

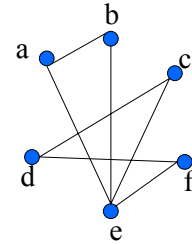


# Graphs



## Model Real World Problems

- Computer Networks
- Airline Connections
- Boston Road Map
- Program Flowchart
- Final Exam Conflicts
- Transistor Layouts
- Game Trees



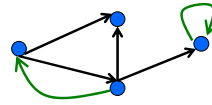
## Topics for this week

- Applications of Graphs
- Proving things about Graphs
  - Graph Coloring and Isomorphism
  - Connectivity and Paths
  - Trees, Tours

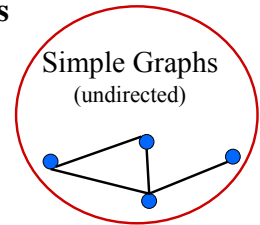


## Types of Graphs

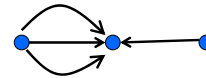
Directed Graphs



Simple Graphs  
(undirected)

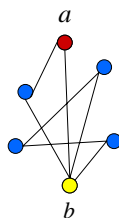


Multi-Graphs



## Definition

- A **Simple Graph** is a set of **vertices  $V$**  and a set of **edges  $E$**  where each edge is an unordered pair of distinct vertices  $a$  and  $b$ .

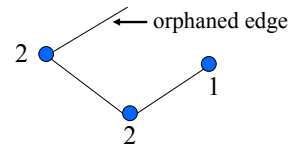


- **Degree** of a vertex  $v$  is the number of edges it connects to.

*Exercise:* Draw a graph of **3 vertices** with **degrees 2, 2 and 1**



## An Impossible Graph



$$\sum_{v \in V} \text{degree}(v) = 2 \times \text{number of edges}$$

- **Theorem:** Sum of degrees must be even
- **Handshaking Lemma**



## Problem : Airline Gate Allocation

### Airline Scheduling:

- Given a set of airline flights arriving at different times, **how many different gates** do I need in order to accommodate them?

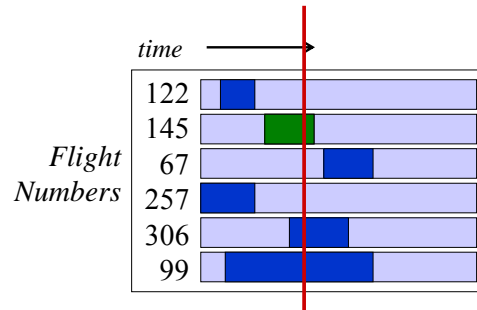


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## Airline Schedule

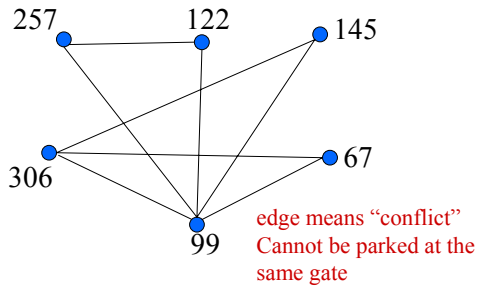


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## Model as a Graph



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## Solution: Graph Coloring

Color each vertex such that  
**no two adjacent vertices**  
have the **same color**

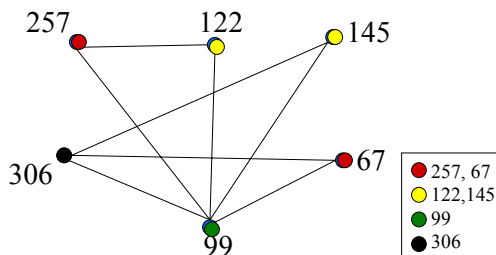


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## Assignment of Gates (or Colors)



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## Resource Allocation Problems

- How should I **schedule the final exams** to avoid two exams being scheduled at the same time if there are students taking both
- How many different **habitats** do I need to house several species of animals, some of which cannot coexist with others
- Frequency allocation: **How many different frequencies** do I need for overlapping radio broadcasts of wireless devices?

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## Chromatic Number

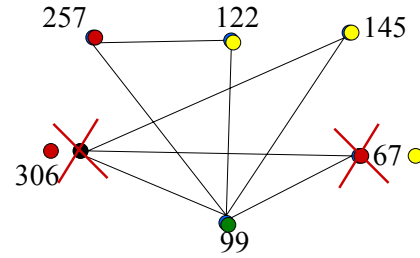
- *Question:* what is the minimum number of colors I need?  
– *Chromatic Number*
- How do I know that this is the minimum?

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## Assignment Using Fewer Gates



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## Class Problem 1

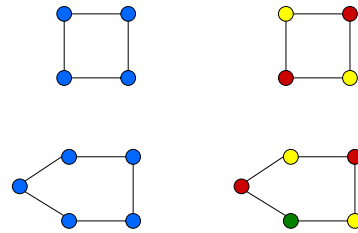


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## Simple Cycles $C_n$



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## Complete Graph $K_n$

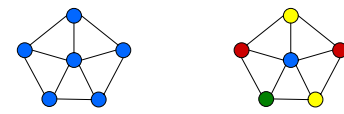


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## Wheel $W_n$



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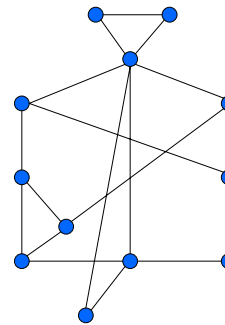


## Chromatic Number

- Odd cycles need 3 colors
- Complete graph needs exactly  $n$  colors
- Wheels can be colored with 4 colors
  - (if outer rim is even, then 3 colors)
- What about arbitrary graphs ?

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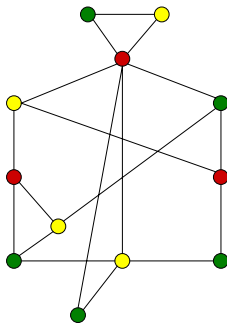
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Courtesy  
Steven Rudich  
CMU

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- Prove that it  
*can* be colored  
with 3 colors

- Prove that it  
*can not* be  
colored with 2 colors

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## Arbitrary Graphs

### Inspiration vs Perspiration

2-colorable?

Odd cycle exists: **NO**

No odd cycle: **YES!**

*Minimum color?*

– Exhausting idea...

3-colorable?

Brute-force approach,  
test all possible  
assignments

$3^n$

Can't do much better!

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## Brute Force vs Inspiration

Scheme saved on Sunday November 21, 1993 at 9:15:23 PM

Release 7.3.1

Microcode 11.146

Runtime 14.166

1 |=> (expt 3 2)

;Value: 9

1 |=> (expt 3 10)

;Value: 59049

1 |=> (expt 3 20)

;Value: 3486784401

1 |=> (expt 3 100)

;Value: 515377520732011331036461129765621272702107522001

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## Graph Coloring Theorems

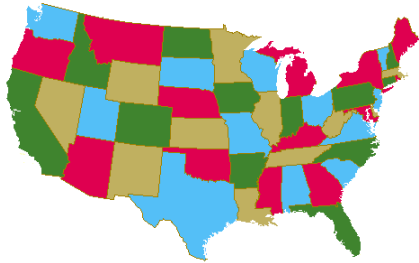
- A graph is 2-colorable iff there are no odd cycles.
- Complete graph  $K_n$  requires  $n$  colors
- If the maximum degree is  $d_{max}$ , then the graph can be colored with  $(d_{max} + 1)$  colors
- All **planar graphs** can be colored with **4 colors**

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## Four Color Theorem

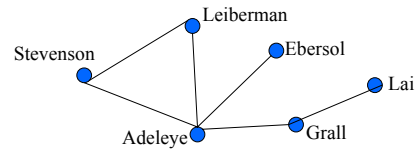


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## Who Knows Who Graph



### Six Degrees of Separation

– Is there a **path of length  $\leq 6$**  between any two nodes?

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## Oracle of Bacon (a.k.a. Six Degrees of Bacon)

(University of Virginia)



Total number of linkable actors: 535914

Bacon Number	# of People
0	1
1	1620
2	124832
3	326664
4	76190
5	5825
6	652
7	104
8	24
9	1
10	1

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## Class Problems

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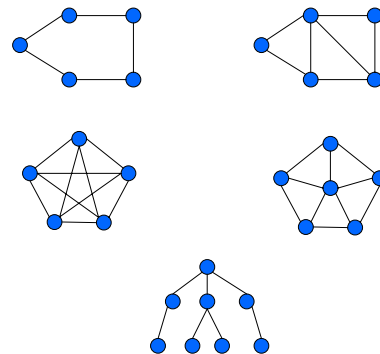


## Coloring with $d_{max}$ colors

- Induction Hypothesis  
 $P(n)$  = a graph with  $n$  vertices and maximum degree  $d_{max}$  can be colored with  $d_{max} + 1$  colors
- Base Case
- Inductive Step
  - Given an  $n+1$  graph, remove one vertex
  - Remaining  $n$  graph is colorable in  $d_{max} + 1$  colors
  - (why? Need to prove that the  $n$  graph has  $d_{max}$  degree)
  - Add vertex back in, must be one left over color
  - (why?)


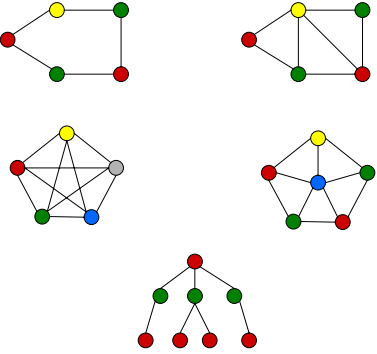
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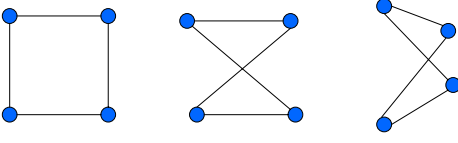



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## Graph Isomorphism

- Topology, not Geometry

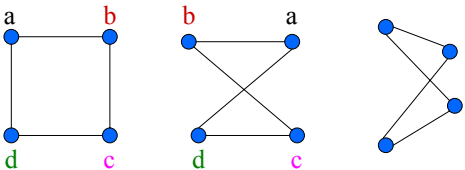


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## Graph Isomorphism

- Bijection between G1 and G2 nodes that preserves the adjacency list




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## Graph Isomorphism

- Isomorphism is an **equivalence relation** on the set of  $n$ -vertex Graphs!
- Can test for some invariants but otherwise another hard problem —  $n!$  possible mappings

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## Representations

List  $((v_1, v_1) (v_1, v_3) (v_3, v_3))$

Matrix

	$v_1$	$v_2$	$v_3$
$v_1$	1	0	1
$v_2$	0	0	0
$v_3$	0	0	1

Adjacency Lists

$v_1 \rightarrow v_1, v_3$   
 $v_3 \rightarrow v_3$

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