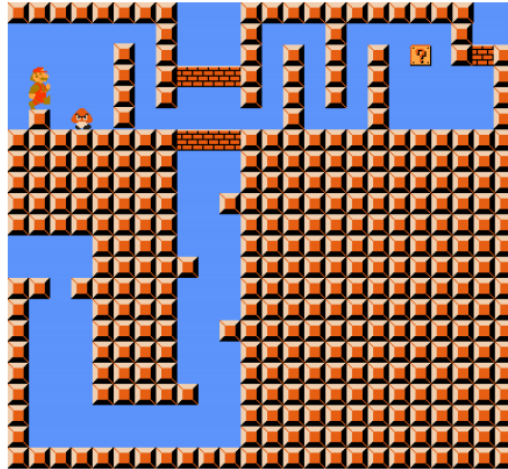


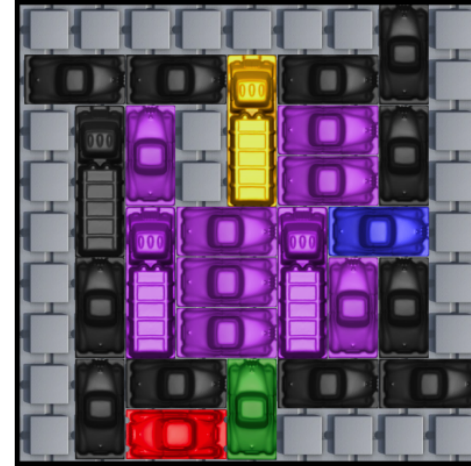
# ALGORITHMIC LOWER BOUNDS: FUN WITH HARDNESS PROOFS

## Super Mario Bros.



Crossover gadget for NP-hardness

## Rush Hour



AND gadget for PSPACE-hardness

## Minesweeper



OR gadget for NP-hardness

## Hardness Made Easy\*

Learn **when to give up** the search for efficient algorithms; see **connections** between computational problems; **solve puzzles** to prove theorems, solve **open problems**, and write papers.

*Topics:* NP, PSPACE, EXPTIME, EXPTIME, EXPSPACE, 3SUM, approximation, fixed parameter, games & puzzles, 3SUM, key problems, gadgets, and proof styles.

Fall 2014



6.890 taught by Professor Erik Demaine

Grad H, AUS, and Theoretical CS Concentration  
Tuesday & Thursday 3:30-5:00pm in room 2-105

<http://courses.csail.mit.edu/6.890/>  
sign up for our mailing list to join the class

*\*Easiness not guaranteed. Side effects such as open problems and a heightened sense of complexity may occur. Ask your advisor if 6.890 is right for you!*

6.890

Lecture 1

Sept. 4, 2014

6.890: Algorithmic Lower Bounds  
/ Fun with Hardness Proofs

"Hardness  
made Easy"

Prof. Erik Demaine

TAs: Sarah Eisenstat & Jayson Lynch

<http://courses.csail.mit.edu/6.890/fall14/>

What is this class?

- practical guide to proving computational problems are formally hard / intractable
- NOT a complexity course  
(but we will use/refer to needed results)
- (anti)algorithmic perspective

Why take this class?

- know your limits in algorithmic design
- master techniques for proving hardness
- cool connections between problems
- fun problems like Mario & Tetris  
(serious problems too)
- solve puzzles → publishable papers

key problems  
- proof styles  
gadgets

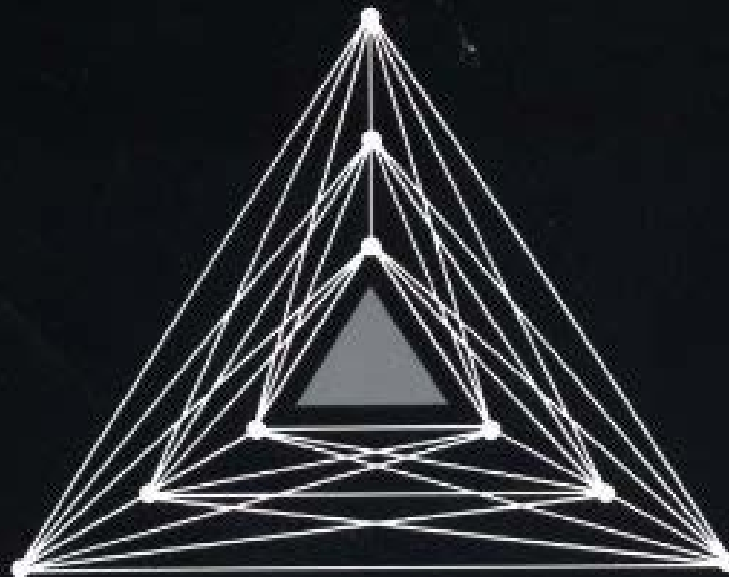
Background: algorithms, asymptotics, combinatorics

- no complexity background needed  
(but also little overlap with a complexity class)

# COMPUTERS AND INTRACTABILITY

## A Guide to the Theory of NP-Completeness

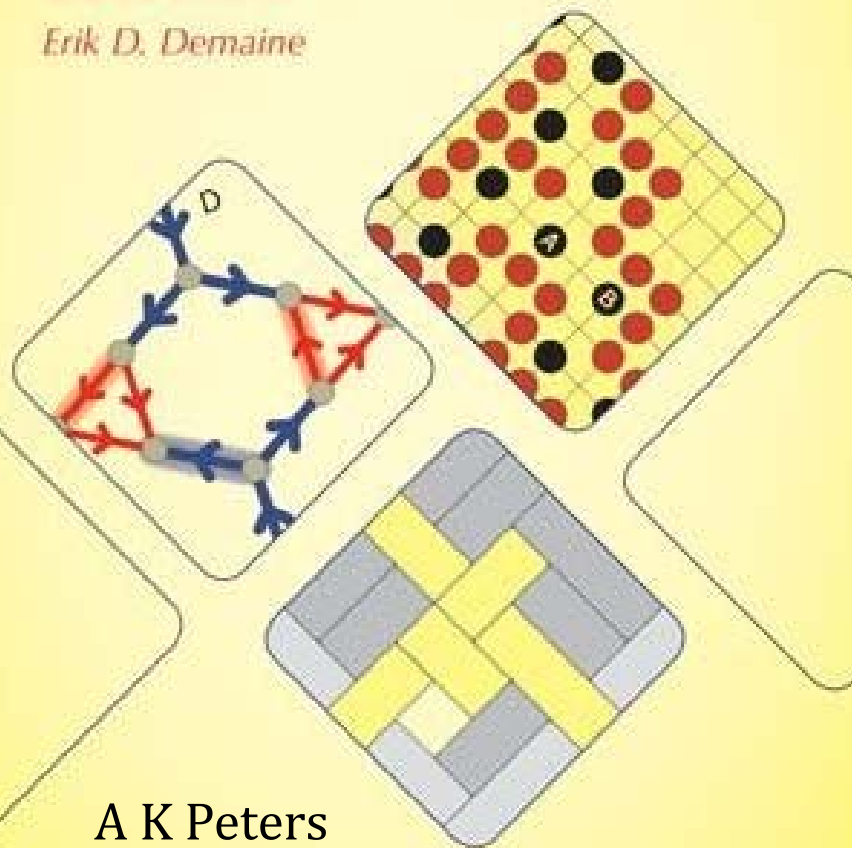
Michael R. Garey / David S. Johnson



W. H. Freeman  
1979

# Games, Puzzles, & Computation

Robert A. Hearn  
Erik D. Demaine



A K Peters  
July 2009

# ゲームと パズルの 計算量

Games, Puzzles,  
& Computation

Robert A. Hearn  
Erik D. Demaine

ロバート・A・ハーン  
エリック・D・ドメイン  
上原隆平



translated by  
Ryuhei Uehara

近代科学社

# Kotaku

TOP STORIES

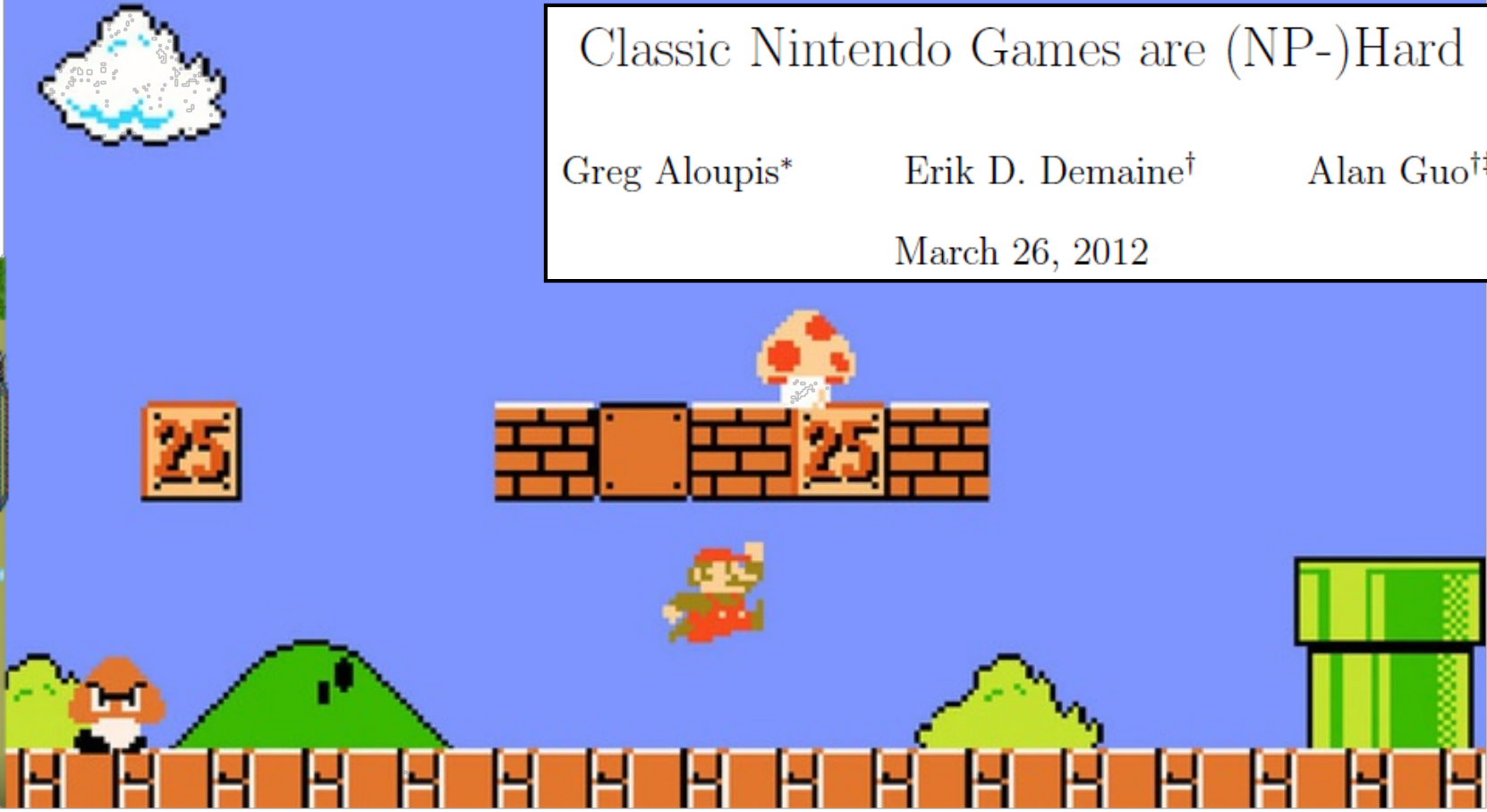
## Classic Nintendo Games are (NP-)Hard

Greg Aloupis\*

Erik D. Demaine†

Alan Guo††

March 26, 2012



HEROES

**TOTAL RECALL**

Mon. - Fri.  
11PM - Mid.  
(EASTERN)

NINTENDO

## Science Proves Old Video Games Were Super Hard

BY LUKE PLUNKETT +

MAR 12, 2012 11:00 PM

Share

g +1

f Like

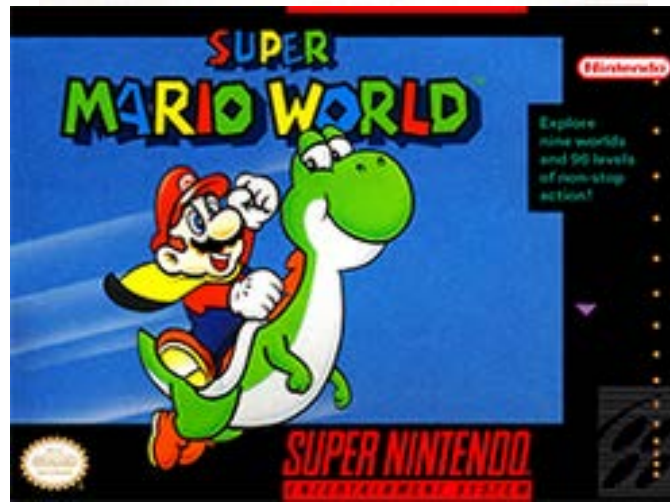
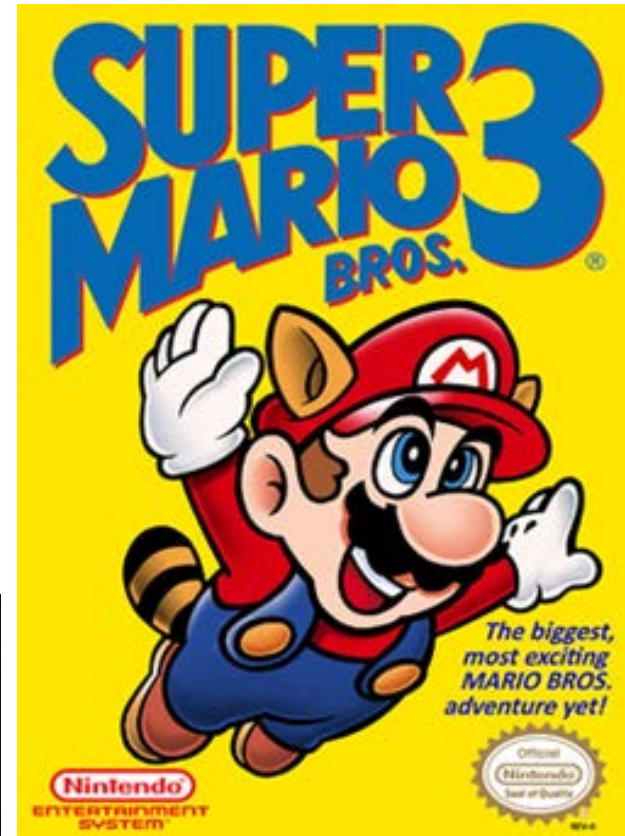
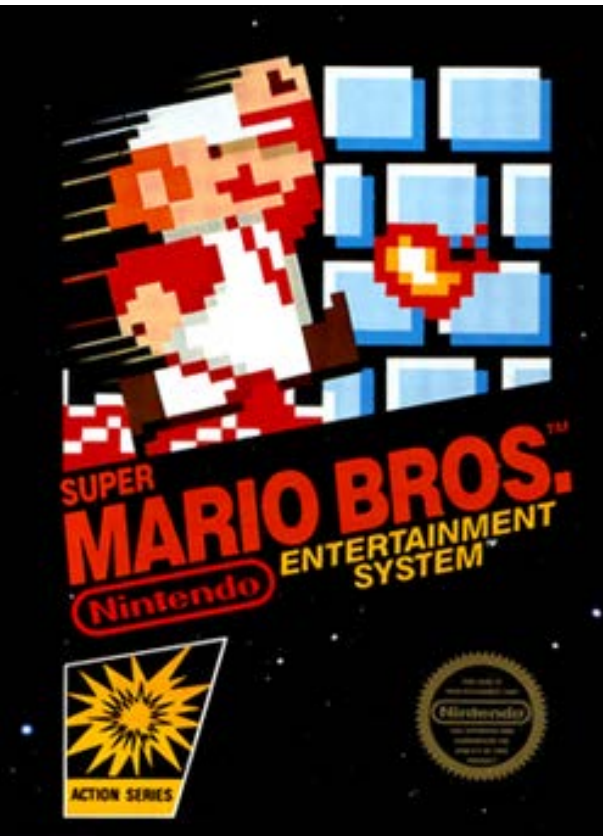
411

33,867 232

# Super Mario Bros. is NP-Hard

[Aloupis, Demaine, Guo 2012]

The Lost Levels

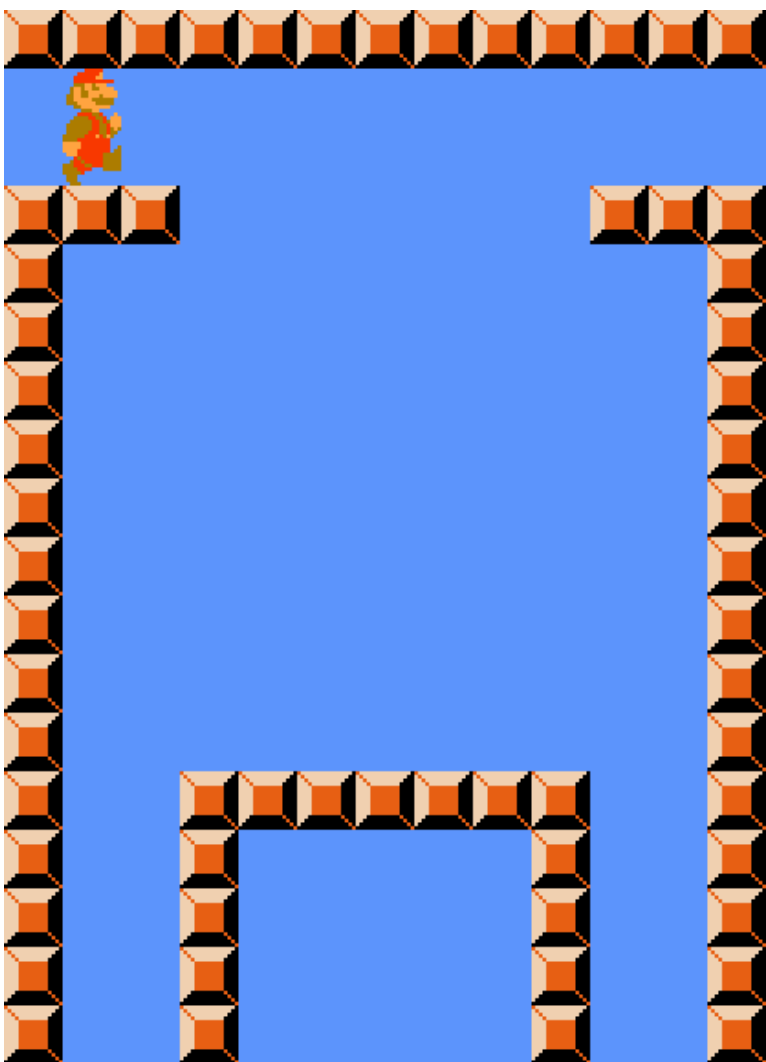




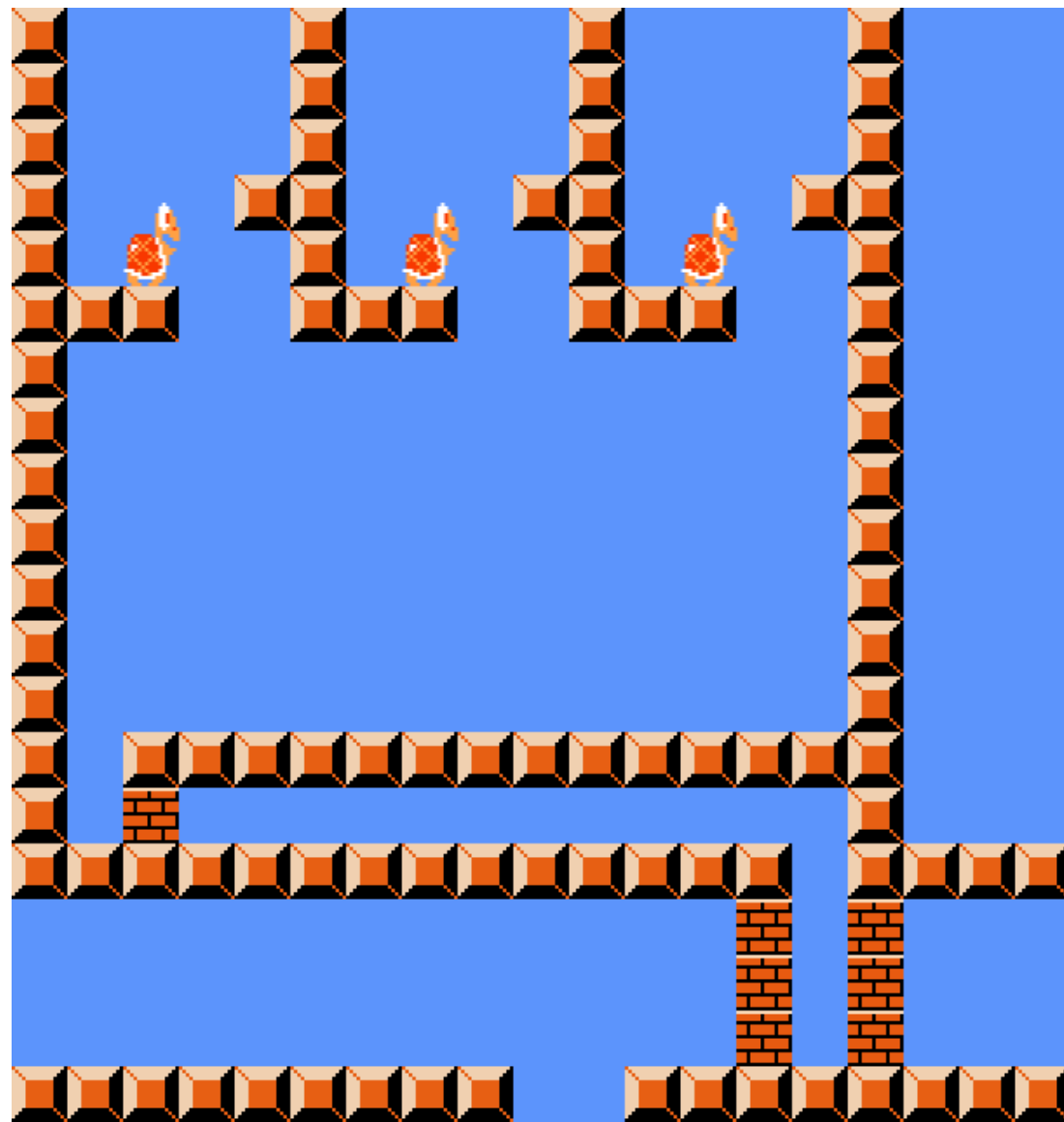
# Super Mario Bros. is NP-Hard

[Aloupis, Demaine, Guo 2012]

**clause**



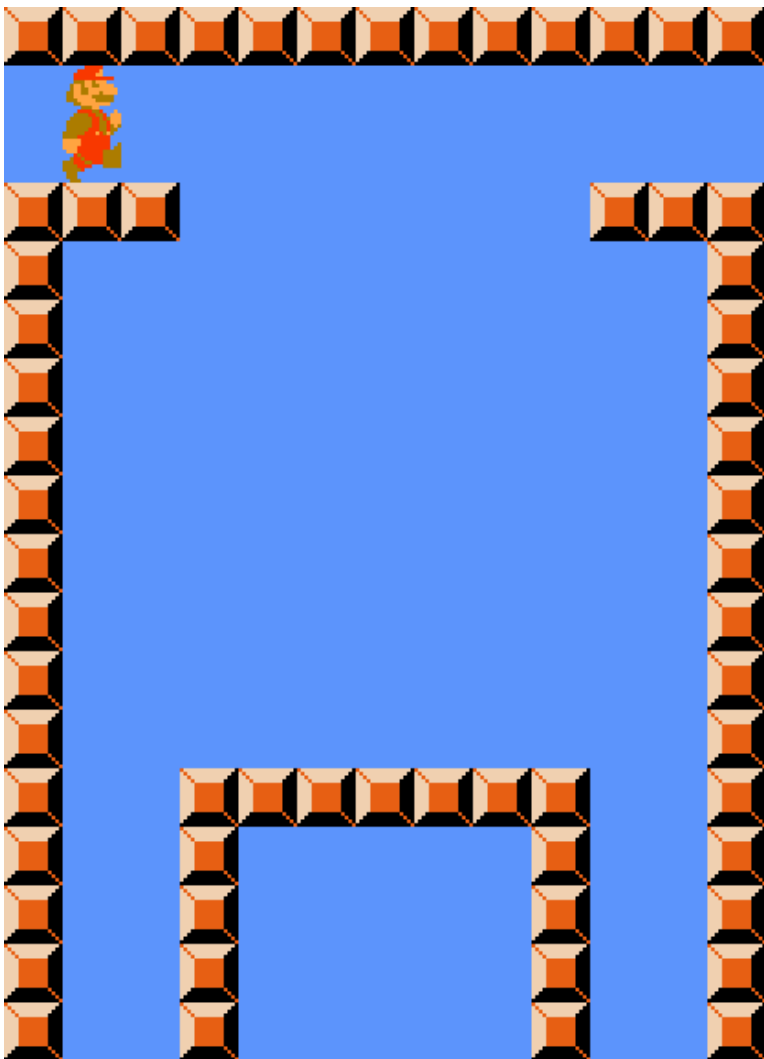
**variable**



# Super Mario Bros. is NP-Hard

[Aloupis, Demaine, Guo 2012]

clause



variable



*Super Mario Bros. 3*

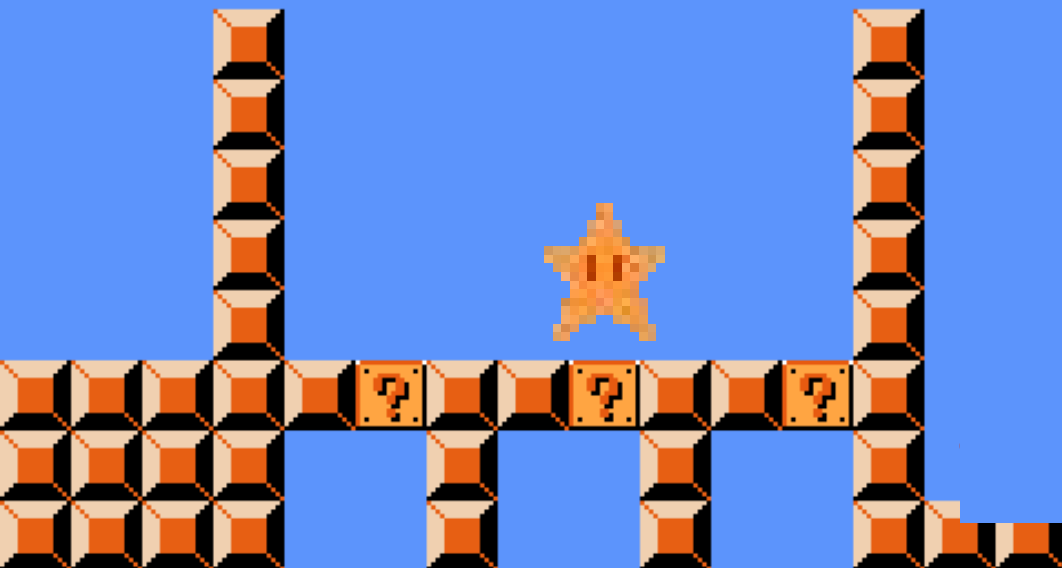




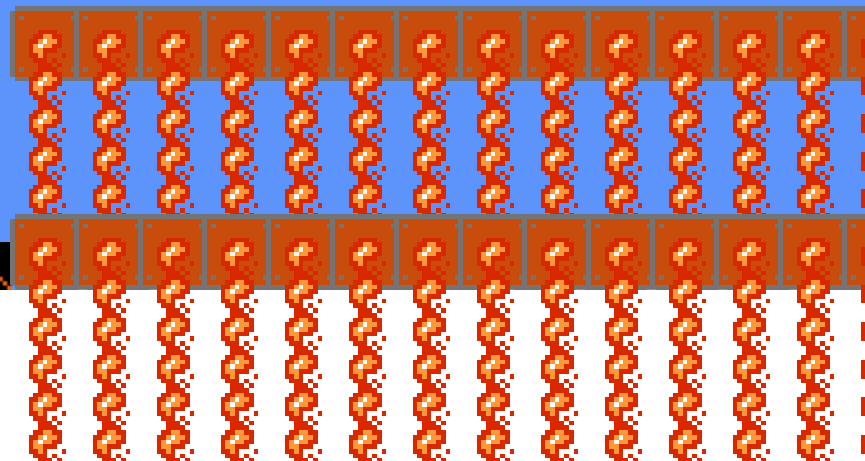
# Super Mario Bros. is NP-Hard

[Aloupis, Demaine, Guo, Viglietta 2014]

**clause**

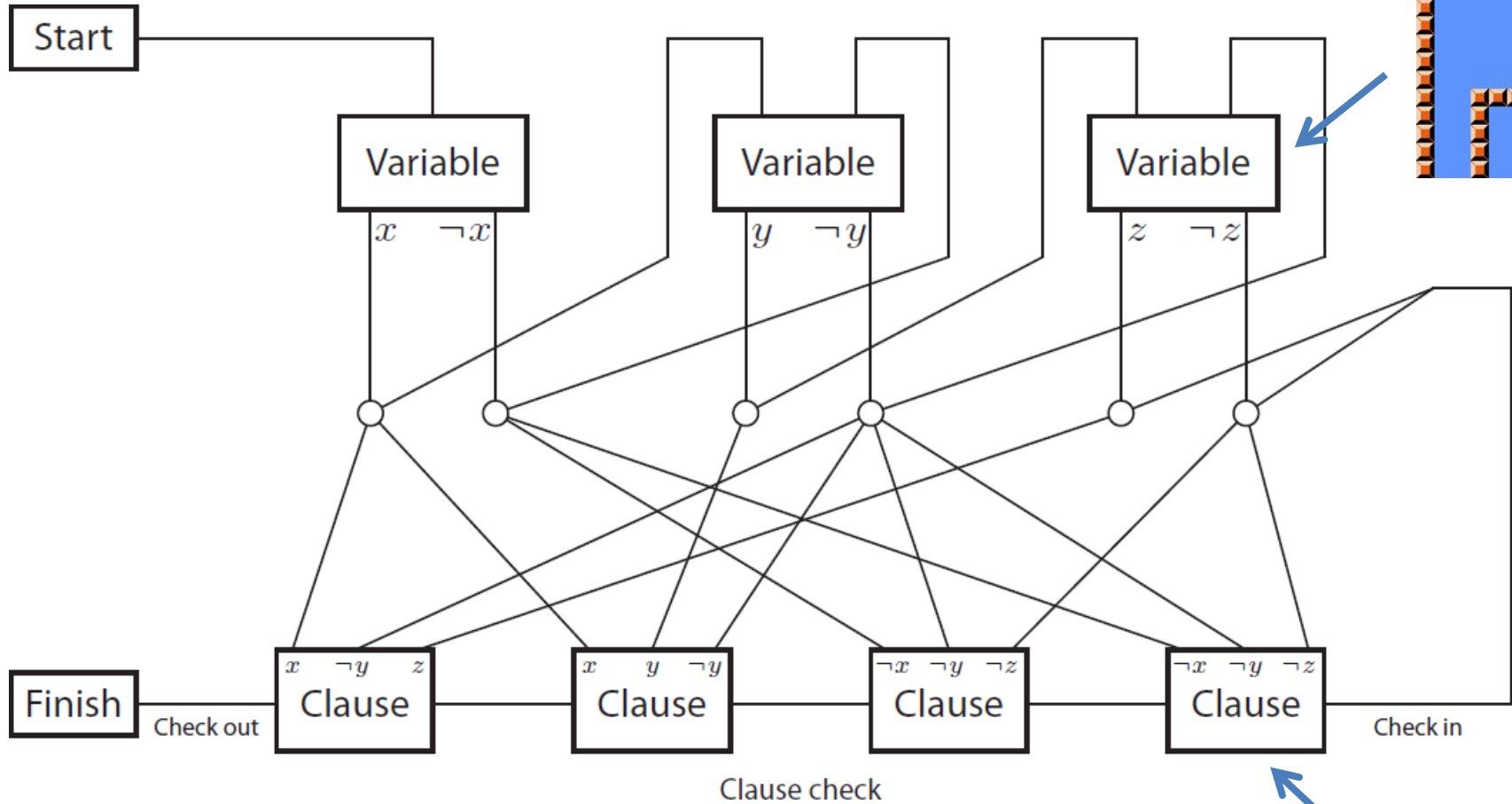


*Super Mario Bros.*



# Super Mario Bros. is NP-Hard

[Aloupis, Demaine, Guo, Viglietta 2014]



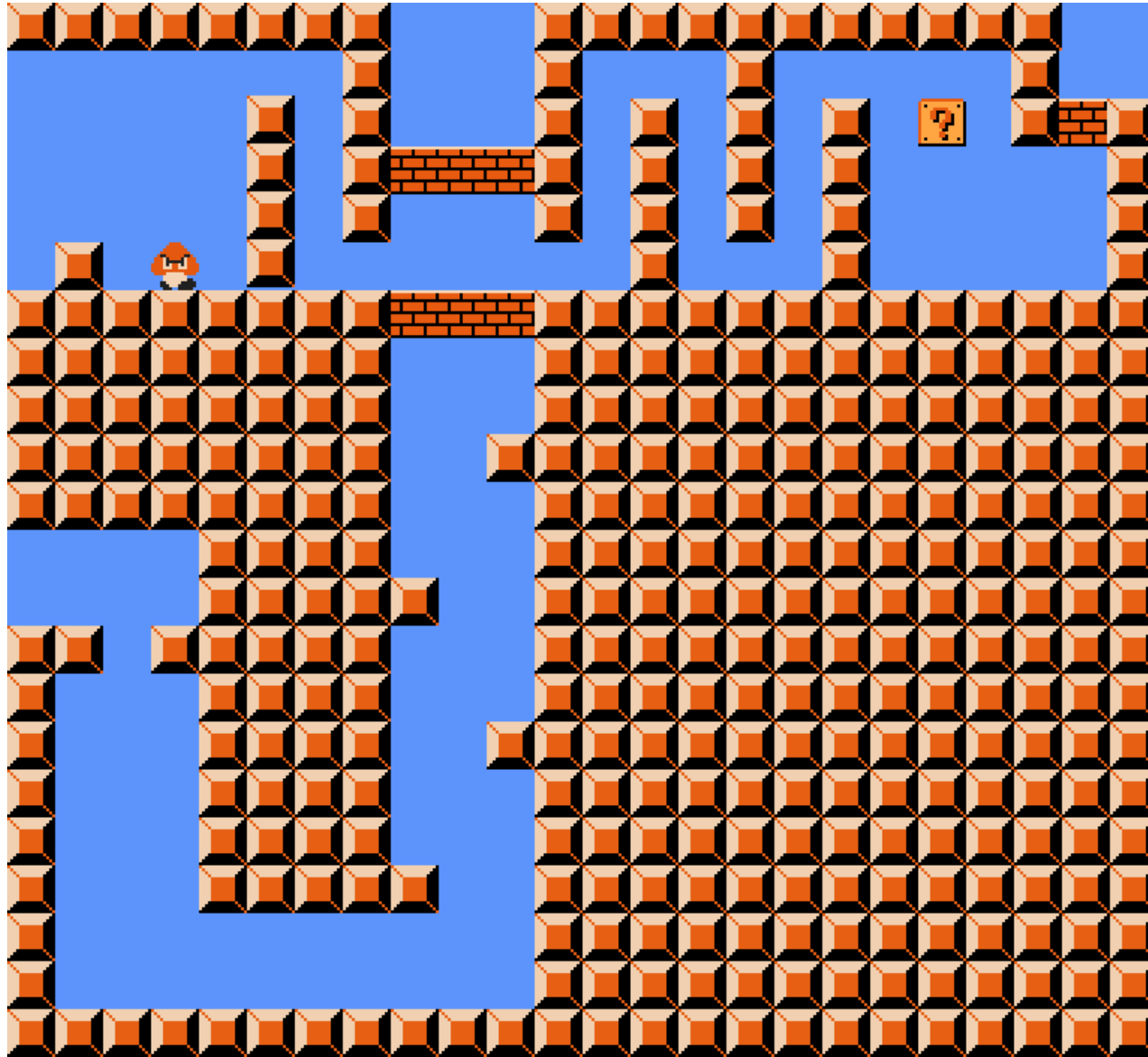
$(x \text{ OR } \neg y \text{ OR } z) \& (x \text{ OR } y \text{ OR } \neg y) \&$   
 $(\neg x \text{ OR } \neg y \text{ OR } \neg z) \& (\neg x \text{ OR } \neg y \text{ OR } \neg z)$





# Super Mario Bros. is NP-Hard

[Aloupis, Demaine, Guo, Viglietta 2014]



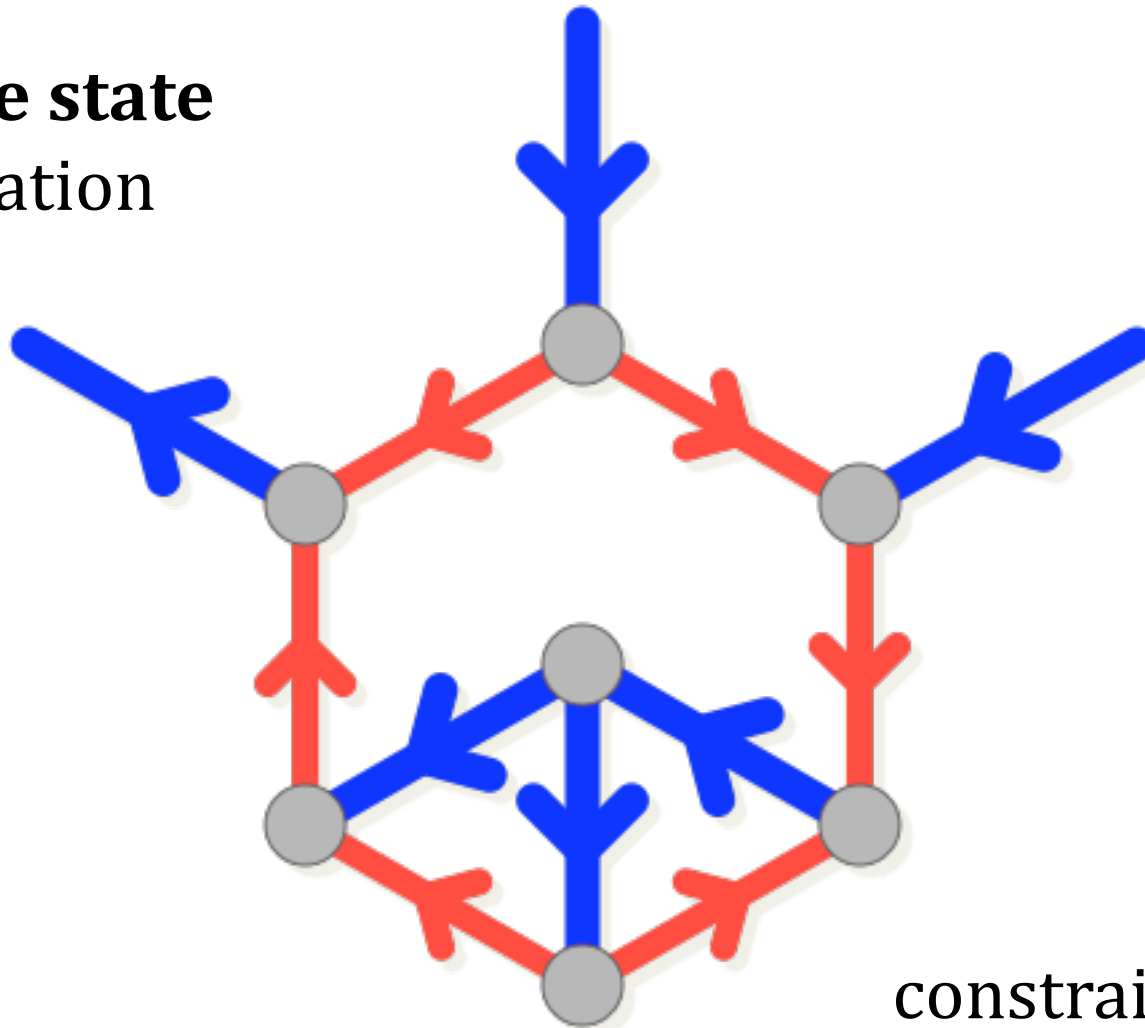
crossover





# Constraint Graphs

**Machine state**  
= orientation



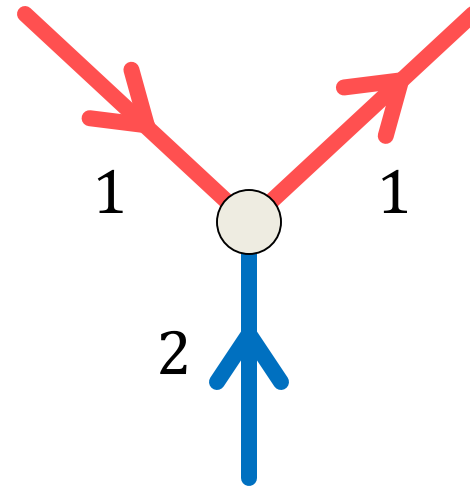
constraint graph



# Constraint Logic

— = 1

— = 2

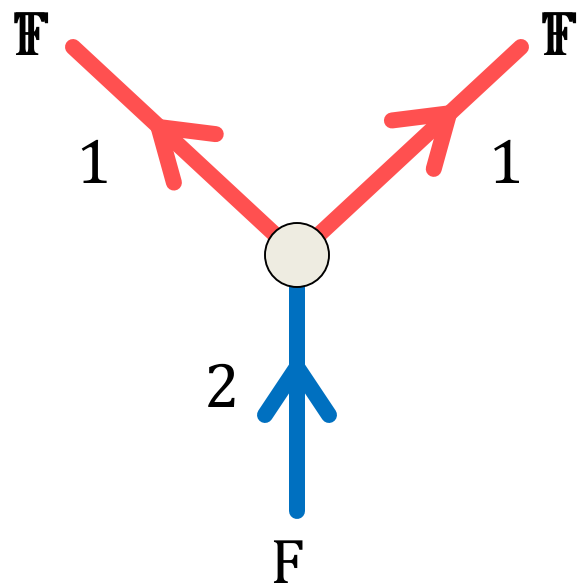


**Rule:** at least 2 units incoming at a vertex

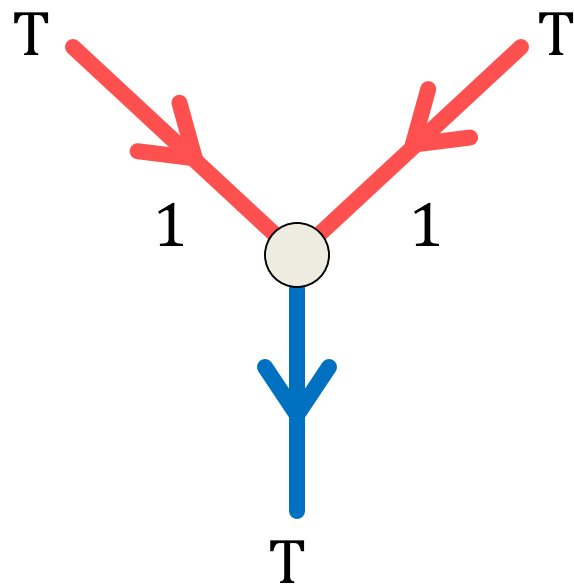
**Move:** reverse an edge, preserving Rule



# AND vertex



not your usual  
AND gate!



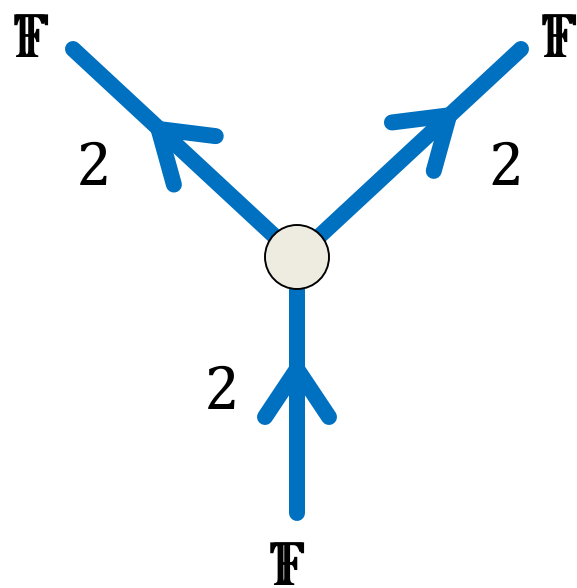
inputs

output

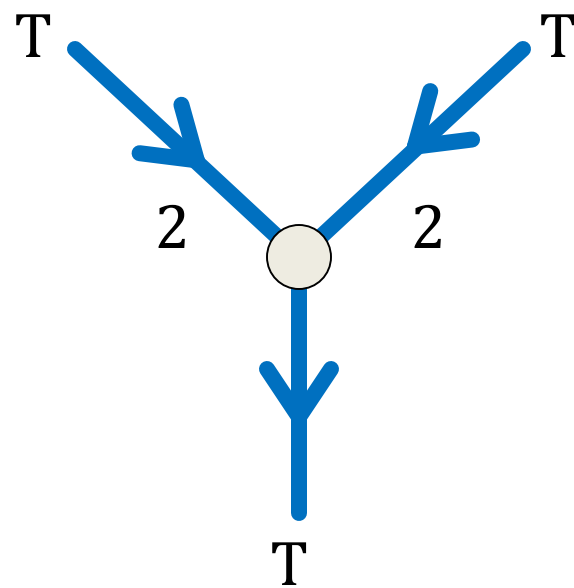
**Rule:** at least 2 units  
incoming at a vertex



# OR vertex



not your usual  
OR gate!



inputs

output

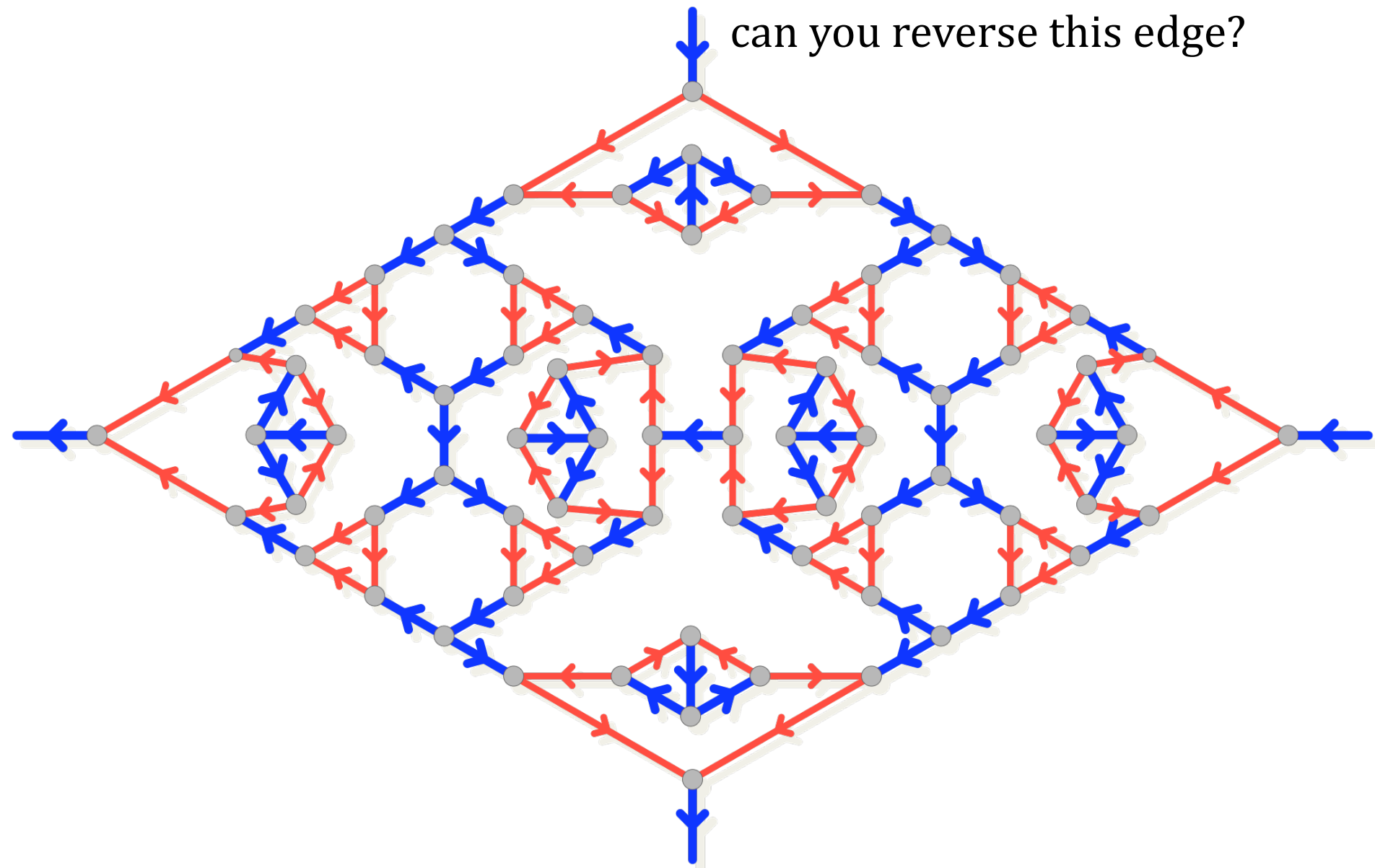
**Rule:** at least 2 units  
incoming at a vertex





# Decision Problem

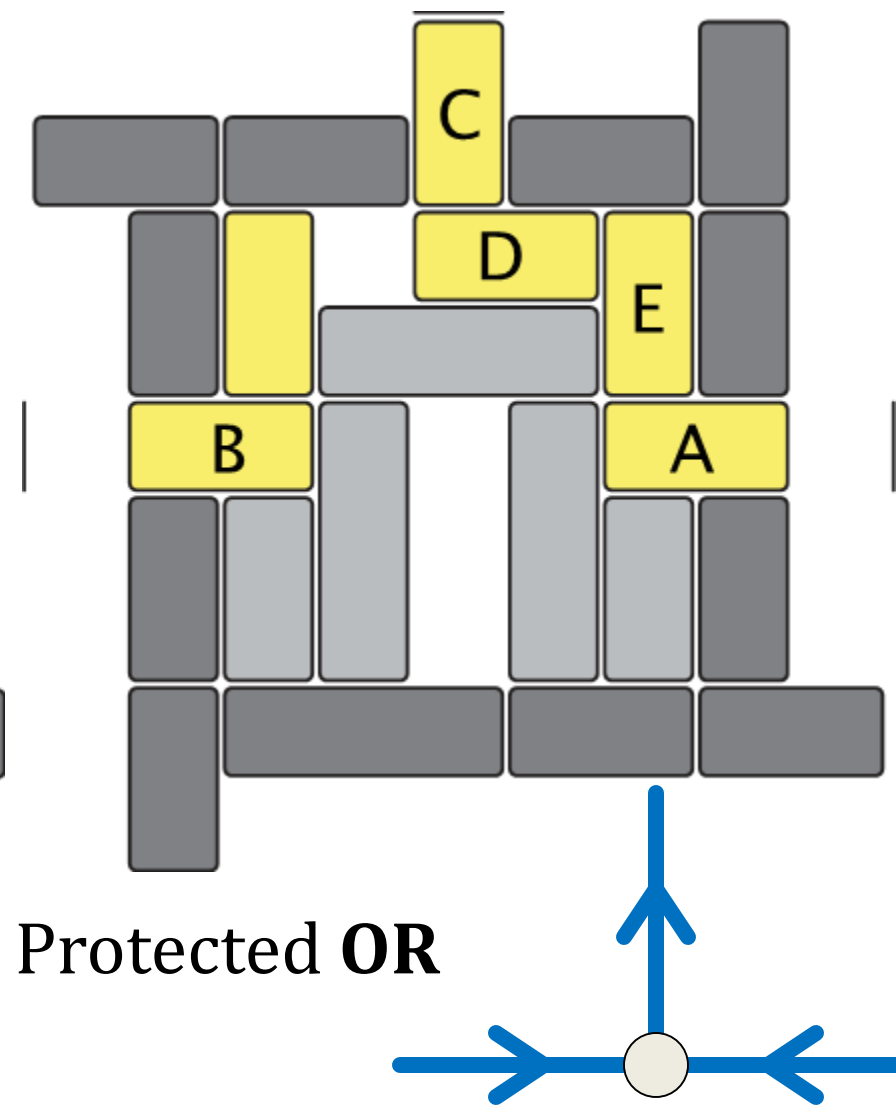
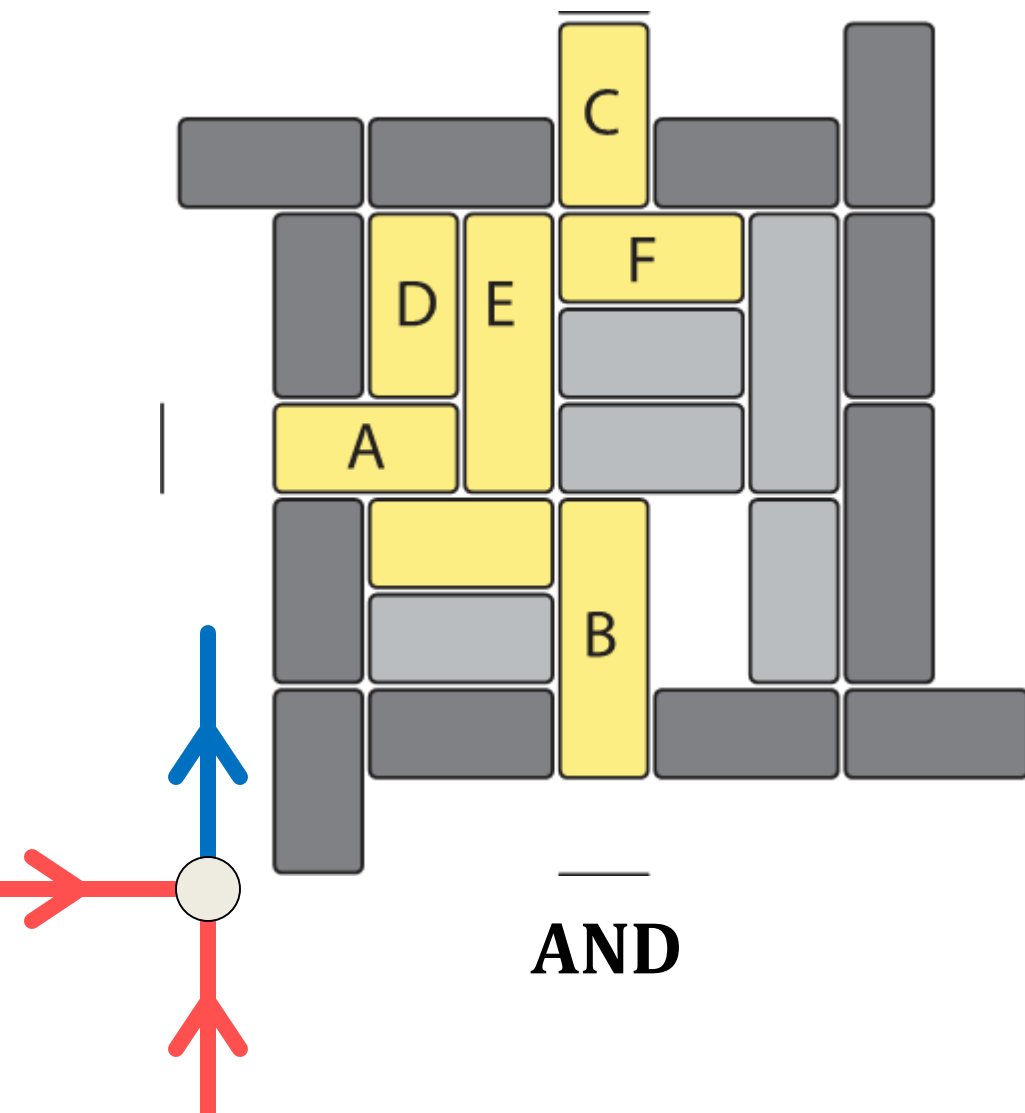
can you reverse this edge?





# Rush Hour is PSPACE-complete

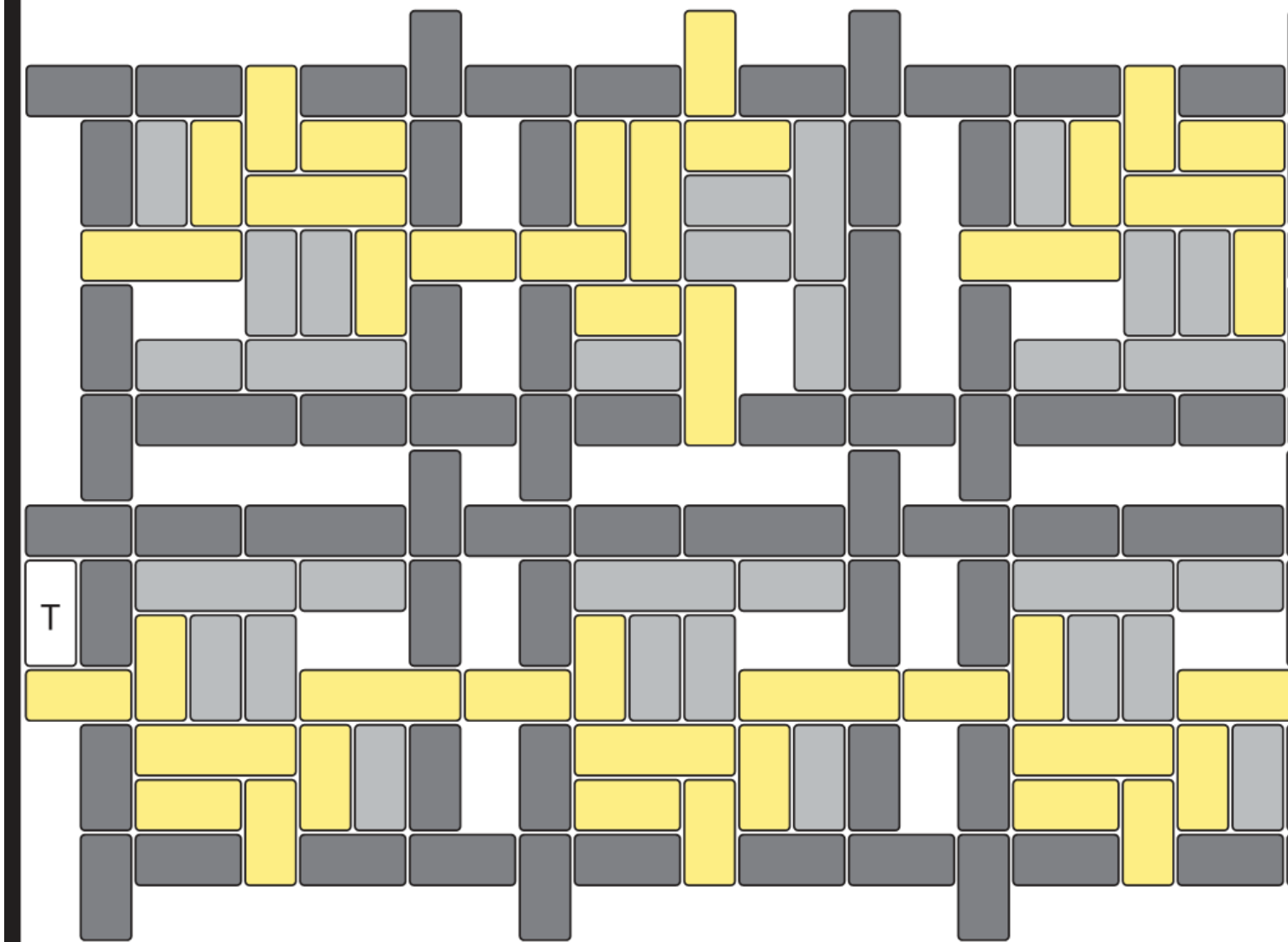
[Flake & Baum 2002; Hearn & Demaine 2002]





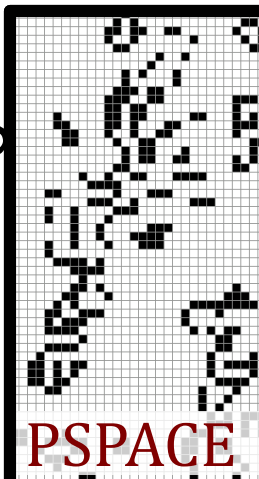
# Rush Hour is PSPACE-complete

[Flake & Baum 2002; Hearn & Demaine 2002]



# Complexity of Games & Puzzles

unbounded



PSPACE



PSPACE



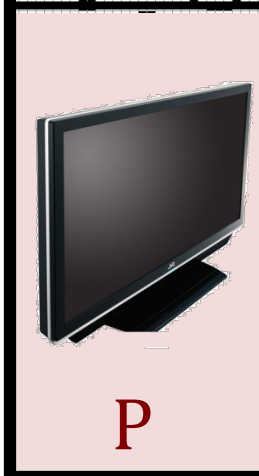
EXPTIME



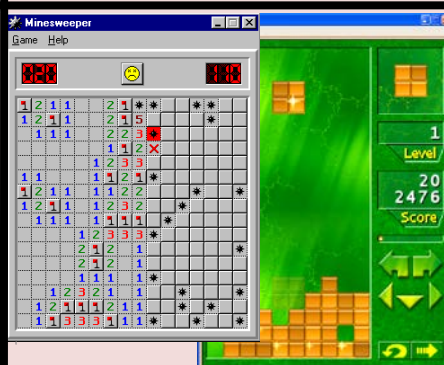
Rengo Kriegspiel?

Undecidable

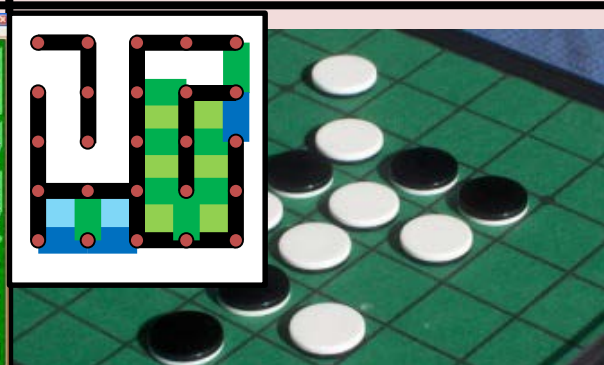
bounded



P



NP



PSPACE



bridge?

NEXPTIME

0 players  
(simulation)

1 player  
(puzzle)

2 players  
(game)

team,  
imperfect info

